ISO/IEC JTC 1/SC 32 N 1246

Date: 2005-04-01

REPLACES: --

corrected

ISO/IEC JTC 1/SC 32

Data Management and Interchange

Secretariat: United States of America (ANSI) Administered by Farance, Inc. on behalf of ANSI

DOCUMENT TYPE	Summary of Voting/Table of Replies
TITLE	Summary of Voting/Table of Replies for 32N1199 - ISO/IEC CD 9075-02 Information technology Database Languages - SQL - Part 2: Foundation (SQL/Foundation)
SOURCE	SC 32 Secretariat
PROJECT NUMBER	1.32.03.06.02.00
STATUS	WG 3 should take and resolve the comments . corrected 2005-04-01
REFERENCES	
ACTION ID.	ACT
REQUESTED ACTION	
DUE DATE	
Number of Pages	130
LANGUAGE USED	English
DISTRIBUTION	P & L Members
	SC Chair
	WG Conveners and Secretaries

Douglas Mann, Secretary, ISO/IEC JTC 1/SC 32

Farance, Inc *, 360 Pelissier Lake Road, Marquette, MI, United States of America

Telephone: +1 906-249-9275; Facsimile; E-mail: MannD@battelle.org

available from the JTC 1/SC 32 WebSite http://staging.jtc1sc32.org/

*Farance, Inc. administers the ISO/IEC JTC 1/SC 32 Secretariat on behalf of ANSI

ISO/IEC JTC 1/SC 32 N1246

Summary of Voting on Document SC 32 N 1199,

Title: ISO/IEC CD 9075-02 Information technology -- Database Languages - SQL - Part 2: Foundation (SQL/Foundation)

"P" Member	Approval	Approval with Comments	Disapproval	Abstention
Australia			X	
Belgium				
Brazil				
Canada			X	
China	X			
Czech Republic	X			
Egypt				
Finland				
Germany		X		
Italy				Х
Japan			X	
Korea, Republic of	X			
Netherlands, The			X	
Norway				
Portugal				
Sweden	Х			
United Kingdom			X	
United States			X	
Total "P"	4	1	6	1
"O" Member				
Austria				
Denmark				
France				
Russian Federation				
Switzerland				
Total "O"				

ITALY

Lack of Experts

Template for comments and secretariat observations

Date: 2005-03-14 Document: **32N1199 9075-2 Foundation**

1	2	(3)	4	5	(6)	(7)
MB ¹	Clause No./ Subclause No./ Annex (e.g. 3.1)	Paragraph/ Figure/Table/ Note (e.g. Table 1)	Type of com- ment ²	Comment (justification for change) by the MB	Proposed change by the MB	Secretariat observations on each comment submitted
AU	N/a	N/a	ge	Address defects in Annex G the Defect Reports in the annex titled Defect reports not addressed in this edition of this part of ISO/IEC 9075 be addressed for the following: 32N1199 9075-2 Foundation - Address defects in annex G		

1 MB = Member body (enter the ISO 3166 two-letter country code, e.g. CN for China; comments from the ISO/CS editing unit are identified by **)

2 Type of comment: ge = general te = technical ed = editorial

NOTE Columns 1, 2, 4, 5 are compulsory.

National Body CAN Comments — 2005-02-03

32N1198, ISO/IEC CD 9075-01 Information technology - Database Languages - SQL - Part 1: Framework (SQL/Framework) 32N1199, ISO/IEC CD 9075-02 Information technology - Database Languages - SQL - Part 2: Foundation (SQL/Foundation) 32N1201, ISO/IEC CD 9075-03 Information technology - Database Languages - SQL - Part 3: Call-Level Interface (SQL/CLI) 32N1202, ISO/IEC CD 9075-04 Information technology - Database Languages - SQL - Part 4: Persistent Stored Modules (SQL/PSM) 32N1203, ISO/IEC CD 9075-09 Information technology - Database Languages - SQL - Part 9: Management of External Data (SQL/MED) 32N1204, ISO/IEC CD 9075-10 Information technology - Database Languages - SQL - Part 10: Object language bindings (SQL/OLB) 32N1205, ISO/IEC CD 9075-11 Information technology - Database Languages - SQL - Part 11: Information and Definition Schemas (SQL/Schemata)

32N1206, ISO/IEC CD 9075-13 Information technology - Database Languages - SQL - Part 13: SQL Routines and Types Using the Java™ Programming Language(SQL/JRT)

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
				CI) SQL/Framework	
	CAN-P01-001		1-Major Technical	P01-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved. Solution None provided with comment.	
				CI	OSQL/Foundation	
	CAN-P02-001		1-Major Technical	P02-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved. Solution	
					CD SOL/CLI	
	CAN-P03-001		1-Major Technical	P03-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved. Solution None provided with comment.	
					CD SQL/PSM	
	CAN-P04-001		1-Major Technical	P04-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					Solution	
					None provided with comment.	
					CD SQL/MED	
	CAN-P09-001		1-Major	P09-No specific	All Possible Problems and Editor's Notes must be satisfactorily resolved and all	
			Technical	location	problems discovered during the course of the ballot resolution process must be	
					satisfactorily resolved.	
					Solution	
					None provided with comment.	
					CD SQL/OLB	
	CAN-P10-001		1-Major	P10-No specific	All Possible Problems and Editor's Notes must be satisfactorily resolved and all	
			Technical	location	problems discovered during the course of the ballot resolution process must be	
					satisfactorily resolved.	
					Solution	
				<u> </u>	D SQL/Schemata	
	CAN-P11-001		1-Major	P11-No specific	All Possible Problems and Editor's Notes must be satisfactorily resolved and all	
			Technical	location	problems discovered during the course of the ballot resolution process must be	
					satisfactorily resolved.	
					Solution None provided with comment	
					CD SQL/JRT	
	CAN-P13-001		1-Major	P13-No specific	All Possible Problems and Editor's Notes must be satisfactorily resolved and all	
			Technical	location	problems discovered during the course of the ballot resolution process must be	
					Solution	
					Solution None provided with comment	
					None provided with comment.	

ISO/IEC JTC1/SC32/WG3:TXL-034 – DIN NI-32 N 0642

Authoritative Version: Adobe Acrobat Portable Document Format (PDF)



ISO

International Organization for Standardization

DIN

Deutsches Institut für Normung

Din NI-32 Database ISO/IEC JTC 1/SC 32 Data Management and Interchange WG 3 Database Languages

Title: German Comments on SC32 N 1156: ISO/IEC CD 9075-1, 2, 3, 4, 9, 10, 11, 13:200x(E)

Status: Consolidated comments to assist with resolution of CD ballot comments

Author: Jörn Bartels (Germany)

National Body DEU Comments — 2005-03-08

SEQ #	Cmn t ID	See Als o	Severity	Reference	Description	Addressed By
		•		ISO/IEC	FCD 9075-01:200x(E) SQL/Framework	
1	DEU- P01- 010		1-Major Technical	P01-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved. Solution None provided with comment.	
				ISO/IEC	FCD 9075-02:200x(E) SOL/Foundation	
2	DEU- P02- 010		1-Major Technical	P02-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved. Solution None provided with comment.	
3	DEU- P02- 020		1-Major Technical	P02-07.6 Table reference	Queries of the form SELECT FROM <joined table=""> Do not seem to be supported anymore. This is due to changes proposed in DRS-077. Solution None provided with comment.</joined>	
				ISO/	IEC FCD 9075-03:200x(E) SQL/CLI	
4	DEU- P03- 010		1-Major Technical	P03-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved. Solution None provided with comment.	
				ISO/I	EC FCD 9075-04:200x(E) SQL/PSM	
5	DEU- P04- 010		1-Major Technical	P04-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved. Solution None provided with comment.	
6	DEU- P11- 020	DEU- P11- 030	1-Major Technical	<i>P04-18.2 "MODULE_PRIVILEGES"</i> <i>Table</i>	The table MODULE_PRIVILEGES stores the privileges granted on a specific module. The same information could be stored in the table USAGE_PRIVILEGES. This would simplify the definition schema and standardise the way, how privileges are stored. Solution None provided with comment.	
7	DEU- P11- 030	DEU- P11- 020	1-Major Technical	P04-18.2 "MODULE_PRIVILEGES" Table	The constraint MODULE_PRIVILEGE_GRANTOR_CHECK and MODULE_PRIVILEGE_GRANTEE_CHECK reference still the tables ROLES and USERS. They are gone!	

SEQ #	Cmn t ID	See Als 0	Severity	Reference	Description	Addressed By
					Solution	
					Replace the check constraints with a foreign key on AUTHORIZATIONS.	
				ISO/I	EC FCD 9075-09:200x(E) SQL/MED	
8	DEU-		1-Major	P09-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered	
	P09-		Technical		during the course of the ballot resolution process must be satisfactorily resolved.	
	010				Solution	
0	DEU		2 Minor	P00-25 Definition Schema	None provided with comment.	
9	P09-		Z-Minor Technical	1 09-25 Definition Schema	should be defined. This are at least 25.4 FOREIGN DATA WRAPPERS have table. 25.8	
	020		Teenneur		FOREIGN TABLES base table and 25.10 ROUTINE MAPPINGS base table	
					Solution	
					None provided with comment.	
10	DEU-		3-Major	P09-25.2	The constraint DATA_TYPE_DESCRIPTOR_DATA_TYPE_CHECK_COMBINATIONS is entirely	
	P09-		Editorial	DATA_TYPE_DESCRIPTOR" table	replaced. This leads to problems of desynchronisation with SQL/Schemata. It does also not allow	
	030				modifications from other parts (like SQL/XML) of the standard.	
					The constraint does also currently not check the NULL applicability of the columns, as described in	
					Description 2) which are inserted by this constraint	
					Solution	
					None provided with comment.	
11	DEU-		3-Major	P09-25.2	The Descriptions 2) and 3) are in conflict with each other. They describe both the NULLability of the	
	P09-		Editorial	DATA_TYPE_DESCRIPTOR" table	newly introduced columns. There is a conflict if both come to different results.	
	040				Solution	
					The Descriptions 2) and 3) should be merged.	
12	DEU-		2-Minor	P09-25.4	There is no constraint, which verifies the existence of the catalog and the authorization Identifier, which is	
	P09-		Technical	FOREIGN_DATA_WRAPPERS table	used.	
	030				Solution	
					None provided with comment	
13	DEU-		2-Minor	P09-25.6 "FOREIGN SERVERS" table	There is no constraint, which verifies the existence of the catalog and the authorization Identifier, which is	
15	P09-		Technical		used.	
	060					
					Solution	
					None provided with comment.	
14	DEU-		2-Minor	P09-25.12 "TABLES" table	There is no constraint, which verifies that for a FOREIGN table there is also an entry in the table	
	P09-		Technical		FOREIGN_TABLES. This could be done as it is done already in constraint	
	0'/0				TABLES_CHECK_NOT_VIEW of the table TABLES.	

SEQ #	Cmn t ID	See Als 0	Severity	Reference	Description	Addressed By
					Solution None provided with comment.	
15	DEU- P09- 080		2-Minor Technical	P09-25.13 "USAGE_PRIVILEGES" table	There is no constraint added, which checks the values of OBJECT_CATALOG and OBJECT_NAME as described in Description 1)	
					USAGE_PRIVILEGES_CHECK_REFERENCES_OBJECT which allows OBJECT_SCHEMA to be the empty string.	
					It is currently also not allowed, that the OBJECT_TYPE is anything except 'DOMAIN', 'CHARACTER SET', 'COLLATION', 'TRANSLATION', 'SEQUENCE'. The use for a foreign-data wrapper or a foreign server requires a modification of constraint USAGE_PRIVILEGES_OBJECT_TYPE_CHECK. Solution	
					None provided with comment.	
16	DEU-		2-Minor	P09-25.15 "USER_MAPPINGS" table	There is no foreign key check for the column AUTHORIZATION_IDENTIFIER.	
	P09-		Technical			
	090				Solution	
	1			ICO/I	$\mathbf{E} = \mathbf{E} \mathbf{C} \mathbf{E} \mathbf{C} \mathbf{D} \mathbf{O} \mathbf{O} \mathbf{F} \mathbf{C} \mathbf{E} \mathbf{O} \mathbf{C} \mathbf{O} \mathbf{C} \mathbf{D}$	
4 -	DEU		1.16			
17	DEU-		I-Major Technical	P10-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered during the course of the ballot resolution process must be satisfactorily resolved.	
	010		Technical		Solution	
	010				None provided with comment	
				ISO/IE/	C FCD 9075-11:200x(E) SOL/Schemata	
18	DEU-		1-Major	P11-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered	
10	P11-		Technical	·····	during the course of the ballot resolution process must be satisfactorily resolved.	
	010				Solution	
					None provided with comment.	
19	DEU-		1-Major	P11-No specific location	It is not clear, which tables should have a direct or indirect relationship to the table SCHEMATA. For	
	P11- 020		Technical		some tables is a foreign key defined, for some is a check constraint defined, which checks the foreign key relationship only when there are schemas in the same catalog.	
					There should be an explanation for this distinction and all relationships need to be checked for correctness.	
					An example for a dubious relationship is the constraint TRIGGERS_REFERENCES_TABLES. Is it really possible, to define a trigger on a table of another catalog. This is especially strange, as the table	

SEQ #	Cmn t ID	See Als 0	Severity	Reference	Description	Addressed By
					TRIGGERED_UPDATE_COLUMNS has a direct foreign key to COLUMNS. Solution None provided with comment.	
20	DEU- P11- 030		3-Major Editorial	P11-5.54 Short name views	The View definitions in 5.54 Short name views should be sorted according to the order of the base views (i.E. position of CONSTR_COL_USAGE). Solution Order them according to base view order	
21	DEU- P11- 040		2-Minor Technical	P11-5.78 "SQL_LANGUAGES" View	The View SQL_LANGUAGES is depricated. In Note 6 is a reference to SQL/Framework Subclause 6.4. There is a backwards reference, in Note 9, which says "The equivalent information is available to the SQL user in the Information Schema." When we delete the view, it is not clear if the Note 9 should also be deleted. Solution None provided with comment.	
22	DEU- P11- 050		2-Minor Technical	P11-6.9 "CHARACTER_ENCODING_FORMS" Table	The NOT NULL Constraints are not needed, as all columns are part of the primary key. Solution Delete the NOT NULL constraints.	
23	DEU- P11- 060		2-Minor Technical	P11-6.10 "CHARACTER_REPERTOIRES" Table	The NOT NULL Constraints on the column CHARACTER_REPERTOIRE_NAME is not needed, as the column is part of the primary key. Solution Delete the NOT NULL constraint.	
24	DEU- P11- 070		4-Minor Editorial	P11-6.11 "CHARACTER_SETS" Table	The column NUMBER_OF_CHARACTERS is in the last Edition of the Standard depricated and should now be deleted. Solution Delete the column. Do the same in the View Definition 5.12 CHARACTER_SETS view and in 5.78 Short name views in the view CHARACTER_SETS_S. Delete the corresponding List Elements 1) and 2) in Annex C.	
25	DEU- P11- 080		2-Minor Technical	P11-6.16 "COLLATIONS" Table	There is no constraint for the column CHARACTER_REPERTOIRE_NAME defined. It needs to reference the Table CHARACTER_REPERTOIRES. Solution Add the constraint COLLATIONS_FOREIGN_KEY_CHARACTER_REPERTOIRES FOREIGN KEY (CHARACTER_REPERTOIRES) REFERENCES CHARACTER_REPERTOIRES.	
26	DEU- P11- 090		4-Minor Editorial	P11-6.16 "COLLATIONS" Table	The columns COLLATION_TYPE, COLLATION_DICTIONARY, and COLLATION_DEFINITION are in the last edition of the Standard depricated and should now be deleted. Solution Delete the columns. Do the same in the View Definition 5.15 COLLATIONS view and in 5.78 Short name views in the view COLLATIONS_S. Delete the corresponding List Elements 3) and 4) in Annex C.	

SEQ #	Cmn t	See Als	Severity	Reference	Description	Addressed By
	ID	0				
27	DEU- P11-		2-Minor Technical	P11-6.16 "COLLATIONS" Table	The column PAD_ATTRIBUTE has no NOT NULL check constraint, even that there is in the description	
	100		reennear		Solution	
	100				None provided with comment	
20	DEU		2 Minor	P11-6 20 "COLUMNS" Table	For the columns IS GENERATED and IDENTITY GENERATION are no check constraints specified	
20	P11-		Technical		but in the description are Lists of allowed values	
	110		reenneur		Solution	
					Add to the column IS GENERATED the following column level constraint:	
					CONSTRAINT COLUMNS IS GENERATED CHECK CHECK (IS GENERATED in ('NEVER'	
					'ALWAYS'))	
					Add to the column IDENTITY GENERATION the following column level constraint:	
					CONSTRAINT COLUMNS IDENTITY GENERATION CHECK CHECK	
					(IDENTITY_GENERATION IN ('ALWAYS', 'BY DEFAULT'))	
29	DEU-		2-Minor	P11-6.21	The constraint DATA_TYPE_DESCRIPTOR_FOREIGN_KEY_SCHEMATA assures that the values of	
	P11-		Technical	"DATA_TYPE_DESCRIPTORS" Table	USER_DEFINED_TYPE_CATALOG and USER_DEFINED_TYPE_SCHEMA have corresponding	
	120				rows in the table SCHEMATA. The constraint	
					DATA_TYPE_DESCRIPTOR_CHECK_REFERENCES_UDT allows that the value for the column	
					USER_DEFINED_TYPE_CATALOG has no corresponding row in SCHEMATA. As this is not possible	
					according to the first constraint, we could rewrite this constraint as a foreign key.	
					It is not clear, if this is intended.	
					Solution	
					None provided with comment.	
30	DEU-		2-Minor	P11-6.21	The columns SCOPE_CATALOG, SCOPE_SCHEMA, and SCOPE_NAME are not checked against the	
	P11-		Technical	DATA_TIPE_DESCRIPTORS Table	possible values in the table TABLES.	
	130				Solution	
					None provided with comment.	
31	DEU-		4-Minor	P11-6.21 "DATA TYPE DESCRIPTORS" Table	The constraint DATA_TYPE_DESCRIPTOR_CHECK_OBJECT_TYPE should be a column constraint, as	
	P11-		Editorial	DATA_TIPE_DESCRIPTORS Table	it references only the column OBJECT_TYPE.	
	140				Solution	
					Remove the preceding comma.	
32	DEU-		4-Minor	P11-6.21	In the constraint	
	P11-		Editorial	DATA_TTPE_DESCRIPTORS Table	DATA_TYPE_DESCRIPTOR_CHECK_REFERENCES_COLLATION_CHARACTER_SET_APPLICA	
	150				BILLITY should be a comma "," at the end of the 10 th line of the constraint.	
					Solution	
					Add the missing comma.	
33	DEU-		2-Minor	P11-6.24 "DOMAIN_CONSTRAINTS"	Should there be a NOT NULL check constraint for the columns IS_DEFERRABLE and	
	P11-		Technical	Table	INITIALLY_DEFERRED?	

SEQ #	Cmn t ID	See Als 0	Severity	Reference	Description	Addressed By
	160				In the description is no explanation of the meaning of a possible null value. Solution	
34	DEU- P11- 170		2-Minor Technical	P11-6.31 "PARAMETERS" Table	The foreign key constraint PARAMETERS_FOREIGN_KEY_SCHEMATA does not check the name of the routine. This constraint should be removed and instead there should be a foreign key to ROUTINES be defined. Solution	
35	DEU- P11- 180		2-Minor Technical	P11-6.31 "PARAMETERS" Table	None provided with comment. There needs to be a unique constraint defined, which guarantees the uniqueness of a parameter name for a routine. Solution Add the constraint: CONSTRAINT PARAMETERS_UNIQUE_CHECK (UNIQUE SPECIFIC_CATALOG, SPECIFIC_SCHEMA, SPECIFIC_NAME PARAMETER, NAME)	
36	DEU- P11- 190		2-Minor Technical	P11-6.31 "PARAMETERS" Table	For the columns FROM_SQL_SPECIFIC_CATALOG, FROM_SQL_SPECIFIC_SCHEMA, and FROM_SQL_SPECIFIC_ NAME and TO_SQL_SPECIFIC_CATALOG, TO_SQL_SPECIFIC_SCHEMA, and TO_SQL_SPECIFIC_NAME is no foreign key check defined. Solution	
37	DEU- P11- 200		1-Major Technical	P11-6.36 "ROUTINE_PRIVILEGES" Table	The table ROUTINE_PRIVILEGES stores the privileges granted on a specific routine. The same information could be stored in the table USAGE_PRIVILEGES. This would simplify the definition schema and standardise the way, how privileges are stored. Solution None provided with comment	
38	DEU- P11- 210		2-Minor Technical	P11-6.41 "SCHEMATA" Table	For the columns DEFAULT_CHARACTER_SET_CATALOG, DEFAULT_CHARACTER_SET_SCHEMA and DEFAULT_CHARACTER_SET_NAME is a foreign key referencing the table CHARACTER_SETS missing. Solution Add the missing Foreign Key constraint: CONSTRAINT SCHEMATA_FOREIGN_KEY_CHARACTER_SETS FOREIGN KEY (DEFAULT_CHARACTER_SET_CATALOG, DEFAULT_CHARACTER_SET_SCHEMA, DEFAULT_CHARACTER_SET_NAME) REFERENCES CHARACTER_SETS	
39	DEU- P11-		4-Minor Editorial	P11-6.48 "TABLE_CONSTRAINTS" Table	The constraint TABLE_CONSTRAINTS_UNIQUE_CHECK is not needed, as the uniqueness of the constraint name is already checked by the assertion UNIQUE_CONSTRAINT_NAME in subclause 6.4	

SEQ #	Cmn t ID	See Als 0	Severity	Reference	Description	Addressed By
	220				Solution	
					A possible solution is to remove the superflouus constraint.	
40	DEU-		4-Minor	P11-6.50 "TABLE_PRIVILEGES"	In the constraint TABLE_PRIVILEGES_TYPE_CHECK is the last element of the inlist misspelled. It	
	P11-		Editorial	Table	should be <u>REFERENCES</u> instead of EFERENCES	
	230				Solution	
					Fix the typo.	
41	DEU-		4-Minor	P11-6.54	The constraint TRIGGERED_UPDATE_COLUMNS_FOREIGN_KEY_TRIGGERS is not needed, as a	
	P11-		Editorial	"TRIGGERED_UPDATE_COLUMNS"	more restrictive relationship is already guaranteed by constraint	
	240			Tuble	TRIGGERED_UPDATE_COLUMNS_EVENT_MANIPULATION_CHECK.	
					Solution	
					A possible solution is to remove the superflouus constraint.	
42	DEU-		2-Minor	P11-6.55	The table TRIGGER_COLUMN_USAGE should have a foreign Key to the table	
	P11-		Technical	TRIGGER_COLUMIN_USAGE Table	TRIGGER_TABLE_USAGE, and not to TRIGGERS.	
	250				Solution	
					Add the following constraint:	
					TRIGGER_COLUMN_USAGE_FOREIGN_KEY_TRIGGER_TABLE_USAGE	
					FOREIGN KEY	
					PEEERENCES TRIGGER TARLE USAGE	
					It might be possible to remove the constraint	
					TRIGGER COLUMN USAGE FOREIGN KEY TRIGGERS.	
43	DEU-		2-Minor	P11-6.62 "USER DEFINED TYPES"	In the last query of the constraint USER DEFINED TYPES CHECK SOURCE TYPE is the column	
	P11-		Technical	Table	OBJECT TYPE not in the reference List of the IN clause.	
	260				Solution	
					None provided with comment.	
44	DEU-		2-Minor	P11-Appendix C 6)	The columns FEATURE_ID and FEATURE_NAME of the view SQL_PACKAGES are in the last Edition	
	P11-		Technical		of the Standard depricated and should now be deleted. But without these columns does the view not	
	270				provide any usefull information. Should the entire View be deleted?	
					Solution	
					None provided with comment.	
				ISO/I	EC FCD 9075-13:200x(E) SQL/JRT	
45	DEU-		1-Major	P13-No specific location	All Possible Problems and Editor's Notes must be satisfactorily resolved and all problems discovered	
	P13-		Technical		during the course of the ballot resolution process must be satisfactorily resolved.	
	010				Solution	
					None provided with comment.	

END OF PAPER

3 March, 2005



ISO

International Organization for Standardization

ISO/IEC JTC 1/SC 32 **Data Management and Interchange** WG 3 Database Languages

- Title: Japan Ballot Comments on CD 9075:2007(E)
- **Status:** Document to accompany ballot response
- Author: Masashi Tsuchida , Takaaki Shiratori, Takashi Kotera
- Abstract: We present the comments of Japan on the CD ballot documents, to accompany our ballot response on that document.

References: [1] WG3:TXL-002 = 32N1198, ISO/IEC CD 9075-1, Information technology. Database languages SQL. Part 1: Framework (SQL/Framework)] [2] WG3:TXL-003 = 32N1199, ISO/IEC CD 9075-2, Information technology. Database languages. SQL. Part 2: Foundation (SOL/Foundation) [3] WG3:TXL-004 = 32N1201, ISO/IEC CD 9075-3, Information technology. Database languages. SQL. Part 3: Call-Level Interface (SQL/CLI) [4] WG3:TXL-005 = 32N1202, ISO/IEC CD 9075-4, Information technology . Database languages . SQL . Part 4: Persistent Stored Modules (SQL/PSM) [5] WG3:TXL-006 = 32N1203, ISO/IEC CD 9075-9, Information technology . Database languages . SQL . Part 9: Management of

External Data (SQL/MED)

[6] WG3:TXL-007 = 32N1204, ISO/IEC CD 9075-10, Information technology . Database languages . SQL . Part 10: Object language bindings (SQL/OLB)

[7] WG3:TXL-008 = 32N1205, ISO/IEC CD 9075-11, Information technology . Database languages . SQL . Part 11: Information and Definition Schemas (SQL/Schemata)

[8] WG3:TXL-009 = 32N1206, ISO/IEC CD 9075-11, Information technology . Database languages . SQL . Part 13: Information and Definition Schemas (SQL/JRT)

SEQ	Cmnt	See				Addressed		
#	ID	Also	Severity	Reference	Description	Ву		
	SQL/Framework							
001	JPN-P01-001		1-Major	P01-No specific	There are quite a few features to discuss a emerging next standard. Japan thinks			
			Technical	location	that we should take enough time to add new features.			
					Solution			
					Solution			
					None provided with comment.			
				S	SQL/Foundation			
002	JPN-P02-002		1-Major	P02-No specific	There are quite a few features to discuss a emerging next standard. Japan thinks			
			Technical	location	that we should take enough time to add new features.			
					Solution			
					Solution			
					None provided with comment.			
003	JPN-P02-003		1-Major	P02-11.3,	It is allowed that which is <as clause="" subquery=""> with</as>			
			Technical	definition>	be materialized at table definition			
					Solution			
					None movided with commont			
004	JPN-P02-004		1-Maior	P02-14.8.	An INSERT statement has no different effects on identity columns specified			
004			Technical	<insert< th=""><th>GENERATED ALWAYS and that specified GENERATED BY DEFAULT.</th><th></th></insert<>	GENERATED ALWAYS and that specified GENERATED BY DEFAULT.			
				statement>				
					Solution			
					None provided with comment.			
			1		SOL/CLI			
005	JPN-P03-001		1-Major	P03-No specific	There are quite a few features to discuss a emerging next standard. Japan thinks			
			Technical	location	that we should take enough time to add new features.			
					Solution			
					50101011			
					None provided with comment.			

SEQ	Cmnt	See	G	Deferre	Description	Addressed	
#	ID	Also	Severity	Reference	Description	Ву	
SQL/PSM							
006	JPN-P04-001		1-Major	P04-No specific	There are quite a few features to discuss a emerging next standard. Japan thinks		
			Technical	location	that we should take enough time to add new features.		
					Solution		
					None provided with comment.	l	
	I		L.		SOL/MED		
007	JPN-P05-001		1-Major	P05-No specific	There are quite a few features to discuss a emerging next standard. Japan thinks		
			Technical	location	that we should take enough time to add new features.		
					Solution		
					None provided with comment	l	
	I		1		SOL/OLB		
008	JPN-P10-001		1-Major	P10-No specific	There are quite a few features to discuss a emerging next standard. Japan thinks		
000			Technical	location	that we should take enough time to add new features.		
						l	
					Solution		
					None provided with comment	l	
					SOI /Schome		
	IDN D11 001		1 1 1	D11 N			
009	JPN-P11-001		1-Major Technical	P11-No specific	There are quite a few features to discuss a emerging next standard. Japan thinks that we should take enough time to add new features		
			reennear	loculion	that we should take chough this to add new readeres.		
					Solution	l	
					None provided with comment.		
	1				SQL/JRT		
009	JPN-P13-001		1-Major	P13-No specific	There are quite a few features to discuss a emerging next standard. Japan thinks		
			Technical	location	that we should take enough time to add new features.	l	
					Solution	l	
					Southon	l	
					None provided with comment.	l	

ISO/IEC JTC1/SC32/WG3 TXL-031 2004-01-23

ISO International Organization for Standardization ISO/IEC JTC 1/SC 32 Data Management and Interchange WG 3 Database Languages

Project: ISO: 1.32.3.5

Title: Ballot Comment on ISO/IEC CD 9075-1, -2, -3, -4, -9, -10, -11, and -13

Status: Netherlands National Body Comments

Author: Stephen Cannan (Editor)

References:

SC32 N01198, CD 9075-1 Information Technology - Database Language SQL - Part 1: Framework (SQL/Framework) Jim Melton (Editor), December, 2004.
 SC32 N01199, CD 9075-2 Information Technology - Database Language SQL - Part 2: Foundation (SQL/Foundation) Jim Melton (Editor), December, 2004.
 SC32 N01201, CD 9075-3 Information Technology - Database Language SQL - Part 3: Call Level Interface (SQL/CLI) Jim Melton (Editor), December, 2004.
 SC32 N01202, CD 9075-4 Information Technology - Database Language SQL - Part 4: Persistent Stored Modules (SQL/PSM) Jim Melton (Editor), December, 2004.
 SC32 N01203, CD 9075-9 Information Technology - Database Language SQL - Part 9: Management of External Data (SQL/MED) Jim Melton (Editor), December, 2004.
 SC32 N01204, CD 9075-10 Information Technology - Database Language SQL - Part 10: Object Language Bindings (SQL/OLB) Jim Melton (Editor), December, 2004.
 SC32 N01205, CD 9075-11 Information Technology - Database Language SQL - Part 11: Schemata (SQL/Schemata) Jim Melton (Editor), December, 2004.
 SC32 N01206, CD 9075-13 Information Technology - Database Language SQL - Part 11: Schemata (SQL/Schemata) Jim Melton (Editor), December, 2004.

The Netherlands vote is:

SQL/Framework No with comments SQL/Foundation No with comments SQL/CLI No with comments SQL/PSM No with comments SQL/MED No with comments SQL/OLB No with comments SQL/Schemata No with comments SQL/JRT Yes with comments

If all problems and technical errors, i.e. those identified in this ballot, and those identified during the editing meeting(s), are resolved to our satisfaction, then the Netherlands will change its NO votes to YES votes.

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
				5	SOL/Framework	
	NLD-P01-001		3-Major	P01-04.04, SQL	FRM-002 The following Language Opportunity has been noted:	
			Editorial	data types	Source: DBL:BBN-167/X3H2-98-386	
					Language Opportunity:	
					Section needs a better organization	
					There should be a section called SQL Data Types. Then a short definition of	
					what is meant by an SQL	
					data type. Then the list of the five types of data types (predefined, row type,	
					user-defined type, collection	
					type, and reference type). Then there should be a definition for each.	
					Solution	
					None provided with comment.	
	NLD-P01-002		2-Minor	P01-06.03.03.03,	WG3-P01-001	
			Technical	Rule evaluation	The referenced subclause includes the following text:	
				oraci	In general, if some syntactic element contains more than one other syntactic	
					element, then the General Rules for contained elements that appear earlier in the	
					Production for the containing syntactic element are applied before the General	
					For example, in the production:	
					<a> ::= <c></c>	
					the Syntax Rules and Access Rules for <a>,,and <c>are effectively applied</c>	
					simultaneously. The General Rules for are applied before the General Rules	
					for <c>, and the General Rules for <a>are applied after the General Rules for</c>	
					both and <c>.</c>	
					In SQL/Foundation, Subclause 13.5, " <sql procedure="" statement="">", is a clear</sql>	
					exception to this general rule for General Rules, for the GRs of the particular	
					contained statement (e.g., an <insert statement="">) are clearly intended to be</insert>	
					invoked only when a GR in Subclause 13.5 explicitly states that the contained	
					statement to be executed.	
					Now, it might be that the introductory words, in general, can be taken to imply that there are some executions, but in that area shouldn't the executions be	
					explicitly mentioned?	
					Solution	
					None provided with comment.	
					SOL/Foundation	
	NLD-P02-001		1-Major	P02-04.32.01	FND-975 The following Possible Problem has been noted:	
	102 001		Technical	General	Source: $WG3:SIA-030 = H2-2004-222$	
				description of	Possible Problem:	
				cursors	Subclause 4.32.1, "General description of cursors", contains:	
					For every <declare cursor=""> in [emphasis added] an SQL-client module, a cursor</declare>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
					is effectively created when an SQL-transaction (see Subclause 4.35, "SQL-	
					transactions") referencing the SQL-client module	
					is initiated.	
					For every <dynamic cursor="" declare=""> in an <sql-client definition="" module="">, a</sql-client></dynamic>	
					cursor is effectively created when an SQL-transaction (see Subclause 4.35,	
					"SQL-transactions") referencing the <sqlclient definition="" module=""> is initiated.</sqlclient>	
					An extended dynamic cursor is also [emphasis added] effectively created when	
					an <allocate cursor="" statement=""> is executed within an SQL-session and destroyed</allocate>	
					when that SQL-session is terminated.	
					This text suffers from several problems, all of which probably need to be	
					addressed at the same time:	
					1) The first paragraph entertains the notion of a piece of SQL syntax appearing	
					inside something that is not a piece of SQL syntax. It seems that either " <declare< td=""><td></td></declare<>	
					cursor>" should be replaced by "cursor", or "SQL-client module" should be	
					replaced by " <sql-client definition="" module="">". In either case there would be</sql-client>	
					knock-on effects on the remaining text. Note that the second paragraph prefers	
					to talk about syntactic containment exclusively, but its text is too suspect for it	
					to be used as a guideline for correcting the first paragraph.	
					2) The first paragraph entertains the notion of an SQL-transaction referencing an	
					SQL-client module.	
					Regardless of whether this should be SQL-client module or <sql-client module<="" td=""><td></td></sql-client>	
					definition>, it is not clear exactly what it means for an SQL-transaction that is	
					the process of being initiated to "reference" that thing. Text elsewhere in	
					SQL/Foundation (for example, in Subclause 16.7, " <commit statement="">"),</commit>	
					strongly suggests that several distinct SQL-client modules can be "associated	
					with the same current SQL-transaction. Can they be associated with the SQL-	
					transaction without also being referenced by it? For that matter, can they be	
					"associated with " are sumerumous, then how can all the SQL alient modules	
					associated with are synonymous, then now can an the SQL-chem modules	
					2) The second personal entertains the notion of creation of a surger, and yet we	
					baye not been able to find any mention of this concept in any General Pule	
					Subclause 19.8 " <deallocate (cp3)="" does="" prepared="" require<="" statements"="" td=""><td></td></deallocate>	
					destruction of certain cursors and this is corroborated (redundantly?) by	
					Subclause 19.15 "callocate cursor statements" GR3)d) However, neither	
					Subclause 19.15, " https://www.com.com/www	
					statements" has any GR requiring a cursor to be created	
					4) As already noted, a cursor, having been created in somewhat mysterious	
					circumstances is never destroyed (unless it happens to be an allocated cursor, or	
					a cursor declared in a <compound 13.1<="" psm_subclause="" see="" sol="" statements="" td=""><td></td></compound>	
					" <compound statement="">" GR3)c)ii)2) and GR5)). It seems then that if n SOL</compound>	
					transactions in the same SOL-session "reference" the same SOL-client module.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
			, i i i i i i i i i i i i i i i i i i i		then each <declare cursor=""> contained in the corresponding <sql-client module<="" th=""><th></th></sql-client></declare>	
					definition> causes the creation of <i>n</i> distinct cursors. And yet Subclause 14.2,	
					" <open statement="">", SR1), says "Let CR be the cursor specified by DC", where</open>	
					<i>DC</i> is a <declare cursor="">. There are two problems with this:</declare>	
					— It is not clear which of those <i>n</i> cursors is the one specified by <i>DC</i> . Of course,	
					if the standard clearly specified that all but one of these had been destroyed by	
					this time, then there would be no ambiguity.	
					— The cited text in Subclause 4.32.1, "General description of cursors", makes it	
					clear that a cursor comes into existence at run-time and therefore, not being a	
					schema object, should not be referred to in a syntax rule. Since the SQL-session	
					context already includes cursor positions, perhaps it should also be defined to	
					include cursors. $(1 + 1)^{1/2} = 1$	
					5) The final sentence contains the word "also", which could be understood as	
					suggesting that some way of creating an extended dynamic cursor has already	
					destroyed as a consequence of its propered statement being deallocated before	
					SOL session termination	
					Solution	
					None provided with comment.	
	NLD-P02-002		2-Minor	P02-03.01,	FND-953 The following Language Opportunity has been noted:	
			Technical	Definitions	Language Opportunity:	
					There has been a discussion about Unicode 4.0 on the ISODBL list. [Ake has]	
					found out that Note 7 in SQL/Foundation will be affected, bescuase it contains	
					explicit code points. U+180E and U+205F have been added to the "Zs" class in	
					Unicode 4.0. Note that U+200B currently is of class "Zs", although it should not	
					be treated as white-space. The Unicode Technical Committee will probably	
					change the class for U+200B (ZERO-WIDTH SPACE) to "Cf" in the near	
					future.	
					Solution	
					None provided with comment.	
	NLD-P02-003		2-Minor	P02-04.10,	FND-845 The following Language Opportunity has been noted:	
			Technical	Collection types	Source: WG3:YYJ-016 (CAN-P02-001, USA-P02-005)	
					Language Opportunity:	
					The next edition of the SQL standard should standardize the syntax and	
					Semantics of one of more additional conection types.	
					Solution	
			1 Maion	P02 04 14 02	None provided with comment.	
	NLD-P02-004		T-Major Technics	Types of tables	FIND-944 The following Possible Problem has been noted: Source: $W(G_2)$ + W	
			recinical	- ypes of idoles	Possible Problem.	
					This Subclause with paragraph numbers added for expository nurnoses, says of	
					created temporary tables:	
					Possible Problem: This Subclause, with paragraph numbers added for expository purposes, says of created temporary tables:	

JJJCU
v

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
			berenzy		global one might have the same , when on reflection it clearly can't. This is not, be it noted, analogous to extended names, as explicitly intended by DBL:LON-156. Furthermore, "unique implementation-dependent name associated with the SQL- client module in which the created local temporary table is referenced" is unclear. Is this name persistent? One interpretation is that every occurrence of the must be contained in the same <sql-client module<br="">definition>. But perhaps it is intended to mean that there is no restriction on where the <table name> can occur; but only occurrences in externally-invoked procedures in the same module refer to the same thing (i.e. those of other modules refer to their own "local" temporary table). In paragraph 2, it is not clear whether a distinction is intended between "global temporary table contents" and "created local temporary tables", but presumably not. In paragraph 2, the meaning of "distinct within SQL-sessions" is unclear, because there is only one SQLsession active at any one time (even though there may be dormant ones).</table </sql-client>	
	NLD-P02-005		1-Major Technical	P02-04.14.02, Types of tables	 None provided with comment. FND-945 The following Possible Problem has been noted: Source: WG3:HBA-042 = H2-2003 Possible Problem: In this Subclause, as modified by [PSM-WD], Subclause 4.3.1, "Types of tables", with paragraph numbers added for expository purposes, says of declared temporary tables: 1) A declared local temporary table is a module local temporary table. A module local temporary table is a named table defined by a <temporary li="" table<=""> declaration> in an SQL-client module. A module local temporary table is effectively materialized the first time it is referenced in an SQL-session, and it persists for that SQL-session. 2) A declared local temporary table may be declared in an SQL-client module. 3) Inserted by SQL/PSM A declared local temporary table may be declared in an SQL-server module. 4) A declared local temporary table that is declared in an SQL-client module is a named table defined by a <temporary declaration="" table=""> that is effectively materialized the first time any <externallyinvoked procedure=""> in the <sql-client definition="" module=""> that contains the <temporary declaration="" table=""> is executed. A declared local temporary table is accessible only by <externally-invoked procedure="">s in the <sql-client definition="" module=""> that contains the <temporary declaration="" table=""> is executed. A declared local temporary table is accessible only by <externally-invoked procedure="">s in the <sql-client definition="" module=""> that contains the <temporary declaration="" table="">. The effective <schema name=""> of the <schema n<="" td=""><td></td></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></schema></temporary></sql-client></externally-invoked></temporary></sql-client></externally-invoked></temporary></sql-client></externallyinvoked></temporary></temporary>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
<i>"</i>					implementation-dependent SQL-session identifier associated with the SQL- session and a unique implementation-dependent name associated with the <sql-client definition="" module=""> that contains the <temporary declaration="" table="">. 5) Inserted by SQL/PSM A declared local temporary table that is declared in an SQL-server module is a named table defined by a <temporary declaration="" table=""> that is effectively materialized the first time any <module routine=""> in the <sql- server module definition> that contains the <temporary declaration="" table=""> is executed. A declared local temporary table is accessible only by <module routine>s in the <sql-server definition="" module=""> that contains the <temporary table declaration>. The effective <schema name=""> of the <schema qualified<br="">name> of the declared local temporary table may be thought of as the implementation-dependent SQL-server module definition> that contains the <temporary declaration="" table="">. The second sentence of paragraph 1 is no longer true when paragraphs 3 and 5 have been inserted by PSM. Moreover, whatever truth is expressed by paragraph 1 is repeated by paragraphs 2 and 4, which are specific to SQL-client modules. Evidently paragraphs 2 and 4 were inserted to correspond to paragraphs 3 and 5, so making paragraph 1 redundant, which should have been deleted at the same time. To say, in paragraphs 4 and 5, that a declared local temporary table has an <i>effective</i> <schema name=""> is misleading, since its name must be prefixed by MODULE. Solution None provided with comment.</schema></temporary></schema></schema></temporary </sql-server></module </temporary></sql- </module></temporary></temporary></sql-client>	
	NLD-P02-006		2-Minor Technical	P02-04.14.02, Types of tables	FND-969 The following Language Opportunity has been noted: Source: WG3:SIA-018 = H2-2004-429 Language Opportunity: Every view component is an underlying table. The reason that underlying table terminology was not used was that the hierarchy of underlying tables does not follow the hierarchy of syntactic containment, owing to the distinctive treatment accorded the tables and derived tables in the FROM clause of a <query specification> compared with other derived tables found in a <query specification>. If this issue can be overcome, it may be possible to eliminate the notion of view component and just use underlying tables. Solution None provided with comment.</query </query 	
	NLD-P02-007		1-Major Technical	P02-04.17, Integrity constraints	FND-703 The following Possible Problem has been noted: Source: WG3:BBN-139/X3H2-98-363 Possible Problem:	
					It seems that SQL3's specification of deferrable constraints is ill-specified. Referential constraints are based on the notion of marking rows for deletion	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					before the rows are effectively deleted at the end of the SQL statement. This is	
					necessary because updates cascaded by referential constraints need to be	
					"propagated" through rows marked for deletion in order to avoid anomalies (non	
					deterministic behavior). If a referential constraint is deferred, then rows that	
					need to be kept around for the execution of referential constraints will not be	
					present at the end of the transaction (or when the referential constraint is turned	
					to immediate). These rows will be deleted at the end of the SQL statements. So,	
					it is unclear how referential constraints are checked in these cases (e.g., are we	
					supposed to maintain multiple versions of the database and check the constraints	
					against those versions? If so, how do the updates are "propagated" to the current	
					version of the database?).	
					Another problem with deterrable constraints is that stored procedures and	
					triggers can never rely on the existence of a consistent database during their	
					execution because the application that caused the invocation of the stored	
					procedure and/or trigger could have deferred the checking of certain constraints	
					prior to the invocation of the procedure or trigger. (Please note that this has also	
					a major impact to the implementation of such concepts because plans generated	
					deferring such constraints)	
					Also it is not clear to me that deformable constraints and triggers work smoothly.	
					First REFORE triggers execute REFORE the SOL statement that activates	
					them. However, the REFORE execution cannot be guaranteed if referential	
					constraints are deferred because the execution of the BEFORE trigger needs to	
					be deferred as well. Second, if the BEFORE trigger is modifying the values of	
					transition variables such that they can be inserted/undated with correct values in	
					the database, what will happen with such values if the BEFORE trigger executes	
					after the database has been undated? Third, triggers are executed in a well	
					defined order. This is important to guarantee that changes to the database are	
					done in a deterministic manner. If constraints are deferred, then one may end up	
					deferring the execution of several instances of the same trigger for which there	
					is no well defined order of execution. This will lead to non-deterministic	
					behavior in the database.	
					Solution	
					None provided with comment.	
	NLD-P02-008		2-Minor	P02-04.27, SQL-	FND-725 The following Language Opportunity has been noted:	
	102 000		Technical	invoked routines	Source: WG3:FRA-122/X3H2-98-688)	
					Language Opportunity:	
					Subclause 4.27, "SQL-invoked routines", does not adequately describe the	
					concepts of dynamic binding and subject function selection.	
					Solution	
					None provided with comment.	
	NLD-P02-009		2-Minor	P02-04.32.	FND-607 The following Language Opportunity has been noted:	
	1111111111111111		2 minor	 ,	11.2 so, the following Language opportunity has been noted.	1

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
			Technical	Cursors	Source: DBL:LGW-146/X3H2-97-349 Language Opportunity: The ability to hold a cursor through rollback will be extremely useful to applications. Yet the second bullet of this Subclause says "a holdable-cursor is closed no matter what its state if the SQL-transaction is terminated with a rollback operation." This provision is not always necessary according to Jim Gray and Andrewas Reuter "Transaction Processing: Concepts and Techniques". Solution None provided with comment.	
	NLD-P02-010		2-Minor Technical	P02-04.32.01, General description of cursors	FND-929 The following Possible Problem has been noted: Source: WG3:HBA-040 Possible Problem: Although the second paragraph of this subclause defines terms to denote both varieties of dynamic cursors, it does not provide a way of referring to a cursor that is <i>not</i> dynamic. Solution None provided with comment.	
	NLD-P02-011		1-Major Technical	P02-04.33.04, SQL-statements and transaction states	 FND-923 The following Possible Problem has been noted: Source: WG3:HBA-029 Possible Problem: Subclause 4.33.4, "SQL-statements and transaction states", includes: If the initiation of an SQL-transaction occurs in an atomic execution context, and an SQL-transaction has already completed in this context, then an exception condition is raised: <i>invalid transaction termination</i>. At first sight it doesn't seem possible for transaction termination to be followed by transaction initiation "in" the same atomic execution context. In general, transaction initiation is caused by execution of an SQL-statement of the transaction-initiating kind and transaction termination is caused by executing an SQL-statement of a different kind (COMMIT or ROLLBACK). Note that Subclause 13.5, "<sql procedure="" statement="">", GR2), specifies that a new statement execution contexts", we are told (last paragraph) that an SQL-transaction contexts", we are told (last paragraph) that an SQL-transaction context. We conclude that the cited paragraph is relevant only when execution of a transaction-initiating statement</sql> (a) actually causes a transaction to be initiated, and (b) causes an exception to be raised of the special transaction, in the light of this observation, we perceive the following problems: 1) There is no General Rule in, for example, Subclause 13.5, "<sql procedure="" statement="">", to confirm the cited text.</sql> 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
					2) If execution of an SQL-statement causes an exception to be raised, then all changes to SQL-data and schemas are cancelled anyway. As the failing statement is also the one that initiated the transaction, the effect seems to be the same as that of a successful rollback, so what's the point in raising an additional exception expressing the fact that the transaction cannot be terminated? After all, the user executing the statement in question wasn't even trying to terminate the current transaction! We wonder if the rule was intended to cater for some eventuality other than the only one we can find. Even if the foregoing analysis proves to be refutable, it might be a good idea to add an explanation to Subclause 4.33.4, "SQL-statements and transaction states". Solution None provided with comment.	
	NLD-P02-012		1-Major Technical	P02-04.33.05, SQL-statement atomicity and statement execution contexts	FND-924 The following Possible Problem has been noted: Source: WG3:HBA-029 Possible Problem: Subclause 4.33.5, "SQL-statement atomicity and statement execution contexts", includes: The statement execution context brought into existence by the execution of an atomic SQL-statement or the evaluation of a <subquery> is an atomic execution context. The inclusion of "or the evaluation of a <subquery>", and the GRs of Subclause 7.15, "<subquery>", that back it up, seem questionable. Isn't expression evaluation always atomic? The question also arises as to whether deletion of the questionable text (and GRs) would make any material difference to the standard. A search of the SQL:2003 Foundation FDIS for the word "atomic" reveals no GRs that are conditional upon the atomicity or non-atomicity of a statement execution context. Instead, there are some special GRs for <subquery> that enforce its atomicity by creating and destroying a savepoint level, and in Subclause 13.5, "<sql procedure statement>", for undoing any changes to SQLdata or schemas made execution of by an atomic statement that terminates with an exception. It seems, then, that the only effects caused by atomicity are to do with savepoints and database updates. But it appears that database updates are not possible during evaluation of a subquery, being outlawed by Subclause 7.13, "<query expression>", SR23): 1) 23) A <query expression=""> QE shall not generally contain a <routine invocation> whose subject routine is an SQL-invoked routine that possible modifies SQL-data. Note that the BNF production for <subquery> is <left paren=""> <query< th=""><th></th></query<></left></subquery></routine </query></query </sql </subquery></subquery></subquery></subquery>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					expression> <right paren=""> and a <query expression=""> cannot contain an SQL</query></right>	
					procedure statement. SR23) in combination with the GRs of Subclause 10.4,	
					" <routine invocation="">", makes it impossible for an evaluation of an <routine< td=""><td></td></routine<></routine>	
					invocation> caused by evaluation of a <subquery> to cause an SQL-data change</subquery>	
					statement to be executed. Therefore it is impossible for evaluation of a	
					<subquery> to have any effect on SQL-data or schemas (possibly explaining the</subquery>	
					lack of a GR in Subclause 7.15, " <subquery>", specifying that changes to SQL-</subquery>	
					data and schemas are to be cancelled). Therefore any savepoints established	
					during evaluation of a <subquery> have to be ineffectual. Therefore there is no</subquery>	
					point in establishing a new, atomic, statement execution context for the	
					evaluation of a <subquery>.</subquery>	
					But that's not all! Consider the <query expression=""> SELECT foo() FROM 1,</query>	
					and suppose that there is some naw in the reasoning that leads to the conclusion	
					that the effect of the invocation of 100() cannot possibly depend on whether the	
					EDOM Travel direction context is atomic. In that case SELECT 100()	
					FROM I would not in general be equivalent to SELECI * FROM (SELECI	
					100() FROM 1) 1, for the shorter expression does not contain a <subquery>,</subquery>	
					whereas the fonger one does. In general, the consequences of the effect of	
					evaluation of a <query expression=""> possibly varying according to whether it is</query>	
					problems for optimisers	
					The foregoing analysis also brings into question the following sentence in	
					Subalausa 4.33.4. "SOL statements and transaction states":	
					1) If an <sol <subauerys="" a="" and<="" causes="" control="" evaluation="" of="" statements="" td="" the=""><td></td></sol>	
					there is no current SOI transaction, then an SOI transaction is initiated before	
					evaluation of the <subauery></subauery>	
					Perhaps this is pointless, too. If it proves not to be, the wisdom of starting a	
					transaction in the middle of executing an SOL-statement at an indeterminate	
					point in that execution, to boot is surely questionable.	
					In any case, we note that the sentence is not borne out by the GRs of Subclause	
					7.15, " <subquery>".</subquery>	
					Solution	
					Delete "or evaluation of a <subauery>" from the cited sentence of Subclause</subauery>	
					4.33.5, "SQL-statement atomicity and statement execution contexts": possibly	
					delete the cited sentence of Subclause 4.33.4. "SOL-statements and transaction	
					states"; delete GRs 1) ("Let OLDSEC "and 4) ("All savepoints") of	
					Subclause 7.15, " <subguery>". A search of the SOL:2003 Foundation FDIS for</subguery>	
					" <subquery>" reveals that a change might also be needed in Subclause 4.33.3.</subquery>	
					"SQL-statements and SQL-data access indication".	
					Other Parts of SQL:2003 have not been checked to see if they might be affected.	
					But see WG3:HBA-041.	
	NLD-P02-013		1-Major	<i>P02-04.35.02</i> ,	FND-972 The following Possible Problem has been noted:	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
			Technical	Savepoints	Source: WG3:SIA-031 = H2-2004-???	
					Possible Problem:	
					Neither here nor anywhere else is there any definition of the term savepoint.	
					Paragraph 5 of this Subclause contains:	
					If a <rollback statement=""> references a savepoint SS, then all changes made to</rollback>	
					SQL-data or schema subsequent to the establishment of the savepoint are	
					canceled, all savepoints established since SS was established are destroyed, and	
					the SQL-transaction is restored to its state as it was immediately following the	
					execution of the <savepoint statement="">.</savepoint>	
					The state of an SQL-transaction is not defined, nor is it referred to in Subclause 16.8, " <rollback statement="">".</rollback>	
					Presumably what is being referred to is some (or all) of the SQL-session context.	
					Which <savepoint statement=""> is not specified. Presumably the one that</savepoint>	
					established SS. So perhaps "	
					following the establishment of SS"; or even " as it was at that time".	
					Perhaps what is really meant is something to the effect of, A savepoint is a	
					preserved copy of (the values of specified elements) of the SQL-session context	
					at the time a <savepoint statement=""> was executed, plus sufficient data to enable</savepoint>	
					all subsequent changes to SQL-data or schemas in the current SQL-transaction	
					to be canceled.	
					When a <rollback statement=""> is executed, that contains a <savepoint specifier=""></savepoint></rollback>	
					SS, then all changes made to SQL-data or schema subsequent to the	
					establishment of SS are canceled, all savepoints established since SS was	
					established are destroyed, and elements of the SQL-session context are restored	
					to the values that were preserved in SS.	
					Paragraph 6 says:	
					It is implementation-defined whether or not, or now, a <rollback statement=""> that</rollback>	
					references a <savepoint specifier=""> affects diagnostics area contents, the contents</savepoint>	
					This implementation defined element is not mentioned in Anney P	
					"Implementation_defined elements" (or in Anney C. "Implementation_defined elements")	
					elements")	
					Solution:	
					Specify what happens in terms of the contents of the SOI -session context	
					Solution	
					None provided with comment.	
	NLD-P02-014		1-Major	P02-04.37, SQL-	FND-954 The following Possible Problem has been noted:	
			Technical	sessions	Source: WG3:ZSH-037R1/H2-2003-???	
					Possible Problem:	
					WG3:FRA-045r4 proposed no changes to what is now WG3:ZSH-013,	
					Subclause 4.37, "SQLsession".	
					However, according to WG3:FRA-045r4, Section 2.1, "Authorization stack":	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					There is a stack of SQL-session contexts. There is one cell on this stack when the SQLsession begins. An additional SQL-session context is pushed on the stack for each <routine invocation>, and is removed when the <routine invocation=""> completes execution. There is no reference to this anywhere in this subclause, although there are various statements of the form "An SQL-session has a". Moreover, the list of SQL-session contents is incorrect and incomplete. The term "current SQL-session identifier" is listed, where the meaning of "current" is indicated in the following NOTE (55 in WG3:ZSH-013) and evidently used to distinguish the "current" SQL-session from dormant SQL-sessions. It is therefore probably intended to refer to the SQL-session identifier of the currently active (as opposed to dormant) SQL-session. If this surmise is correct, then the "current SQL-session user identifier" is missing. There is no reference to the authorization stack, though the two terms used to refer to the components of the only visible cell of that stack are mentioned. Solution None provided with comment.</routine></routine 	
	NLD-P02-015		1-Major Technical	P02-04.37.04, Execution contexts	 FND-955 The following Possible Problem has been noted: Source: WG3:ZSH-037R1/H2-2003-??? Possible Problem: This subclause contains the statement: There is always a statement execution context, a routine execution context, and zero or more trigger execution contexts. There is a significant and unnecessary inconsistency between the descriptions of routine execution contexts and trigger execution contexts. Consider what happens if an SQL-invoked routine R1 invokes another, R2. Are there now one or more than one routine execution contexts? The answer is clearly there is one in each of two levels of the stack of SQL-session contexts, as is made clear by Subclause 10.4, "<routine invocation="">". Whether there is a routine execution context when no routine has been invoked is debatable: it could be (and indeed is) said that there is an empty one; or it could be said that there is none. In which case, it would be true to say that "there are zero or more routine execution contexts.</routine> Consider now how it arises that there is more than one trigger execution context. The only case that springs to mind is that of the triggered action of a trigger 11, causing another trigger T2 to fire. In this case, each trigger will have a trigger execution context. However, it seems fairly clear that the triggered action of T2 cannot access the state changes in the trigger execution context of T1. Therefore, to say that there are, during the execution of T2, two trigger execution contexts, although true in a sense, is unhelpful. 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					the same SQLsession context; unless, of course, T1 invokes a routine that causes	
					T2 to fire, in which case a new SQL-session context is created, containing a new	
					routine execution context. However, whether or not it contains, when created,	
					the trigger execution context of T1, we are unable to discover.	
					Solution	
					None provided with comment.	
	NLD-P02-016		1-Major	<i>P02-05.04</i> ,	FND-932 The following Possible Problem has been noted:	
			Technical	Names and	Source: WG3:HBA-050R1	
				identifiers	Possible Problem:	
					SR19) of this subclause is:	
					19) An <identifier> that is a <correlation name=""> is associated with a table within</correlation></identifier>	
					a particular scope.	
					The scope of a <correlation name=""> is either a <select row="" single="" statement:="">,</select></correlation>	
					<subquery>, or <query specification=""> (see Subclause 7.6, ""),</query></subquery>	
					or is a <trigger definition=""> (see Subclause 11.39, "<trigger definition="">"). Scopes</trigger></trigger>	
					may be nested. In different scopes, the same <correlation name=""> may be</correlation>	
					associated with different tables or with the same table.	
					The inclusion of <subquery> is puzzling. For consider that if such a scope is</subquery>	
					contained in a <subquery>, then it must also be wholly contained in some</subquery>	
					<query specification=""> contained in that <subquery>. Furthermore, a <subquery></subquery></subquery></query>	
					that contains more than one <query specification=""> cannot possibly constitute the</query>	
					scope of any <correlation name="">. For example:</correlation>	
					(SELECT * FROM 11 UNION SELECT * FROM 12)	
					confined to the particular $\langle query specification \rangle$ in which it is defined. A scalar	
					expression could be added to the <subauery> that includes an outer reference</subauery>	
					but the <u>correlation</u> names used in that reference would have a wider scope than	
					the <subauery></subauery>	
					The inclusion of $<$ query specification> is also suspect, because the scope of a	
					<pre><correlation name=""> is not necessarily a whole <ouery specification=""></ouery></correlation></pre>	
					Solution	
					None provided with comment	
	NLD-P02-017		1-Major	P02-05.04	FND-946 The following Possible Problem has been noted:	
	102 017		Technical	Names and	Source: WG3:HBA-042 = H_2 -2003-	
			reenneur	identifiers	Possible Problem:	
					This subclause says (regarding local temporary tables):	
					Something needs to be said, in either or both of the Syntax Rules and General	
					Rules of Subclause 5.4, "Names and identifiers", about how a	
					identifies a created temporary table.	
					As a minimum, a reference to a created local temporary table must be prohibited	
					in any <schema routine=""> R, because, by the time R is invoked, it cannot be</schema>	
					regarded as being syntactically contained in the <sql client="" definition="" module=""></sql>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					that contained the <externally-invoked procedure=""> which created it. Solution As a minimum, a Syntax Rule should be added, to the effect that: 1) If identifies a created local temporary table, then shall not be contained in a <schema routine="">. Note: This appears to remove the need for any reference to created local temporary tables in Subclause 10.4, "<routine invocation="">", General Rule 5) d) i), which is a problem for PSM.</routine></schema></externally-invoked>	
	NLD-P02-018		2-Minor Technical	P02-06.01, <data type></data 	FND-729 The following Language Opportunity has been noted: Source: WG3:YGJ-112 (SQL/MM YGJ-023), Paul Cotton for WG4, July 6, 1999, and Paul Scarponcini via email on 6 July 1999 Language Opportunity: According to YGJ-112: "REF types need to be scoped; i.e., the table(s) they refer to must be explicitly provided. If a column is of type REF type, the scope may be defined at table creation time. If the column is of type UDT which contains REF type attributes, then the scope must be declared when the UDT is created. The SQL/MM Part 3: Spatial standard defines the UDTs for spatial data. The standard is unable to predict in which tables the referenced information will be stored; this is a function of database design. Therefore, column scoping must be expanded to support deeply nested references, i.e., REF types within a UDT or ARRAY. This would allow a user, when creating tables, to define the scope of a UDTs REF type as part of the column definition for a column of type UDT." When a <reference type=""> is used as the data type of an attribute of a structured type, the <scope clause=""> must be specified when the encompassing user-defined type is defined. It is a Language Opportunity to be able to specify the <scope clause> of the "nested" <reference type="">s when a column is defined on the encompassing user-defined type. Paul Scarponcini added: This applies to ARRAYs as well (e.g., an ARRY of REF, and ARRAY of UDTs having REF attributes. The resultant syntax may be quite messy, as different REFs within the column may have different scopes. Would it be worth considering reversing the scope specification: when the reference dtable is created, specify that it shall be included in the scope for a particular column, rahter than specifying the referenced table when the referencing column is specified? Solution</reference></scope </scope></reference>	
	NLD-P02-019		2-Minor Technical	P02-06.01, <data type></data 	FND-730 The following Language Opportunity has been noted: Source: WG3:YGJ-112 (SQL/MM YGJ-023) and Paul Cotton for WG4, July 6, 1999	
SEQ	Cmnt	See				Addressed
-----	-------------	------	----------------------	---	---	-----------
#	ID	Also	Severity	Reference	Description	By
					Language Opportunity: According to YGJ-112: "A second limitation of SQL 99 with respect to REF types is that they only achieve uni-directional "pointers"." A REF type value may be de-referenced to obtain the instance to which it refers. It is a Language Opportunity to provide direct support for determining all instances of a REF type which refer to a particular instance. Solution	
	NLD-P02-020		2-Minor Technical	P02-06.01, <data type></data 	FND-812 The following Language Opportunity has been noted: Source: WG3:PER-098R1/H2-2001-059 Language Opportunity: Perhaps Feature S096, "Optional array bounds", can be folded in Feature S091, "Basic array support". Solution None provided with comment.	
	NLD-P02-021		2-Minor Technical	P02-06.04, <value specification> and <target specification></target </value 	FND-692 The following Language Opportunity has been noted: Source: DBL:CWB-081/X3H2-98-068 Language Opportunity: Although there is provision for refining a <value expression=""> of row type or structured type, there is no provision for refining a <target specification="">. As a result, a field of a row or an attribute of a structured type cannot be passed as output or in/out argument of an SQL-invoked routine, or used in other target contexts. This problem is partially remedied in PSM <assignment statement="">. Possibly the support for refined targets can be adapted from PSM and moved to Foundation. Solution None provided with comment.</assignment></target></value>	
	NLD-P02-022		2-Minor Technical	P02-06.04, <value specification> and <target specification></target </value 	FND-723 The following Language Opportunity has been noted: Source: WG3:FRA-132/X3H2-98-694 Language Opportunity: Currently we have no capability to treat an <element reference=""> as a <target specification>. This precludes their use as output arguments of routine invocations, for example. The same observation can be made of <field reference>, <dereference operation="">, <reference resolution="">, and <method invocation> (some of these subject to the restriction that the method must be a mutator). (Lest you object that [Fred is] thinking of allowing surreptitious updates to column values by referencing them as output arguments of a routine invocation, be it noted that these expressions can also be used with parameters and variables.) However, [Fred believes] that the general solution to this problem is to introduce a notion of l-values and r-values, as in the specification of C.</method </reference></dereference></field </target </element>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
					Solution	
					None provided with comment	
	NI D-P02-023		1-Major	P02-06.06	WC3-P02-001	
	1120 1 02 025		Technical	<identifer< td=""><td>Consider the expression SELECT $*$ FROM T T1 WHERE C1 = (SELECT</td><td></td></identifer<>	Consider the expression SELECT $*$ FROM T T1 WHERE C1 = (SELECT	
			reenneur	chain>	MAX (C1) FROM T T2 WHERE T1 C2 $>$ C1) It is surely indisputable that	
					the two references to C1 in the subguery are syntactically legal and are	
					references to T2.C1, according to the normal block-scoping rules that are	
					commonly used in SQL implementations. And yet SR8) appears to make them	
					illegal. SR8)a)ii) is applicable:	
					ii) [the <identifier chain="">] shall be contained in the scope of one or more</identifier>	
					range variables whose associated tables include a column whose <column< td=""><td></td></column<>	
					name> is equivalent to I1 or in the scope of a <routine name=""> whose associated</routine>	
					<sql declaration="" list="" parameter=""> includes an SQL parameter whose <sql< td=""><td></td></sql<></sql>	
					parameter name> is equivalent to <i>I</i> ₁ . Let the phrase <i>possible scope tags</i> denote	
					those range variables and <routine name="">s.</routine>	
					In the example, C1 is contained in the scope of both T1 and T2. The	
					continuation of this subrule is a Case whose first subrule is:	
					1) If the number of possible scope tags in the innermost scope containing a	
					possible scope tag is 1 (one), then let <i>IPST</i> be that possible scope tag.	
					Now, if this condition were true in our example, and the single possible scope	
					tag were 12, then all would be well, but unfortunately that does not appear to be	
					the case. The innermost scope containing a possible scope tag for CT consists of two fragments. SELECT MAX (C1) and WHERE T1 C2 > C1 (construction)	
					two fragments: SELECT MAX (C1) and wHERE $11.22 > C1$ (see Subclause 7.6 "stable references" SB5). How many of the two possible score tags for C1	
					7.0, <table c1<="" for="" how="" many="" of="" possible="" reference,="" scope="" sks).="" tags="" td="" the="" two=""><td></td></table>	
					one but the one in question is T1 (contained in the $\langle where clause \rangle$) not T2	
					If on the other hand "in" means "that are in scope in" then the answer is two	
					for both T1 and T2 are in scope. Of course "in" is not intended to mean either	
					of those things: in fact, it is clear under this close examination that "in the	
					innermost scope" is not an appropriate phrase here at all.	
					Having shown that "in the innermost scope" is not appropriate, we now show	
					that "containing a possible scope tag" isn't appropriate either. Consider the	
					following slightly simpler example: SELECT * FROM T WHERE C1 = (
					SELECT MAX (C1) FROM T). How many possible scope tags do we have	
					now? Well, MAX (C1) is in the scope of the T that is defined in the outer	
					<from clause=""> and it is also in the scope of the other T that is defined in the</from>	
					<subquery>'s <from clause="">. Do we have two possible scope tags that are both</from></subquery>	
					named T, or do we have just one possible scope tag with two distinct reasons for	
					it being a possible scope tag? In any case, whether we have one or two, how	
					many are "in the innermost scope containing a possible scope tag"?	
					The scope of the T defined in the <subquery> is just SELECT MAX (C1),</subquery>	
					which contains no possible scope tags at all. The scope of the T defined in the	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
#	ID	Also	Severity	Reference	Description outer <from clause=""> consists of the fragments SELECT * and WHERE C1 = (SELECT MAX (C1) FROM T), which happens to contain T, though not the T that has this scope! It seems that when the same range variable name is used for two or more different purposes (and necessarily in that case with different scopes in each case), and when a column reference lies within each of those scopes, only the one <i>having</i> the innermost of those scopes is applicable (and so that one is applied). And when two or more <i>different</i> range variables are used, as in our first example, then they are all applicable but it is again the one <i>having</i> the innermost scope that is applied, provided, of course, that there is exactly one range variable qualifying as a possible scope tag, whose scope is the innermost of the scopes containing the column reference. Solution</from>	Ву
					None provided with comment.	
	NLD-P02-024		2-Minor Technical	P02-06.09, <set function specification></set 	FND-819 The following Language Opportunity has been noted: Source: WG3:PER-044R1/H2-2000-619 Language Opportunity: The proponents of multiargument GROUPING function believe that it is a trivial extension of the single argument function, and therefore does not warrant a separate feature. This could be achieved by simply deleting the Conformance Rule that creates Feature T433, "Multiargument GROUPING function", thereby allowing all GROUPING functions to fall under Feature T431, "Extended grouping capabilities". Solution None provided with comment.	
	NLD-P02-025		2-Minor Technical	P02-06.12, <cast specification></cast 	WG3-P02-002 SR10) prohibits the containment of a <collate clause=""> in the target <data type="">, <i>TD</i>. When a <data type=""> is specified, this is clear, but it is not so clear when a <domain name=""> is specified, in which case SR1) defines <i>TD</i> to be "the <data type> of the domain". The BNF for <domain definition=""> (Subclause 11.24) doesn't include a <data type="">, though it does include a <predefined type="">. A domain descriptor is said to include a data type descriptor, but note carefully that <i>every</i> character data type descriptor contains the fully qualified name of a collation. Note that a <predefined type=""> might include a <collate clause="">, and also that if the <domain definition=""> contains a <collate clause="">, then that is considered to be equivalent to the containment of that <collate clause=""> in the <predefined type="">. Some tidying appears to be needed, but what the intended rule is in the case of casting to a domain needs to be determined before a precise redrafting can be proposed. Note that a related problem exists in the SQL:2007 WD for SQL/XML, Subclause 6.4, <xml cast="" specification="">, SR10), so this P.P. might eventually need to be cloned as a CD ballot comment against Part 14.</xml></predefined></collate></collate></domain></collate></predefined></predefined></data></domain></data </domain></data></data></collate>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					Solution	
					None provided with comment.	
	NLD-P02-026		2-Minor	P02-06.12, <set< td=""><td>FND-693 The following Language Opportunity has been noted:</td><td></td></set<>	FND-693 The following Language Opportunity has been noted:	
			Technical	function	Source: FCD1/1998 NLD-P02-017, DBL:CWB-132/X3H2-98-187	
				specification>	Language Opportunity:	
					We do not understand SR 4). If an outer reference is permitted at all, surely it	
					should be permitted any number of times, just as literals and host variable names	
					can occur any number of times. We would add that we see no reason to prohibit	
					outer references altogether. For example, if SUM(OUTER.C1) is legal, surely	
					SUM(OUTER.C1+OUTER.C1) is also legal. Besides, why should column	
					references that are not outer references be prohibited as soon as there is an outer	
					reference? SR 4) of Subclause 6.9, " <set function="" specification="">", says:</set>	
					4) The <value expression=""> simply contained in <set function="" specification=""></set></value>	
					shall not contain a \langle set function specification \rangle or a \langle subquery \rangle . If the \langle value	
					expression> contains a column reference that is an outer reference, then that	
					outer reference shall be the only column reference contained in the <value< td=""><td></td></value<>	
					expression.	
					we agree that the above fulle is overly restrictive. However, we believe this fulle was adopted in SOL 02 to prohibit query formulations of the form:	
					SELECT *	
					FROM t1	
					GROUP BY	
					HAVING (SELECT c21	
					FROM t2 CROUD BY	
					WHERE (SELECT c3	
					FROM t3	
					WHERE SUM (t1.c12 + t2.c22) >	
) In the above example, outer references from multiple levels are being referenced.	
					in the same aggregate function. Semantically, this does not make sense and must	
					he prohibited	
					Solution	
					None provided with comment	
	NLD-P02-027		2-Minor	P02-06.15,	FND-816 The following Language Opportunity has been noted:	
			Technical	<subtype< td=""><td>Source: WG3:PER-099/H2-2001-061</td><td></td></subtype<>	Source: WG3:PER-099/H2-2001-061	
				treatment>	Language Opportunity:	
					Perhaps Feature S162, "Subtype treatment for references", can be folded into	
					Feature S161, "Subtype treatment".	
					Solution	
					None provided with comment.	
	NLD-P02-028		2-Minor	P02-06.15,	FND-829 The following Language Opportunity has been noted:	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
			Technical	<subtype treatment></subtype 	Source: WG3:PER-186/H2-2001-??? Language Opportunity: WG3:PER-099 extended <subtype treatment=""> so that an expression of type REF($t1$) would be TREATed as one of type REF($t2$) if $t2$ is a subtype of $T1$. It was noted that, in that case, it should also be possible to TREAT: — An expression of type $t1$ ARRAY[n] as one of type $t2$ ARRAY[n]. — An expression of type $t1$ MULTISET as one of type $t2$ MULTISET. — An expression of type ROW(, $f1 t1$,) as one of type ROW(, $f1 t2$,). In the ROW case, it might even be possible to support TREATment over more than one field. For example, an expression of the type ROW(, $f1 t1$,, $f2 t1$,) might be TREATable as ROW(, $f1 t1$,, $f2 t2$,), as ROW(, $f1 t2$,, $f2$ t1,), or as ROW(, $f1 t2$,, $f2 t2$,), even though SQL does not (at the time of writing this Language Opportunity) support multiple inheritance in general. In the ROW case, it would also be necessary to decide whether field names must match as indicated in these examples. Solution</subtype>	
	NLD-P02-029		2-Minor Technical	P02-06.28, <string value<br="">expression></string>	None provided with comment. FND-858 The following Language Opportunity has been noted: Source: WG3:ICN-054R2 = H2-2002 Language Opportunity: The term "character string operands" was used to replace a previously undefined term "components" in SR2. Is this the correct terminology to use? Solution None provided with comment.	
	NLD-P02-030		2-Minor Technical	P02-06.34, <boolean value<br="">expression></boolean>	FND-920 The following Language Opportunity has been noted: Source: WG3:ZSH-129 = H2-2002 Language Opportunity: The rules for known-not-null conditions in SR3) are more complicated than most implementations are prepared to implement, and not necessary for most users. The full implementation of known not null should be placed in a conformance feature. Without the feature, a much simpler definition should apply. Solution None provided with comment.	
	NLD-P02-031		2-Minor Technical	P02-06.35, <array value<br="">expression></array>	 FND-808 The following Language Opportunity has been noted: Source: (was Possible Problem FND736) WG3:PER-171/H2-2001-??? (FCD1/2000 NLD-P02-027), from WG3:YGJ-074/X3H2-99-164R1 Language Opportunity: The ability to extract a subarray of an array would be useful. Such an ability would also satisfy a separate Language Opportunity to be able to truncate an array. 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
			j		Solution	— j
					None provided with comment.	
	NLD-P02-032		2-Minor	P02-07.04,	FND-756 The following Language Opportunity has been noted:	
			Technical	<table< td=""><td>Source: WG3:YGJ-069r1 = H2-99-155r3 and WG3:BHX-096/H2-2000-248R1</td><td></td></table<>	Source: WG3:YGJ-069r1 = H2-99-155r3 and WG3:BHX-096/H2-2000-248R1	
				expression>	Language Opportunity:	
					It might be useful to be able to filter windowed results based on the values of	
					<olap function="">, most likely through a new clause analogous to <where< td=""><td></td></where<></olap>	
					clause> and <having clause="">, but following <window clause=""></window></having>	
					Solution	
					None provided with comment.	
	NLD-P02-033		2-Minor	P02-07.09,	FND-610 The following Language Opportunity has been noted:	
			Technical	<group by<="" td=""><td>Source: DBL:LGW-146/X3H2-97-349</td><td></td></group>	Source: DBL:LGW-146/X3H2-97-349	
				clause>	Language Opportunity:	
					Continuing work is needed to complete object support as outlined in "Providing	
					Rich Query Functionality"	
					(DBL:LHR-0/8 = X3H2-95-462) with regard to expanding GROUP BY to	
					the query. The ability to group the result of a table expression by the value of	
					avpressions is important to many applications. The ability to name these	
					grouping expressions and use those names to retrieve the results of the grouping	
					column cum expression in the select list of the table expression is equally	
					important to avoid applications having to repeat the expression (giving	
					opportunity for errors) in the select list.	
					Solution	
					None provided with comment.	
	NLD-P02-034		2-Minor	P02-07.12,	FND-528 The following Language Opportunity has been noted:	
			Technical	<query< td=""><td>Source: DBL:MAD-170/X3H2-96-544R1, point 2.1, FCD1/1998 CAN-P02-</td><td></td></query<>	Source: DBL:MAD-170/X3H2-96-544R1, point 2.1, FCD1/1998 CAN-P02-	
				specification>	031, DBL:CWB-132/X3H2-98-187	
					Language Opportunity:	
					DBL:MAD-170/X3H2-96-544R1, point 2.1, noted:	
					The definition of a possibly nullable result column in the Syntax Rules of	
					Subclause 7.12, " <query specification="">", is broader than necessary, in that an</query>	
					aggregate of a column that is known not nullable is regarded as possibly	
					nullable. For example, SUM(EMP.EMPNO) is defined as possibly nullable,	
					even II EMP.EMPNU IS declared NUT NULL.	
					DDL. $WD-152/A5\Pi 2-90-107$ added.	
					specification for example SUM(EMPNO) is known not nullable when	
					FMPNO is known not nullable. However, GR 3(b) of Subclause 6.9 "-set	
					function specification>", makes it clear that (with the exception of COUNT)	
					<pre><set function="" specification="">s return NULL when they are applied to an empty</set></pre>	
					table. Hence, we assume that <set function="" specification="">s are possibly nullable,</set>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
SEQ #	Cmnt ID	See Also	Severity 2-Minor Technical	Reference P02-07.12, <query specification></query 	Description except for COUNT. And, that is what SR 12) of Subclause 7.12, " <query specification="">", specifies. Hence, we believe that there is no problem with SR 12) of Subclause 7.12, "<query specification="">". Solution None provided with comment. FND-908 The following Language Opportunity has been noted: Source: P02, SQL/Foundation, Subclause 7.12, "<query specification="">", CR 4) and CR 8) Language Opportunity: Conformance Rule 4) as formulated does not impose a restriction on the user writing SQL and as such does not follow the required model for Conformance Rules. Fred Zemke in an email to Stephen Cannan dated 2002-10-17 wrote: Subclause 7.12, "<query specification="">" CR 4) - this is an example of the occasional practice of using the CRs to alter the definition of a defined term. This practice seems borderline to me. On the one hand, the CRs are regarded as merged with the SRs whenever the designated feature is absent, and definitions appear in the SRs, so it would seem possible to make a redefinition in a CR. On the other hand, does a redefinition constitute a limitation on the user? I think the better approach is the one taken regarding</query></query></query></query>	Addressed By See Comment
					CR 4) - this is an example of the occasional practice of using the CRs to alter the definition of a defined term. This practice seems borderline to me. On the one hand, the CRs are regarded as merged with the SRs whenever the designated feature is absent, and definitions appear in the SRs, so it would seem possible to make a redefinition in a CR. On the other hand, does a redefinition constitute a limitation on the user? I think the better approach is the one taken regarding functional dependencies, for example, Subclause 7.12, " <query specification="">", CR 3). This could have been done by defining a term such as 'group-invariant <value expression="">' in the SRs, saying that all derived columns in the SELECT list of a grouped query must be group-invariant <value expression="">. Instead the approach taken is essentially to carve out two categories of derived column: the kind permitted in the SELECT list of a grouped query by the SRs, and the more restricted kind permitted by the CRs. Returning to 7.12 CR 4), the path would be to define two notions of updatable, using the more liberal one in the SRs and restricting to the more conservative one in the CRs. In fact, we already have two terms, updatable and simply updatable. The conclusion I am coming to is that this CR should be a CR that restricts to 'simply updatable' unless Feature</value></value></query>	
	NLD 202 026		2 Minor	P02-07 15	T111 is present. Conformance Rule 8) as formulated does not impose a restriction on the user writing SQL and as such does not follow the required model for Conformance Rules. This rule should be deleted, and, in any subclause that uses 'updatable' in an SR, there should be a CR that restricts the use of UNION unless Feature T111 is present. Solution None provided with comment, but the body of the comment outlines a solution	
	INLD-P02-030		Z-Minor Technical	<subquery></subquery>	Source: WG3:HBA-050R1	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
	NLD-P02-037		2-Minor Technical	P02-08.02, <comparison predicate></comparison 	LecomptionLecomptionLecomptionLecomptionLanguage Opportunity:HBA-050 shows that the need for the BNF term <subquery> is nothing like as strong as it once might have been, and has given rise to a certain amount of difficulty and confusion. Perhaps it would be better to dispose of the term altogether (though <scalar subquery="">, <row subquery="">, and almost certainly need to be retained) and treat parenthesized <query </query expression>s in similar style to our treatment of parenthesized <value </value expression>s.Any proposal to address this Language Opportunity should of course check for existing uses of <subquery> in Parts other than Foundation.SolutionNone provided with comment.FND-909 The following Language Opportunity has been noted:Source: WG3:ZSH-155 = H2-2002Language Opportunity:The Syntax Rules convert all comparison predicates so that they only use < and =. The GRs for comparison of user-defined types spell out rules for > and other comparisons even though they have been transformed away. NOTE 167 following the GR claims that these unreachable GRs are there for informational purposes. In the case of RELATIVE order, there are some strong assumptions being made that $RF(X,Y) = -RF(Y,X)$; otherwise, the system breaks down. We should document what are the expectations for the relative order function somewhere. We do not find such documentation either in <user-defined ordering<br=""></user-defined>function> or in Concepts.</br></br></subquery></row></scalar></subquery>	
	NLD-P02-038		2-Minor Technical	P02-09.03, Data types of results of aggregations	FND-836 The following Language Opportunity has been noted: Source: WG3:YYJ-030R2 = H2-2001 and WG3:ZSH-155 = H2-2002 Language Opportunity: This subclause uses terms that are less precise than they should be. Specifically, the term result data type and data type of the result, without specifying the result of what. The first sentence of Function says: "Specify the result data type of the result of an aggregation". Moreover the term aggregation does not suggest the sense in which it is used here, having since been used extensively in the context of OLAP, see subclause 04.17.03 "Aggregate functions". A better title would be Data types of results of n-adic operations. Were this title adopted, the first sentence could be rewritten as, for example, Let IDTS be a set of data types specified in an application of this Subclause, and let O be the operation. Solution None provided with comment.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
	NLD-P02-039		2-Minor Technical	P02-09.05, Type precedence list determination	FND-709 The following Language Opportunity has been noted: Source: WG3:YGJ-021 Language Opportunity: Paper DBL:BBN-168 added a Syntax Rule to Subclause 11.50, " <sql-invoked routine>", to prohibit the use of ROW because there is nothing in Subclause 9.5, "Type precedence list determination", to handle the type precedence requirements of anonymous row types. Solution</sql-invoked 	_,
					None provided with comment.	
	NLD-P02-040		1-Major Technical	P02-10.04, <routine invocation></routine 	FND-857 The following Possible Problem has been noted: Source: DCOR/2002, USA-STC-031 Possible Problem: There is no definition of how to pass booleans or LOBs to external programs. More generally, the question of how to convert any SQL type to a host language type at the interface to an SQL-invoked routine has never been addressed. Probably it was assumed that the same mechanism as was already defined for module language and embedded language applied, but in fact there are no rules to back up this assumption. If this assumption is correct, then the rules in Subclause 13.4, "Calls to an <externally-invoked procedure="">", are probably appropriate. Perhaps they should be placed in a separate subclause so they can be referenced by both <routine invocation> and also <externally invoked="" procedure="">. See also paper WG3:PER-176. Solution</externally></routine </externally-invoked>	
	NLD-P02-041		1-Major Technical	P02-10.04, <routine invocation></routine 	 FND-956 The following Possible Problem has been noted: Source: WG3:ZSH-037R1/H2-2003-??? Possible Problem: 5) Preserve the current SQL-session context CSC and create a new SQL-session context RSC derived from CSC as follows: This appears to specify what happens to every element of an SQL-session context when a new SQL-session context is created. However, it does not say what happens to: The zero or more trigger execution contexts The values of all valid locators The text defining the SQL-path (which in any case seems somewhat redundant, since the SQL-path is taken care of) The text defining the default transform group name The text defining the user-defined type name-transform group name pair for each userdefined type explicitly set by the user 	

SEQ	Cmnt	See	a t	D.A		Addressed
#	ID	Also	Severity	Reference	Description	Ву
					5) Preserve the current SQL-session context CSC and create a new SQL-session	
					context RSC as follows:	
					···	
					Solution	
	NI D D02 042		2.16	DO2 11 Sahama	None provided with comment.	
	NLD-P02-042		2-Minor	P02-11, Schema definition and	FND-710 The following Language Opportunity has been noted:	
			Technical	manipulation	Source: wG5.1GJ-021	
				1	A RENAME TABLE statement has been strongly desired for a very long time	
					and any users will be expecting to see it in SOL3	
					Solution	
					None provided with comment	
	NLD-P02-043		2-Minor	P02-11. Schema	FND-694 The following Language Opportunity has been noted:	
	1120 102 013		Technical	dewfinition and	Source: DBL:CWB-114/X3H2-98-169	
				manipulation	Language Opportunity:	
					The current choices for <drop behavior="">, RESTRICT and CASCADE, are too</drop>	
					limiting. CASCADE is so sweeping that the user must hesitate to use it, not	
					knowing what may be dropped. RESTRICT, on the other hand, is so limited that	
					the user must find all dependencies and drop them in the proper order. There is a	
					third model, based on the notion of invalidation. With this model, a dependent	
					definition does not block a drop; instead, the dependent object is simply marked	
					invalid. Later usage of an invalid object causes its recompilation, which may	
					Solution	
					Solution None provided with comment	
	NI D D02 044		2 Minor	P02-11.03	None provided with comment.	
	NLD-102-044	Z-IVIINOr Technical	Z-Ivinio Technical	<table< td=""><td>Source: WG3:PFR_104/H2_2001_085R1</td><td></td></table<>	Source: WG3:PFR_104/H2_2001_085R1	
			reenneur	definition>	Language Opportunity:	
					The ability to specify options for inheriting column default and identity column	
					properties, as in the ke clause>, would also be beneficial for the <as subquery<="" td=""><td></td></as>	
					clause>.	
					Solution	
					None provided with comment.	
	NLD-P02-045		2-Minor	P02-11.03,	FND-874 The following Language Opportunity has been noted:	
			Technical	<table< td=""><td>Source: WG3:DRS-095</td><td></td></table<>	Source: WG3:DRS-095	
				aejinition>	Language Opportunity:	
					Since in section 1.1.2 [of WG3:DRS-095] we gave reasons for determining the	
					<reterence generation=""> implicitly, it would be most convenient if the <column< td=""><td></td></column<></reterence>	
					constraint definition>s necessary for derived reference representations were	
					descriptor	
	NLD-P02-045		2-Minor Technical	P02-11.03, <table definition></table 	Solution None provided with comment. FND-874 The following Language Opportunity has been noted: Source: WG3:DRS-095 Language Opportunity: Since in section 1.1.2 [of WG3:DRS-095] we gave reasons for determining the <reference generation=""> implicitly, it would be most convenient if the <column constraint="" definition="">s necessary for derived reference representations were implicit, and determined by examination of the corresponding user-defined type descriptor.</column></reference>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					Solution	
					None provided with comment.	
	NLD-P02-046		2-Minor	P02-11.05,	FND-642 The following Language Opportunity has been noted:	
			Technical	<default clause=""></default>	Source: DBL:LGW-152/X3H2-97-352 (also DBL:LGW-023/X3H2-97-044,	
					SEQ# 222, CAN-F-062, converted to LO by WG3:BHX-038/H2-2000-018R3)	
					Language Opportunity:	
					It might be useful to allow default values for row types, perhaps by using row	
					constructors.	
					Solution	
-					None provided with comment.	
	NLD-P02-047		2-Minor	<i>P02-11.05</i> ,	FND-712 The following Language Opportunity has been noted:	
			Technical	<default clause=""></default>	Source: WG3:YGJ-021 and WG3:PER-098R1/H2-2001-059	
					Language Opportunity:	
					It is not possible to specify default values for columns or attributes of an array	
					type, a multiset type, a reference type, a row type, or a user-defined type.	
					Solution	
				D00 11 00	None provided with comment.	
	NLD-P02-048		2-Minor	P02-11.08,	FND-349 The following Language Opportunity has been noted:	
			Technical	constraint	Source: WG3:YGJ-0/4/X3H2-99-164R1 (Bill Kelley noted the following	
				definition>	Language Opportunity, which has been modified by Fred Zemke)	
				,	Language Opportunity:	
					collections types, referential integrity is not definable for elements of	
					Evample: Assume table EMPLOYEE has PRIMARY KEV EMP. ID of type	
					INTEGER:	
					CREATE TABLE MANAGER (
					EMPNO INTEGER,	
					MANAGES INTEGER ARRAY [20])	
					Here "MANAGES" refers to a set of employees, but there is no way to say that	
					they should reference employees. That is, if one were to write:	
					EMDNO INTEGER	
					MANAGES INTEGER ARRAY[20] REFERENCES EMPLOYEE)	
					then EMPLOYEE.EMPNO must be a column of array type, and teh constraint	
					says that the array value in MANAGER.MANAGES must either be null or be	
					equal to an array value in EMPLOYEE.EMPNO.	
					What is needed is a new syntax, perhaps:	
					CREATE TABLE MANAGER (
					EMPNO INTEGER ARRAY [20] FLEMENT REFERENCES	
					EMPLOYEE (EMPNO))	
					ELEMENT REFERENCES would mean that each array element of	
					MANAGER.MANAGES must either be null or equal value in	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
#	NLD-P02-049		2-Minor Technical	P02-11.10, <alter table<br="">statement></alter>	Description EMPLOYEE.EMPNO. ** Editor's Note (number 15) ** (Editor's note: In my opinion, Bill is simply trying to solve the problem using the wrong tools. INTEGER ARRAY[n] is meant to have elements of integers, not elements of employee IDswhich is a different thing altogether.) Solution None provided with comment. FND-747 The following Language Opportunity has been noted: Source: WG3:RTM-028/X3H2-99-252R1 Language Opportunity: It might be useful to have an option so that a conventional (SQL-92) table can evolve to become a table of type. However, any such proposal must avoid the pitfalls noted during development of SQL:1999 for evolution to a table of "named row type" (to use the terminology current before structured types were introduced). The proposal must account for the <reference specification="" type=""> of the user- defined type. If <reference generation=""> is DERIVED, it may be necessary to require a unique constraint or primary key constriant on the appropriate columns. If <references generation=""> is USER GENERATED, it may be</references></reference></reference>	Ву
					necessary to require that the table has no rows. Probably the self-referencing column must be added to the table as part of its evolution to a table of structured type. It is unlikely that the unaltered table will have as its first column a reference to the very type to which the table will be evolving. And, if perchance that condition were met, what would be do with the previously existing values in that column? Solution None provided with comment.	
	NLD-P02-050		1-Major Technical	P02-11.22, <view definition></view 	 FND-933 The following Possible Problem has been noted: Source: WG3:HBA-050R1 CR4) of this subclause is: 4) Without Feature F751, "View CHECK enhancements", conforming SQL language shall not contain <view definition=""> that contains a <subquery> and contains CHECK OPTION.</subquery></view> This is suspect. Even if it really was intended to rule out, as it apparently does, examples of the following form (note the <subquery>):</subquery> CREATE VIEW V AS SELECT FROM (SELECT FROM) AS T WHERE WITH CHECK OPTION then surely it should also be ruling out examples of the following equivalent form: CREATE VIEW V AS 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					WITH A AS (SELECT FROM) SELECT FROM A WHERE WITH CHECK OPTION But it doesn't. (Note the lack of any <subquery>.) A minor additional point is that "WITH CHECK OPTION" would be safer than just "CHECK OPTION", in case WITHOUT CHECK OPTION is ever added to the language. Solution None provided with comment.</subquery>	
	NLD-P02-051		1-Major Technical	P02-11.30, <drop domain<br="">statement></drop>	FND-938 The following Possible Problem has been noted: Source: WG3:HBA-028 Possible Problem: GR1)c) refers to "the explicit or implicit <constraint list="" name="">". The BNF production for <drop definition="" domain=""> does not included a <constraint name<br="">list>, nor do the Syntax Rules specify an implicit one in any circumstances. Solution Delete GR1)c) and edit the lead-in of GR1)d) as shown here: d) For every domain constraint descriptor included in the domain descriptor of D [begin deletion] whose <constraint name=""> is not contained in the excluded constraint list[end deletion]:</constraint></constraint></drop></constraint>	
	NLD-P02-052		2-Minor Technical	P02-11.39, <trigger definition></trigger 	FND-611 The following Language Opportunity has been noted: Source: DBL:LGW-146/X3H2-97-349 Language Opportunity: SQL3 should consider adding syntax to allow the user to specify the ordering in which triggers on the same effect should be fired. Solution None provided with comment.	
	NLD-P02-053		2-Minor Technical	P02-11.41, <user-defined type definition></user-defined 	 FND-603 The following Language Opportunity has been noted: Source: DBL:LGW-131/X3H2-97-293, 24 July, 1997; also USA-081 in first CD ballot for SQL/Foundation and WG3:YGJ-074/X3H2-99-164R1 Language Opportunity: Subclause 11.41, " user-defined type definition>", contains a Syntax Rule reading: 6)g) [A user-defined type] shall not be based on itself. This syntax rule prevents the UDT facility from modeling a recursively-defined data type such as "Tree". Here is a simple example of a UDT definition that is not possible because of that SR: CREATE TYPE Tree (node value INTEGER, 	

SEQ	Cmnt	See	Samarita.	Deferrer ee	Description	Addressed
#	ID	AISO	Severity	Keierence		Ву
					right subtree Tree)	
					Solution	
					None provided with comment.	
	NLD-P02-054		2-Minor	P02-11.50,	FND-713 The following Language Opportunity has been noted:	
			Technical	<sql-invoked< td=""><td>Source: WG3:YGJ-021</td><td></td></sql-invoked<>	Source: WG3:YGJ-021	
				routine>	Language Opportunity:	
					Currently all parameters must be of some specific concrete type. There needs to	
					be a mechanism to declare that a parameter is a character string of arbitrary,	
					unspecified type, at least when invoking PSM. (And there should be some	
					mechanism within PSM to interrogate the character set and length of a character	
					suring parameter). Otherwise the subject routine rules allow you to resolve to the	
					function is invoked, you will get an error when trying to assign the input	
					argument to the parameter's type if the input argument's character set is different	
					from the one declared in the function's signature. There should also be a	
					mechanism to declare that the return type of a function is determined by a	
					parameter's type.	
					Solution	
					None provided with comment.	
			1-Major	<i>P02-12.01</i> ,	WG3-P02-003	
			Technical	<grant< td=""><td>GRs 4)b), 4)c), 4)d), 4)e), 5), 6), though curiously not 7) all contain the phrase</td><td></td></grant<>	GRs 4)b), 4)c), 4)d), 4)e), 5), 6), though curiously not 7) all contain the phrase	
				statement>	"[f]ollowing the successful execution of the <grant statement="">". Given that</grant>	
					GRs are to be evaluated in the order in which they are written, and that	
					successful execution usually (flough flot always) means that the last GK has	
					Note that the BNF for $\langle grant statement \rangle$ specifies that it is either a $\langle grant$	
					privilege statement> or a <grant role="" statement="">. The rule evaluation order</grant>	
					specified in Framework, Subclause 6.3.3.3, makes it clear that the rules for the	
					contained statements are applied before the rules for the containing statement.	
					Perhaps, then, the wording we have questioned should be changed to something	
					to the effect of "following the successful execution of the contained statement";	
					but if the contained statement fails, then doesn't evaluation of the rules of the	
					containing statement end too? In that case, each of these phrases can simply be	
					detected. Note also that $(CP_{0}, 1), 2), 2), 4), 5), 6), and 7) all apositive the eventual of "the$	
					following (grant statement) "Whoever drafts a solution to this problem	
					might like to check that there is no infinite recursion going on here. It might he	
					that "the following <grant statement="">" should better be "the following <grant< td=""><td></td></grant<></grant>	
					privilege statement>" or "the following <grant role="" statement="">", as applicable.</grant>	
					Solution	
					None provided with comment.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
	NLD-P02-055		2-Minor Technical	P02-12.07, <revoke statement></revoke 	FND-734 The following Language Opportunity has been noted: Source: Email from Fred Zemke, 1999-06-09 and WG3:ZSH-155 = H2-2002-	
					Language Opportunity: The OLAP Amendment has created a new kind of dependency, of a view, etc., containing an OLAP function that references a user-defined ordering in its ORDER BY clause, which is dependent on the userdefined ordering. <drop routine="" statement=""> has been edited to account for this dependency; does any other statement need to be edited? Solution None provided with comment.</drop>	
	NLD-P02-056		2-Minor Technical	P02-12.07, <revoke statement></revoke 	 FND-911 The following Language Opportunity has been noted: Source: WG3:ZSH-155 = H2-2002 Language Opportunity: Syntax Rule 36) says: 36) If RESTRICT is specified, then there shall be no abandoned privilege descriptor, abandoned view, abandoned table constraint, abandoned assertion, abandoned domain constraint, lost domain, lost column, lost schema, and no descriptor that includes an impacted data type descriptor, impacted collation, impacted charater set, abandoned user-defined type, forsaken column decriptor, forsaken domain descriptor, or abandoned routine descriptor. This SR has several problems: It is unclear whether there should be a comma following "schema", though we recognize that a schems is a descriptor. (Note: This problem has been fixed by the addition of "and no" between "schema," and "descriptor".) It is unclear whether the object of "includes" is a nested list. (Note: This problem has been resolved by making it clear that it is a nested list.) The terms used to refer to impacted, etc., objects are inconsistent with those used to so designate them. While it is descriptors that are said to be abandoned, impacted, etc., this rule referes to "impacted columns", etc. Several possible candidates for inclusion in the list are absent for no obvious reason; they include abandoned table descriptor, abandoned trigger descriptor, and contaiminated column descriptor. 	
	NLD-P02-057		1-Major	P02-12.07,	FND-979 The following Possible Problem has been noted:	
	1.22 102 007		Technical	<revoke< td=""><td>Source: WG3:SIA-018 = H2-2003-429</td><td></td></revoke<>	Source: WG3:SIA-018 = H2-2003-429	
				statement>	Possible Problem:	
					WG3:SIA-018 introduced the notions of view components, view component privilege descriptors, and view privilege dependency descriptors pertaining to a	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					given view. WG3:SIA-018 proposed explicit rules that specify the creation of	
					view privilege dependency descriptors, but failed to specify explicit rules that	
					specify the destruction of view privilege dependency descriptors. Although a	
					view privilege dependency descriptor can be assumed to be destroyed whenever	
					either its supporting privilege descriptor or the dependent privilege descriptor is	
					destroyed, the standard would be clearer if this were done in the appropriate	
					place(s) in the GRs of Subclause 12.7, " <revoke statement="">".</revoke>	
					Solution	
					None provided with comment.	
	NLD-P02-058		2-Minor	P02-13.01,	FND-921 The following Language Opportunity has been noted:	
			Technical	<sql-client< td=""><td>Source: FCD1/2002, GBR-P02-485</td><td></td></sql-client<>	Source: FCD1/2002, GBR-P02-485	
				module	Language Opportunity:	
				definition>	None of the GRs in this Subclause relate to the creation of an SQL module.	
					Moreover, General Rule 4) relates to the invocation of an externally-invoked	
			procedure.			
					Solution	
					None provided with comment.	
	NLD-P02-059		2-Minor	P02-13.03,	FND-844 The following Language Opportunity has been noted:	
			Technical	<externally-< td=""><td>Source: WG3:YYJ-034 = H2-2001</td><td></td></externally-<>	Source: WG3:YYJ-034 = H2-2001	
				invoked	Language Opportunity:	
				proceaure>	The use of savepoint levels, introduced by WG3:PER-061 and extended by	
					WG3:YYJ-034, still does not cover the case of externally-invoked procedures.	
					Solution	
					None provided with comment.	
	NLD-P02-060		2-Minor	P02-13.05, <sql< td=""><td>FND-925 The following Possible Problem has been noted:</td><td></td></sql<>	FND-925 The following Possible Problem has been noted:	
			Technical	procedure	Source: WG3:HBA-029	
				statement>	Possible Problem:	
					Subclause 13.5, " <sql procedure="" statement="">", includes two GRs to the effect</sql>	
					that if the statement being executed is an atomic one, then all changes to SQL-	
					data and schemas are cancelled. Shouldn't this be conditional on whether the	
					current execution context is atomic, rather than on the statement type? Not that	
					this would make any material difference, but as things stand there appears to be	
					no point in the final sentence of GR2) of this Subclause:	
					1) 2) A statement execution context NEWSEC is established for the execution of	
					5. Let OLDSEC be the most recent statement execution context. NEWSEC	
					becomes the most recent statement execution context. NEWSEC is an atomic	
					execution context, and therefore the most recent atomic execution context, if and	
					Only it S is an atomic SQL-statement.	
					Autough there are OKS in various subclauses that to enforce atomicity where it is required, none of these rules references the atomicity or non-atomicity of an	
					avecution context	
		l			execution context.	

SEQ #	Cmnt	See	Soverity	Deference	Description	Addressed
#	ID	AISO	Severity	Kelerence	Description	Бу
					Solution Either delete all references to atomicity of execution contexts, or change GRs 5)a)ii) and 5)b)ii)1) of Subclause 13.5, " <sql procedure="" statement="">", to be conditional on the atomicity of the current statement execution context. Probably the former solution is to be preferred, in view of the specific GRs in several places that refer to savepoint levels and undoing changes to SQL-data and schemas. But see WG3:HBA-041.</sql>	
	NLD-P02-061		2-Minor Technical	P02-13.06, Data type correspondences	FND-815 The following Language Opportunity has been noted: Source: WG3:PER-107/H2-2001-115 Language Opportunity: Table 18, "Data type correspondences for COBOL", maintains that the COBOL type corresponding to BOOLEAN is PICTURE X. Before the deletion of the BIT type (by paper WG3:PER-107/H2-2001-115), Subclause 20.5, " <embedded SQL COBOL program>", maintained that the declaration "PIC X USAGE IS BIT" could be used either to correspond to a bit string or to a BOOLEAN. This was flawed because the embedded COBOL processor needs to know what SQL type to assign to an embedded variable declaration. After the deletion of the BIT type, there appears to be no support for BOOLEAN in Subclause 20.5, "<embedded cobol="" program="" sql="">", not even in a buggy Syntax Rule. Note that it will not do to overload "PICTURE X" as either CHAR(1) or BOOLEAN, for the same reason that it was not acceptable to overload "PIC X USAGE IS BIT" as either BIT(1) or BOOLEAN. Perhaps "USAGE IS BOOLEAN" is in order. Solution</embedded></embedded 	
					None provided with comment.	
	NLD-P02-062		I-Major Technical	P02-14.07, <delete statement: searched></delete 	 FND-939 The following Possible Problem has been noted: Source: WG3:HBA-028 Possible Problem: GR9) is as follows: 9) Each <subquery> in the <search condition=""> is effectively executed for each row of T and the results are used in the application of the <search condition=""> to the given row of T. If any executed <subquery> contains an outer reference to a column of T, then the reference is to the value of that column in the given row of T.</subquery></search></search></subquery> NOTE 496 — 368 "outer reference" is defined in Subclause 6.7, "<column reference="">".</column> As GR5) already says that the <search condition=""> is "applied to [sic evaluated for might be better] each row of T", perhaps GR9) isn't needed at all. If its existence is justified by the apparently inadequate definition of outer reference in the references duclause 6.7, then surely it would be better to fix SR4) of that Subclause to cater for outer references that are not contained in <table< li=""> </table<></search>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
					expression>s. The first sentence of GR9) refers to each <subquery> in the <search condition="">. It is questionable whether what it says is really applicable to every <subquery> contained in the <search condition="">, regardless of how deeply nested it is. In any case, the sentence is imprecise and inappropriately worded. We don't execute subqueries, and the meaning of "results are used in the application of" is unclear. Similar problems exist in the following rules: • Subclause 7.8, "<where clause="">", GR3) • Subclause 7.10, "<having clause="">", GR2) • Subclause 14.9, "<merge statement="">", GR6)a)i)1) • Subclause 14.11, "<update searched="" statement:="">", GR5)a)ii) and GR5)b)ii) Solution None provided with comment</update></merge></having></where></search></subquery></search></subquery>	
	NLD-P02-063		2-Minor Technical	P02-14.08, <insert statement></insert 	FND-715 The following Language Opportunity has been noted: Source: WG3:YGJ-021 Language Opportunity: When a row of a table that has a system-generated column is inserted, the application has no way to access the newly generated value. This was not an issue when only explicit values were inserted by the application. Solution None provided with comment.	
	NLD-P02-064		2-Minor Technical	P02-14.10, <update statement: positioned></update 	FND-717 The following Language Opportunity has been noted: Source: WG3:YGJ-021 Language Opportunity: The Format for <update target=""> does not provide a way to set a field of an anonymous row type. Seemingly the only way to update column of an anonymous row type is to replace the entire column, which will be awkward in many instances. For example, suppose I only want to update the STREET portion of an ADDRESS column. Looks like I have to use UPDATE T SET ADDRESS = ROW (:STREETVAR, T.CITY, T.STATE, T.ZIP); This means the query writer has to repeat the entire definition of the anonymous row in the query, which can be quite laborious, as well as hiding the simplicity of what the user is actually doing. Also, we must support all kinds of nesting of anonymous rows and UDTs. Solution None provided with comment.</update>	
	NLD-P02-065		2-Minor Technical	P02-14.10, <update statement: positioned></update 	FND-724 The following Language Opportunity has been noted:Source: WG3:FRA-093/X3H2-98-628)Language Opportunity:The <simple specification="" value=""> immediately contained in an <update target="">of a <set clause=""> specifying the array element of the target column to be updated</set></update></simple>	See comment

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					should be a <value specification=""> rather than a <simple specification="" value="">. This would allow the use of a <dynamic parameter="" specification=""> which is currently prohibited because a <simple specification="" value=""> cannot be a <dynamic parameter="" specification="">. General Rules 14)a)ii)5)c) of <update positioned="" statement:=""> and <update statement: searched> will cause an exception to be raised if a null value is passed as a <value specification=""> so no change is necessary to preclude a null value.</value></update </update></dynamic></simple></dynamic></simple></value>	
					Solution — Changes to Subclause 14.10, " <update positioned="" statement:="">": • Revise the BNF for <update target="">, replacing <simple specification="" value=""> with <value specification="">. • Replace <simple specification="" value=""> with <value specification=""> in Syntax Rule 10), General Rule 14) and Conformance Rule 2). — Changes to Subclause 14.11, "<update searched="" statement:="">": • Replace <simple specification="" value=""> with <value specification=""> in Syntax Rule 9) and General Rule 14).</value></simple></update></value></simple></value></simple></update></update>	
	NLD-P02-066		2-Minor Technical	P02-14.10, <update statement: positioned></update 	FND-809 The following Language Opportunity has been noted: Source: (was Possible Problem FND-737) WG3:PER-171/H2-2001-???, FCD1/2000 NLD-P02-063 (from WG3:YGJ-074/X3H2-99-164R1) Language Opportunity: There is no ability to truncate an array. Assigning NULL to the last element of an array does not decrease the length of the array. Solution None provided with comment.	
	NLD-P02-067		2-Minor Technical	P02-14.12, <set clause list></set 	FND-922 The following Language Opportunity has been noted: Source: WG3:ZSH-163 = H2-2003 Language Opportunity: Impossible to Update Different Parts of the Same Column SR 7) prohibits the same column name from appearing more than once in the list of SET clauses. This means that the user who wishes to use the shorthands available for assigning to fields of rows is rather severely restricted, unacceptably so, in our opinion. The problem does not arise in connection with assignment to attributes of UDT values, thanks to the ingenious SR 6). Solution None provided with comment.	
	NLD-P02-068		2-Minor Technical	P02-16.02, <set transaction statement></set 	FND-912 The following Language Opportunity has been noted: Source: WG3:ZSH-155 = H2-2002 Language Opportunity: The standard does not specify a maximum for <number conditions="" of="">. Presumably there is an implementation-defined or -dependent maximum value of <number conditions="" of="">. For example, we could add the following GR after</number></number>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					 GR 2): 2) If <number conditions="" of=""> exceeds an implementation-dependent maximum number of conditions, then an exception condition is raised: <i>invalid condition number</i>.</number> We must also add an entry in either the implementation-defined or the implementation-dependent Annex. Note: WG3:ICN-001 recorded "After some discussion, the consensus was that the condition should be a warning and that a good solution to the comment should involve adding an extra field to the diagnostics area, giving the current transaction's maximum number of conditions." 	
	NLD-P02-069		1-Major	P02-16.04. <set< td=""><td>FND-919 The following Possible Problem has been noted</td><td></td></set<>	FND-919 The following Possible Problem has been noted	
	NLD-1 02-009		Technical	constraints mode statement>	Source: WG3:ZSH-031R3 = H2-2002 Possible Problem: The subclause is silent with regard to the checking of constraints when the constraints mode is set to IMMEDIATE. Turning to Subclause 16.7, " <commit statement>", we see that there is an expectation that SET CONSTRAINTS ALL IMMEDIATE has the effect of checking all constraints and that this effect takes place between GR5) and GR6) of that subclause (as opposed to any vague notion of "at the end of the statement"). The implications for referential constraints that specify referential actions are not clear, especially in the case of referential actions that are triggering events. Solution None provided with comment.</commit 	
	NLD-P02-070		1-Major Technical	P02-16.04, <set constraints mode statement></set 	FND-940 The following Possible Problem has been noted: Source: WG3:HBA-028 Possible Problem: If a <set constraints="" mode="" statement=""> is used to change the current mode of some constraint from deferred to immediate, it might happen that the database fails to satisfy that constraint. In this case, an exception is raised, but the database remains unchanged, so every subsequent statement will fail with the same exception, apart from one that sets the relevant constraint's mode back to DEFERRED or one that makes some change to the database to return it to a consistent state. (One such statement is COMMIT, which turns itself into ROLLBACK if constraints are not satisfied.) Solution None provided with comment.</set>	
	NLD-P02-071		4-Minor Editorial	P02-16.05, <savepoint statement></savepoint 	 FND-973 The following Possible Problem has been noted: Source: WG3:SIA-031 = H2-2004-??? Possible Problem: General Rule 4) of this Subclause is: 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					 A savepoint is established in the current savepoint level and at the current point in the current SQLtransaction. S is assigned as the identifier of that savepoint. This is not sufficiently specific. It doesn't even say that sufficient data is preserved for the successful execution of a subsequent <rollback statement="">.</rollback> Solution Specify what happens in terms of the contents of the SQL-session context. 	
	NLD-P02-072		1-Major	P02-16.07,	FND-941 The following Possible Problem has been noted:	
			Technical	<commit statement></commit 	 Source: WG3:HBA-028 Possible Problem: Just before submitting this paper we discovered, in Subclause 16.7, "<commit statement="">", the following GR: 6) Case: a) If any constraint is not satisfied, then any changes to SQL-data or schemas that were made by the current SQL-transaction are canceled and an exception condition is raised: <i>transaction rollback — integrity constraint violation</i>. b) If the execution of any <triggered sql="" statement=""> is unsuccessful, then any changes to SQL-data or schemas that were made by the current SQL-transaction are canceled and an exception condition is raised: <i>transaction rollback — integrity constraint violation</i>.</triggered> c) If any other error preventing commitment of the SQLtransaction has occurred, then any changes to SQL-data or schemas that were made by the current SQL-transaction rollback with an implementation-defined subclass value. d) Otherwise, any changes to SQL-data or schemas that were made by the current SQL-transaction rollback with an implementation-defined subclass value. d) Otherwise, any changes to SQL-data or schemas that were made by the current SQL-transactions. This seems problematical. Case (a) is possibly okay, catering for any deferred constraints, though there is an opportunity to make it more precise using text similar to what HBA-028 proposed for constraint checking in Subclause 13.5. Regarding case (b), it is not clear how a <commit statement=""> can possibly cause a <triggered sql="" statement=""> to be invoked. Regarding case c), it is not clear what "other error preventing commitment" refers to. Perhaps an informative note is needed.</triggered></commit> </commit> 	
	NLD-P02-073		1-Major	P02-16.07,	FND-970 The following Possible Problem has been noted:	
			Technical	<commit statement></commit 	Source: WG3:SIA-023 = H2-2004-??? Possible Problem: After acceptance of WG3:SIA-023, Subclause 16.7, " <commit statement="">", GR 9)a) is: a) If <commit statement=""> contains AND CHAIN, then an SQL-transaction is</commit></commit>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					initiated. Any branch transactions of the SQL-transaction are initiated with the	
					same access mode, isolation level, and diagnostics area limit as the	
					corresponding branch of the SQL-transaction just terminated.	
					Before acceptance of WG3:SIA-023, it was:	
					a) If AND CHAIN was specified, then a new SQL-transaction is initiated with	
					the same access mode, isolation level, and diagnostics area limit as the SQL-	
					transaction just terminated. Any branch transactions of the SQL-transaction are	
					initiated with the same access mode, isolation level, and diagnostics area limit as	
					the corresponding branch of the SQL-transaction just terminated.	
					The simplification of the first sentence was made possible by WG3:SIA-023's	
					simpler approach to the setting of transaction characteristics. The problems lie in	
					the second sentence. It is not clear what "[a]ny branch transactions of the	
					SQLtransaction" refers to. It seems as if it refers to things that exist, and yet	
					"initiated" suggests that they are to be brought into existence. Also, no mention	
					is made of the initial constraint modes of branch transactions. Maybe the intent	
					is to specify that in every SQL-session containing a branch transaction of the	
					transaction just terminated, a branch transaction is initiated. But it is not clear	
					now branch transactions come into existence in the first place. Subclause 4.55,	
					sQL-transactions, mentions the possibility of their existence without	
					A second point that might need to be considered by environmenting to	
					A second point that high need to be considered by anybody attempting to address this P , concerns the initial constraint modes for the new transaction	
					initiated by AND CHAIN, GP5) clearly specifies that all constraint modes are	
					immediate by AND CHAIN. OKS) clearly specifies that an constraint modes are	
					was not really intended. It seems more intuitive to have constraint modes	
					reinitialised to their declared initial states, as when AND CHAIN is not	
					specified. If that was really the intent (and perhaps what has actually been	
					implemented by implementations, supporting Feature F721, "Deferred	
					constraints" then we might have to consider accepting an incompatible change	
					Solution	
					None provided with comment.	
	NLD-P02-074		1-Major	P02-16.07,	FND-976 The following Possible Problem has been noted:	
			Technical	<commit< th=""><th>Source: WG3:SIA-030 = H2-2004-???</th><th></th></commit<>	Source: WG3:SIA-030 = H2-2004-???	
				statement>	Possible Problem:	
					Subclause 16.7, " <commit statement="">", SR 3) is:</commit>	
					1) For every open cursor that is not a holdable cursor <i>CR</i> in any SQL-client	
					module associated with the current SQL-transaction, the following statement is	
					implicitly executed:	
					CLOSE CR	
					Exactly which cursors are to be closed under this rule is not clear, even when we	
					safely assume that it does not mean cursors other than holdable ones in any	
					SQL-client module associated with the current SQL-transaction. We surmise	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
#		AISU			that "in any SQL-client module associated with the current SQL-transaction" is redundant. We suspect that it is intended to refer to every (non-holdable) cursor open in the current SQL-session, but as it stands it appears to exclude any global extended dynamic cursor allocated during execution of the body of an SQL- invoked routine. Although Subclause 4.22, "SQL-client modules", does say that an SQL-client module includes "Zero or more cursors", there is no statement, there or elsewhere in Clause 4, "Concepts", to the effect that every cursor is in some sense "in" some SQL-client module. Subclause 4.35, "SQL-transactions", does contain the statement that "Each SQL-client module that executes an SQL-statement of an SQL-transaction is associated with that SQL-transaction", but it's not clear what it means for an SQL-client module to be associated (or not) with a particular SQL-transaction, nor is it clear what it means to say "Each SQL-client module that executes an SQL-statement", since statements executed in SQL-invoked procedures might or might not be included. The General Rules of Subclause 16.8, " <rollback statement="">", suffer from similar problems, and we additionally note a curious difference between GR2)e) ("All open cursors are closed") and its counterpart for the ROLLBACK TO SAVEPOINT case, GR3)g), which explicitly specifies execution of certain <close statement="">s. Possible Problem FND-975 describes another problem with the cited General Rule. It might be desirable to address both problems in a single change proposal. Solution None provided with comment.</close></rollback>	<u>Бу</u>
	NLD-P02-075		1-Major Technical	P02-18.01, <set session characteristics statement></set 	<pre>FND-971 The following Possible Problem has been noted: Source: WG3:SIA-023 = H2-2004-??? Possible Problem: The BNF production for <set characteristics="" session="" statement=""> is: <set characteristics="" session="" statement=""> ::= SET SESSION CHARACTERISTICS AS <session characteristic<br="">list> <session characteristic="" list=""> ::= <session characteristic=""> [{ <comma> <session characteristic> }] <session characteristic=""> ::= <transaction characteristics> According to this BNF, the following are both legal <set characteristics<br="" session="">statement>s: SET SESSION CHARACTERISTICS AS TRANSACTION READ ONLY, ISOLATION LEVEL SERIALIZABLE, DIAGNOSTICS SIZE 2 SET SESSION CHARACTERISTICS AS</set></transaction </session></session </comma></session></session></session></set></set></pre>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
#			Severity	Kelerence	DescriptionTRANSACTION READ ONLY, TRANSACTION ISOLATION LEVEL SERIALIZABLE, TRANSACTION DIAGNOSTICS SIZE 2If this strange-looking syntax was not actually intended, then it should be corrected.Solution <session characteristics="" transaction=""> shall contain at most one <isolation level="">, at most one <transaction access="" mode="">, and at most one <diagnostics size="">. <set characteristics="" session="" statement=""> ::= SET SESSION CHARACTERISTICS AS <session transaction<br=""></session>characteristics> <session characteristics="" transaction=""> ::= TRANSACTION <transaction characteristics=""> ::= TRANSACTION <transaction mode=""> [<comma> <transaction </transaction mode>] Replace the Syntax Rules by: 1) <session characteristics="" transaction=""> shall contain at most one <isolation </isolation level>, at most one <transaction access="" mode="">, and at most one <isolation </isolation level>, at most one <transaction access="" mode="">, and at most one <isolation </isolation level>, at most one <transaction access="" mode="">, and at most one <isolation </isolation level>, at most one <transaction characteristics=""> shall contain at most one <isolation </isolation level>, at most one <transaction access="" mode="">, and at most one <isolation </isolation level>, at most one <transaction characteristics="">. Let ESC be the enduring session characteristics of the current SQL-session. 2) If STC contains an <isolation level=""> IL, then the isolation level of ESC is set to the <level isolation="" of=""> contained in IL. 3) If STC contains an <a>access mode> AM, then the access mode of ESC is set read-only or read-write according to whether AM contains READ ONLY or READ WRITE, respectively. 4) If STC contains a <diagnostics size=""> DS, then the condition area limit of ESC</diagnostics></level></isolation></transaction></transaction></transaction></transaction></transaction></transaction></session></comma></transaction></transaction></session></set></diagnostics></transaction></isolation></session>	Бу
	NLD-P02-076		1-Major Technical	P02-18.02, <set session user identifier statement></set 	 Is set to the <number conditions="" of=""> contained in DS.</number> FND-977 The following Possible Problem has been noted: Source: WG3:SIA-026R3 = H2-2004-??? Possible Problem: In SQL/Foundation, GR5) is: 5) If the current user identifier and the current role name are restricted from setting the user identifier to V, then an exception condition is raised: <i>invalid authorization specification</i>. It is not clear how to interpret GR5) in the case where current user and current role do not both exist. Furthermore, suppose they do both exist and just one of them is restricted from setting "the user identifier" to V. The rule is written to require both of them to be so restricted for that exception condition to be raised. This seems a little arbitrary and we wonder if that was what was really intended. Solution None provided with comment. 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
#	ID NLD-P02-077 NLD-P02-078	Also	Severity 2-Minor Technical 2-Minor Technical	Reference P02-19, Dynamic SQL P02-19.06, <prepare statement=""></prepare>	Description FND-726 The following Language Opportunity has been noted: Source: WG3:FRA-126R1 and WG3:PER-098R1/H2-2001-059 Language Opportunity: There is no way to retrieve a locator for an array, a multiset, or a UDT without having pre-knowledge of the type of data to be accessed because the rules for <get descriptor="" statement=""> require that the data type of the <simple specification="" target=""> "match" that represented by the item descriptor area when retrieving DATA. For UDT locators, "match" implies that the UDT for which the locator was declared be the same as that specified in the SQL item descriptor area. For array locators and multiset locators, "match" implies that the element data types be the same. The only way to declare a host variable appropriately is to know in advance what UDTs, arrays, or multisets will be accessed. This is unacceptable for dynamic SQL. A similar problem exists with reference types. Solution None provided with comment. FND-926 The following Possible Problem has been noted: Source: WG3:HBA-040 Possible Problem: General Rule 10) is: 1) 10) If <statement name=""> is specified for the <sql name="" statement="">, P is not a <cursor specification="">, and <statement name=""> is associated with a cursor C through a <dynamic cursor="" declare="">, then an exception condition is raised: dynamic SQL error — prepared statement name or a cursor specification</dynamic></statement></cursor></sql></statement></simple></get>	By
					dynamic SQL error — prepared statement not a cursor specification. This rule is redundant: all it does is warn the user that he won't be able to open the dynamic cursor; unless, of course, he subsequently executes a <prepare statement> with the same <statement name=""> and an <sql statement="" variable=""> whose value is a <cursor specification="">. The check belongs on <dynamic open<br="">statement>. But see WG3:HBA-041. Solution None provided with comment.</dynamic></cursor></sql></statement></prepare 	
	NLD-P02-079		1-Major	<i>P02-19.11</i> ,	FND-949 The following Possible Problem has been noted:	
			Technical	<output using<br="">clause></output>	Source: WG3:HBA-048 = H2-2003 Possible Problem: General Rule 3) of this Subclause is: Case: a) If PS is a <dynamic select="" statement=""> or a <dynamic p="" row="" select<="" single=""> statement>, then the <output clause="" using=""> describes the <target specification="">s for the <dynamic fetch="" statement=""> or the <execute statement="">. Let <i>D</i> be the degree of the table specified by <i>PS</i>. The use here of the BNF non-terminal <target specification=""> is inappropriate in the case that <into descriptor=""> is specified.</into></target></execute></dynamic></target></output></dynamic></dynamic>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					Solution	
					None provided with comment.	
	NLD-P02-080		1-Major	<i>P02-19.11</i> ,	FND-950 The following Possible Problem has been noted:	
			Technical	<i><output i="" using<=""></output></i>	Source: WG3:HBA-048 = H2-2003	
				clause>	Possible Problem:	
					General Rule 6)c) of this Subclause is:	
					If the <output clause="" using=""> is used in a <dynamic fetch="" statement="">, then let</dynamic></output>	
					LTDT be the data type on the most recently executed <dynamic fetch<="" td=""><td></td></dynamic>	
					statement>, if any, for the cursor CR. It is implementation-defined whether or	
					not an exception condition is raised: <i>dynamic SQL error</i> — <i>restricted data typ</i>	
					attribute violation if any of the following are true:	
					fatch statements, if any for the cursor CP" is intended to refer to There is no	
					data type on (or even in or of) a <dynamic fetch="" statement=""></dynamic>	
					We suspect that what is meant is:	
					It is implementation-defined whether or not you're allowed to fetch into a locator	
					on one fetch from CR, but not on the next, or vice versa.	
					Moreover, Annex B, "Implementation-defined elements" contains no entry for	
					this Subclause.	
					Solution	
					None provided with comment.	
	NLD-P02-081		1-Major	<i>P02-19.11</i> ,	FND-951 The following Possible Problem has been noted:	
			Technical	<i><output i="" using<=""></output></i>	Source: WG3:HBA-048 = H2-2003	
				clause>	Possible Problem:	
					General Rule 6) of this Subclause contains two subrules that cause a locator to	
					be generated. The structure is:	
					$\begin{array}{c} 6) \text{ For } 1 \text{ (one) } \delta 1 \delta \text{ D}; \\ C \end{array}$	
					Case:	
					1) If SV is not the null value, then a locator L that uniquely identifies SV is	
					generated and is the value TV of the i-th \leq target specification>	
					Case:	
					If <into descriptor=""> is specified, then</into>	
					Case:	
					Otherwise, [TVT is assumed to be a locator of some sort]	
					Case:	
					If TV is not the null value, then:	
					Case:	
					If TYPE indicates a locator type, then a locator L that uniquely identifies TV is	
					generated and the value of DATA is set to an implementation-dependent four-	
					octet value that represents L.	
					This appears to be generating a locator of a locator. And why " an	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					implementation-dependent four-octet value that represents"?	
					Solution	
					None provided with comment.	
	NLD-P02-082		2-Minor	P02-19.11,	FND-952 The following Language Opportunity has been noted:	
			Technical	<i><output i="" using<=""></output></i>	Source: WG3:HBA-048 = H2-2003	
				19.10. <input< td=""><td>Language Opportunity:</td><td></td></input<>	Language Opportunity:	
				using clause>	Subclause 19.10, " <input clause="" using=""/> ", Syntax Rule 1) is:	
					arguments shall be either a host parameter specifications or an <embedded< td=""><td></td></embedded<>	
					variable specification>.	
					and Subclause 19.11, " <output clause="" using="">", Syntax Rule 1) is:</output>	
					1) The <target specification=""> immediately contained in <into argument=""> shall be</into></target>	
					either a <host parameter="" specification=""> or an <embedded td="" variable<=""><td></td></embedded></host>	
					specification>.	
					It is thus not currently possible for an SQL parameter to be either a <using< td=""><td></td></using<>	
					argument> or an <into argument="">.</into>	
					Solution	
	NI D D02 002		1 Maina	D 02 10 17	None provided with comment.	
	NLD-P02-083		1-Major Technical	<pre>P02-19.17, <dvnamic fetch<="" pre=""></dvnamic></pre>	FND-948 The following Possible Problem has been noted: Source: WG2: HPA $0.48 = H2,2003$	
			Teennear	statement>	Possible Problem:	
					General Rule 2) of this Subclause is:	
					2) The General Rules of Subclause 19.11, " <output clause="" using="">", are applied</output>	
					to the <dynamic fetch="" statement=""> and the current row of <i>CR</i> as the retrieved</dynamic>	
					row.	
					Subclause 19.11, " <output clause="" using="">", doesn't mention either "the current</output>	
					row" or "the retrieved row".	
					Solution The investions of the Connect Dates of this Sate large should be received	
	NI D P02 084		1 Major	P02-19-22	END 030 The following Possible Problem has been noted:	
	NLD-1 02-004		Technical	<preparable< pre=""></preparable<>	Source: WG3:HBA-040	
				dynamic delete	Possible Problem:	
				statement:	Both subclauses contain <scope option=""> in the Format, yet say nothing about it</scope>	
				<i>P02-19.23</i>	in either Syntax Rules or General Rules.	
				<preparable< td=""><td>Both subclauses contain a Syntax Rule:</td><td></td></preparable<>	Both subclauses contain a Syntax Rule:	
				dynamic update	2) All Syntax Rules of Subclause 14.n, " <xx positioned="" statement:="">", apply to</xx>	
				statement:	the <preparable dynamic="" positioned="" statement:="" xx="">, replacing "<declare gurgers " with "<dupamic "="" and<="" collegate="" declare="" gurger="" gurgers="" or="" statements="" td=""><td></td></dupamic></declare </preparable>	
				positionea>	" <xx nositioned="" statement:="">" with "<pre>cursor</pre> or <anocate cursor="" statement:<br="">"<xx nositioned="" statement:="">" with "<pre>cursor</pre> or <anocate cursor="" pre="" statement:<=""></anocate></xx></anocate></xx>	
					positioned>".	
					Neither <xx positioned="" statement:=""> refers to a <declare cursor=""> (they did once),</declare></xx>	
					and if they did, it is difficult to understand how <allocate cursor="" statement=""></allocate>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					could be relevant, since the cursor it creates cannot be referenced by a <cursor< td=""><td></td></cursor<>	
					name>.	
					Solution	
					None provided with comment.	
	NLD-P02-085		2-Minor	P02-20.01,	FND-364 The following Language Opportunity has been noted:	
			Technical	<embedded sql<="" td=""><td>Source: WG3:YGJ-074/X3H2-99-164R1</td><td></td></embedded>	Source: WG3:YGJ-074/X3H2-99-164R1	
				host program>	Language Opportunity:	
					There is a problem for precompilers when the issue of overlapping and non-	
					disjoint scopes for host variables, etc. comes into play. In addition, there are	
					problems caused by things like C macros and the C #ifdef conditional facilities.	
					Solution	
					None provided with comment.	
	NLD-P02-086		1-Major	P02-20.01,	FND-770 The following Possible Problem has been noted:	
			Technical	<embedded sql<="" td=""><td>Source: WG3:BHX-166</td><td></td></embedded>	Source: WG3:BHX-166	
				nosi program>	Possible Problem:	
					Since multiple SQL data types map onto the same C data type in Table 17,	
					"Data type correspondences for C", in Subclause 13.6, "Data type	
					correspondences", SR22) of Subclause 20.1, " <embedded host="" program="" sql="">",</embedded>	
					cannot correctly identify the corresponding SQL data type of a given C data	
					type. The problem identified is severed by Table 17 "Date type correspondences for	
					C" in Subclause 13.6 "Data type correspondences" that defines the mapping of	
					C data types onto SOL data types. The table maps more than one SOL data type	
					onto the same C data type. Hence, when the mapping table is used in reverse, a	
					single C data types mans onto more than one SOL data type. Now in case of	
					syntax rule 22) of Subclause 20.1. " <embedded host="" program="" sol="">", the SOL</embedded>	
					data type has to be determined while an <embedded host="" program="" sol=""> is</embedded>	
					processed. Thus, the SQL data types can only be derived syntactically from the	
					C data types based on Table 17, "Data type correspondences for C", in	
					Subclause 13.6, "Data type correspondences".	
					The solution of the problems would require a change of Table 17, "Data type	
					correspondences for C", in Subclause 13.6, "Data type correspondences", such	
					that a single SQL data type maps onto a single C data type. There might be an	
					alternative solution which accesses the definition of a routine to find out the	
					SQL data types rather than using the mentioned table. Both solutions result in	
					major changes of the document and might also lead to compatibility issue.	
					Hence, a real solution of the identified problems cannot be developed in the	
					given timeirame.	
					Solution	
	NI D D02 007		2.14	DO2 20 05	None provided with comment.	
	NLD-P02-087		2-Minor	PUZ-20.05,	FND-94/ The following Language Opportunity has been noted:	
			Technical	<embedded sql<="" td=""><td>Source: LO arising from WG3:HBA-038 = H2-2003-294</td><td></td></embedded>	Source: LO arising from WG3:HBA-038 = H2-2003-294	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
				COBOL program>	Language Opportunity: With the publication of COBOL 2002, there are opportunities for exploiting the new features in COBOL 2002 in specifying the data type correspondences for COBOL. Mapping SQL user-defined types to object capabilities in COBOL 2002 should also be investigated. Solution None provided with comment.	
	NLD-P02-088		2-Minor Technical	P02-22.01, <get diagnostics statement></get 	WG3-P02-004 GR6)b) appears to assume that a <get diagnostics="" statement=""> specifies a single assignment, whereas in general it can specify several, these being possibly of both statement information items and condition information items. Some kind of "for each" construct is needed in the phrasing of this rule. Arguably such treatment should really be applied to GRs 2) onwards, so that they become subrules of a single outermost rule, but it might be considered acceptable to let them stand and just fix GR6). Solution None provided with comment.</get>	
	NLD-P02-089	1	-Major	<i>P02-24,</i>	Feature F121 Basic diagnostics management (or at least sufficient to return the	
		ſ	Fechnical	Conformance	information inherent in F491) should be included in Core SQL. Solution None submitted with comment	
	NLD-P02-090	1 1 1	-Major Fechnical	P02-24, Conformance	Feature F391 Long Identifiers should be included in Core SQL. Solution None submitted with comment	
	NLD-P02-091	1 7	-Major Fechnical	P02-24, Conformance	Feature F491 Constraint management should be included in Core SQL. Solution None submitted with comment	
	NLD-P02-092	1 	-Major Fechnical	P02-24, Conformance	Feature T051 Row types should be included in Core SQL. Solution None submitted with comment	
	NLD-P02-093	1 1 1	-Major Fechnical	P02-24, Conformance	Feature T141 SIMILAR predicate should be included in Core SQL. Solution None submitted with comment	
	NLD-P02-094	1 7	-Major Fechnical	P02-24, Conformance	Feature T351 Bracketed SQL comments should be included in Core SQL. Solution None submitted with comment	
	NLD-P02-095	1	-Major Fechnical	P02-24, Conformance	UNICODE as a mandatory character set should be included in Core SQL. Solution	
			1 Major	PO2-E SOL	INORE SUDDIFIELD WITH COMMENT	
	NLD-P02-090		Technical	feature taxonomy	Source: WG3:HBA-050R1	
			Teennear		Possible Problem:	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					In Table 35, "Feature taxonomy and definition for mandatory features", row 134, the Description for Feature F131, "Grouped operations" is: — A grouped view is a view whose <query expression=""> contains a <group by<br="">clause> This contradicts the definition of grouped view that existed vacuously in SQL:1999 and has since been deleted. Furthermore, it doesn't seem to be an accurate summary of what Feature F131, "Grouped operations" really is. See FIPS 127-2, feature 13, for the proper definition. Solution</group></query>	
	NLD-P02-097		1-Major Technical	P02-No particular location	FND-772 The following Possible Problem has been noted: Source: WG3:BHX-118 Possible Problem: The proposal accepted in WG3:BHX-118 creates a new problem. It makes is possible for an externallyinvoked procedure invoked directly from the SQL- client to define a WITH RETURN cursor that is left open when the externally- invoked procedure returns to the SQL-client. This is at best meaningless, since the SQL-client has no way to do anything with that cursor, and at worst causes a problem with resource "leaks" related to unclosed cursors. Solution None provided with comment.	
	NLD-P02-098		2-Minor Technical	P02-No particular location	FND-918 The following Possible Problem has been noted: Source: WG3:ZSH-034R1 = H2-2002 Possible Problem: What does CURRENT_ROLE tell us? During execution of an SQL routine <i>R</i> whose security characteristic is DEFINER, an invocation of CURRENT_ROLE will return the authorization identifier (i.e., the role name) of the owner of <i>R</i> . If it were considered that a user might be interested in knowing what role was actually set by the most recent <set role="" statement="">, then we would need a SESSION_ROLE, analogous to SESSION_USER. Solution None provided with comment.</set>	
	NLD-P02-099		2-Minor Technical	P02-No specific location	FND-014 The following Language Opportunity has been noted: Language Opportunity: It was noted in conjunction with CAN-106 discussions that if one inserts a row in a view V1 but do not have INSERT privilege on the underlying view V2 that has a WITH CHECK OPTION constraint, then a <i>constraint violation</i> exception is raised; however, one can then not discover anything about that constraint! Solution	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					None provided with comment.	
	NLD-P02-100		2-Minor	P02-No specific	FND-055 The following Language Opportunity has been noted:	
			Technical	location	Language Opportunity:	
					It has been noted that schema manipulation requires no privileges, but depends	
					directly on ownership of the schema.	
					Solution	
					None provided with comment.	
	NLD-P02-101		2-Minor	P02-No specific	FND-129 The following Language Opportunity has been noted:	
			Technical	location	Language Opportunity:	
					[Note from SLC] We use the terms "destroyed", "deallocated", "deleted",	
					"released", and perhaps others in various places. Are these terms used	
					consistently and could the number of such terms be reduced?	
					Solution	
					None provided with comment	
	NLD-P02-102		2-Minor	P02-No specific	FND-134 The following Language Opportunity has been noted:	
	1122 102 102		Technical	location	Language Opportunity:	
			reenneur		[Note from SLC] The functions LOWER and UPPER might be better defined in	
					terms of translations and collations so that they properly account for all	
					character sets instead of only <simple character="" latin="">s.</simple>	
					Solution	
					None provided with comment	
	NI D P02 103		2 Minor	P02-No specific	FND 100 The following Language Opportunity has been noted:	
	NLD-102-105		Z-Winton Technical	location	Source: Jim Melton	
			reennear		I anguage Apportunity:	
					Lim Melton said in his response to TC LB X3H2-90-267.	
					We believe that many implementations will have schema objects other than	
					those specified in SOL2 (e.g. indexes stored < module>s etc.) that may depend	
					on schema objects defined in SQL2. The DROP semantics for such	
					implementations will depend on those implementation-defined objects as well as	
					those specified in SOL ₂ , yet the SOL ₂ DROP rules do not appear to make	
					allowances for additional restrictions on DROP statements. The wording in	
					SOL2 must be enhanced to allow for such additional restrictions.	
					Paper X3H2-90-373 addressed this, but failed, X3H2 suggested that a broader	
					proposal that addresses the general concept of implementation-defined objects	
					that might restrict CASCADE operations would be acceptable.	
					Solution	
					None provided with comment	
	NLD-P02-104		2-Minor	P02-No specific	FND-212 The following Language Opportunity has been noted:	
	102-104		Technical	location	Source I ON-034/X3H2-90-333 1	
			recimical		Language Onnortunity	
					The ISO SOL2 Editing Meeting in London noted that with the advent of a	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					default character set for domains and columns in a schema, there is an	
					opportunity to change that default character set for the schema. This might, for	
					example, involve an ALTER SCHEMA CHANGE CHARACTER SET	
					statement.	
					Solution	
					None provided with comment.	
	NLD-P02-105		2-Minor	P02-No specific	FND-217 The following Language Opportunity has been noted:	
			Technical	location	Source: Stephen Cannan	
					Language Opportunity:	
					Steve Cannan has noted:	
					It might be necessary to redefine the actions of triggers so that certain actions	
					survive an <i>unsuccessful</i> execution of an SQL statement. For example, a	
					BEFORE DELETE trigger might be used to record <i>attempts</i> to alter a table for	
					security reasons. It would therefore be necessary that the triggered action	
					survive an error in the original statement.	
					Solution	
	NH D D00 104		2.25	D02 N :C	None provided with comment.	
	NLD-P02-106		2-Minor	P02-No specific	FND-241 The following Language Opportunity has been noted:	
			Technical	iocuiion	Language Opportunity:	
					[From London] The following Opportunity exists:	
					when counting the number of rows affected by an <sql affected="" are="" by="" consider="" counting="" might="" one="" rows="" statements,="" td="" that="" the="" too<="" triggered=""><td></td></sql>	
					(a g triggers and referential constraints)	
					(c.g., triggers and referential constraints).	
					None provided with comment.	
	NLD-P02-107		2-Minor	P02-No specific	FND-242 The following Language Opportunity has been noted:	
			Technical	location	Language Opportunity:	
					[From London] The following Opportunity exists:	
					For language consistency, a correlation name should be permitted for the	
					modified table in positioned and searched update and delete statements.	
					Solution	
					None provided with comment.	
	NLD-P02-108		2-Minor	P02-No specific	FND-268 The following Language Opportunity has been noted:	
			Technical	location	Language Opportunity:	
					During consideration of YOK-023/X3H2-92-252, following language	
					opportunity was identified:	
					The set of <identifier>s available as <regular character="" identifier="" set="">s in the</regular></identifier>	
					<similar predicate=""> (see Subclause 8.6, "<similar predicate="">") could profitably</similar></similar>	
					be enhanced to support additional character attributes (e.g., ideographs,	
					syllables, etc., as a result of internationalization work subh as that going on in	
					SC22/WG20.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					Solution	
					None provided with comment.	
	NLD-P02-109		2-Minor	P02-No specific	FND-309 The following Language Opportunity has been noted:	
			Technical	location	Source: Phil Shaw	
					Language Opportunity:	
					Local declarations of dynamic cursor names would seem like a straightforward	
					extension to X3H2-93-056/YOK-034rev.	
					Solution	
					None provided with comment.	
	NLD-P02-110		2-Minor	P02-No specific	FND-317 The following Language Opportunity has been noted:	
			Technical	location	Source: X3H2-93-445/MUN-160	
					Language Opportunity:	
					The representation of SQL-paths in the Information Schema needs to be	
					specified.	
					Solution	
	NI D D00 111		2.14	DO2 No model(None provided with comment.	
	NLD-P02-111		2-Minor	P02-No specific	FND-32/ The following Language Opportunity has been noted:	
			Technical	iocuion	Source: A SH2-95-570R1/MUIN-170	
					Chiege Opportunity:	
					need the ability to add an existing object to a type or to remove it from a type	
					without destroying the object. Existing persons become employees and later stop	
					being employees while continuaing to exist as persons. This can be achieved	
					with a modest extension of current facilities.	
					The paper went on to add that a simple extension would be allow a constructor	
					such as STUDENT() to accept an optional parameter whose value is an existing	
					object that is to be made an instance of STUDENT (but only if it is in the type	
					hierarchy with STUDENTs).	
					Solution	
					None provided with comment.	
	NLD-P02-112		2-Minor	P02-No specific	FND-426 The following Language Opportunity has been noted:	
			Technical	location	Source: Paper X3H2-94-528/DBL:RIO-081 noted the following Possible	
					Problem;	
					WG3:BBN-155/X3H2-98-378 changed it to a Language Opportunity:	
					Language Opportunity: This possibility (factoring out parts of coolumn definition) chief definition	
) was pointed out as an opportunity in SOLL076, and we considered	
					attempting it However although there seemed to be no problem with the RNF	
					we were unsure how to specify a default character set. Consider Syntax Rule 6)	
					of <column definition="">, which reads:</column>	
					6) If a <data type=""> is specified, then:</data>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					 a) Let <i>DT</i> be the <data type="">.</data> b) If DT is CHARACTER, CHARACTER VARYING, or CHARACTER LARGE OBJECT and does not specify a <character set="" specification="">, then the <character set="" specification=""> specified or implicit in the <schema character="" set="" specification=""> of the <schema definition=""> that created the schema identified by the <schema name=""> immediately contained in the of the containing or <alter statement="" table=""> is implicit.</alter></schema></schema></schema></character></character> c) If <i>DT</i> is a <character string="" type=""> that identifies a character set that specifies a <collate clause=""> and the <column definition=""> does not contain a <collate clause="">, then the <column definition=""> does not contain a <collate clause="">, then the <collate clause=""> of the <character string="" type=""> is implicit in the <column definition="">.</column></character></collate></collate></column></collate></column></collate></character> Now, apart from the fact that this masterpiece of prolicity probably has more angle brackets than it should have, it just doesn't seem to work anyway for a LOCAL DECLARED TABLE (which has MODULE instead of a <schema name="">).</schema> Furthermore, the Syntax Rules for <sql declaration="" variable=""> (in RIO-006, SQL/PSM) contain nothing corresponding to this rule. If it's needed here, is it not also needed there?</sql> We seem to need something rather more generic, such as "the character set of the relevant schema". The difficulty is specifying what we mean by "relevant" so as to cover all cases, but it should surely be possible. 	
	NLD-P02-113		2-Minor Technical	P02-No specific location	FND-440 The following Language Opportunity has been noted: Source: Paul Cotton noted the following Language Opportunity in Ottawa, July, 1995 Language Opportunity: DBL:YOW-027 changed Subclause 13.4, "Calls to an <externally-invoked procedure>", to define BOOLEAN parameters as zero (0) for FALSE and one (1) for TRUE for the C language. However, Subclause 6.12, "<cast specification="">", does not currently permit BOOLEAN source values to be cast to a target value of type exact numeric. This would appear to be inconsistent with the abovereferenced change. An opportunity exists to permit this cast. Solution None provided with comment.</cast></externally-invoked 	
	NLD-P02-114		2-Minor Technical	P02-No specific location	 FND-452 The following Language Opportunity has been noted: Source: DBL:YOW-102/X3H2-95-244 discussion Language Opportunity: The specification of the isolation levels is less precise and rigorous than it should be; as a result, the intent is somtimes misperceived and the details are often imsinterpreted. 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					Solution	
					None provided with comment.	
	NLD-P02-115		2-Minor	P02-No specific	FND-453 The following Language Opportunity has been noted:	
			Technical	location	Source: Steve Cannan noted the following Language Opportunity during	
					discussion of DBL:YOW-055/X3H2-95-140:	
					Language Opportunity:	
					Rules such as Subclause 11.10, " subclause 11.10, "key subclause 11.10, "	
					if the phrase "identified by" was defined to require evistence.	
					Solution	
					Solution	
	NI D D02 116		2 Minor	P02-No specific	FND 468 The following Language Opportunity has been noted:	
	NLD-F02-110		Z-Willor Technical	location	Source: X3H2-94-103/DBL:SOU-076	
			Teennear		Language Onnortunity	
					X3H2-94-103/DBL:SOU-076 only introduced a ROW TYPE for SOL (i.e., for	
					SQL variables, parameters, results, and columns). The host language data types	
					are still the scalar types specified in SQL-86, SQL-89, and SQL-92. Thus, the	
					proposal doesn't add the new SQL ROW_TYPE to the host language mappings	
					for module language, embedded syntax, or external routine parameters.	
					Support for host language ROW_TYPEs would require specifying the forms of	
					host language record declarations that are recognized in embedded syntax, and	
					adding such host language record types to the data type correspondences for	
					embedded syntax, module language, and external routines.	
					Such a proposal would presumably include the ability to reference such host	
					language variables as targets of FETCH, SELECT, and assignment statements,	
					as sources of INSERT, UPDATE, and assignment statements, and as arguments	
					See also I anguage Opportunities PSM-078 and CI I-003 BIND-003	
					See also Language Opportunities 1 SM-078, and CLI-005, DIAD-005.	
					None provided with comment	
	NLD-P02-117		2-Minor	P02-No specific	FND-469 The following Language Opportunity has been noted:	
	NLD 102 117		Technical	location	Source: X3H2-94-103/DBL:SOU-076	
			Teenneur		Language Opportunity:	
					SQL3 table definitions include a new LIKE clause that lets you "copy" column	
					definitions from existing tables:	
					CREATE TABLE EMP_DEPT (LIKE EMP, LIKE DEPT,	
					OTHER_COLUMN CHAR(5))	
					A similar clause would seem useful for ROW_1YPE declarations. The clause	
					would, nowever, need to be generalized somewhat to allow for specifying fow	
					Calution	
					Solution	
			1		None provided with comment.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
	NLD-P02-118		2-Minor	P02-No specific	FND-470 The following Language Opportunity has been noted:	_
			Technical	location	Source: X3H2-94-103/DBL:SOU-076	
					Language Opportunity:	
					As noted in PP FND-469, the LIKE clause provides a shorthand for creating	
					tables of similar formats.	
					As described in X3H2-94-105/DBL:SOU-0/0, this proposal includes the ability to specify a POW. TVPE as a DOMAIN or a DISTINCT TVPE (this results	
					from defining ROW_TYPE as a \angle data type>). A possible follow-on proposal	
					could extent CREATE TABLE to allow reference to ROW TYPE domains	
					and/or types:	
					CREATE DOMAIN NAME AS ROW_TYPE (FIRST CHAR(10), LAST	
					CHAR(10));	
					CREATE TABLE OF NAME; There are several detailed questions that such a proposal would need to address	
					FOr example, can domain names and LIKE both be used in a CREATE	
					TABLE? Can a DISTINCT TYPE he used in a CREATE	
					TABLE?	
					Solution	
					None provided with comment.	
	NLD-P02-119		2-Minor	P02-No specific	FND-471 The following Language Opportunity has been noted:	
			Technical	location	Source: X3H2-94-103/DBL:SOU-076	
					Language Opportunity:	
					Given two rows, R1 and R2, a "concatenation" or "join" operator could be	
					defined. For discussion, assume that it would be written with the operator .	
					Then, if R1 has F1 fields and R2 has F2 fields, R1 R2 would yield a row	
					with $F1+F2$ fields, where the values of the first F1 fields are the values of the	
					Solution	
					Solution	
	NI D D02 120		2 Minor	PO2 No specific	FND 472 The following Language Opportunity has been noted:	
	NLD-102-120		Technical	location	Source: X3H2-94-103/DBL:SOU-076	
			reenneur		Language Opportunity:	
					According to this paper, two ROW_TYPEs are equivalent (and assignable) if	
					both have the same number of fields and every pair of fields in the same position	
					have compatible types.	
					A possible follow-on could consider an option for assignment and type	
					equivalence rules based on the names (instead of the positions) of the fields,	
					similar to the <corresponding specification=""> of <query expression="">s.</query></corresponding>	
					Solution	
			2.16	DO2 N	None provided with comment.	
	NLD-P02-121		2-Minor	P02-No specific	FND-4/4 The following Language Opportunity has been noted:	
		1	recnnical	iocunon	Source: A3H2-94-105/DBL:SOU-0/0	1
SEQ	Cmnt	See				Addressed
-----	-------------	------	----------------------	-----------------------------	--	-----------
#	ID	Also	Severity	Reference	Description	By
					Language Opportunity: A possible follow-on paper could extend the definition of ROW_TYPEs to allow constraints and default values. Solution	
					None provided with comment.	
	NLD-P02-122		2-Minor Technical	P02-No specific location	FND-475 The following Language Opportunity has been noted: Source: X3H2-94-103/DBL:SOU-076 Language Opportunity: A possible follow-on paper could integrate the rules for ROW_TYPE comparisons in predicates into one single Subclause. Solution None provided with comment	
	NLD-P02-123		2-Minor Technical	P02-No specific location	None provided with comment. FND-519 The following Language Opportunity has been noted: Source: X3H2-96-111/DBL:MCI-098 Language Opportunity: The TRIGGERED_COLUMNS base table in the Definition Schema misses an opportunity to capture both the explicit UPDATE columns of a trigger and other explicit or implicit "referenced" columns of the trigger. Consider replacing the "TRIGGERED_COLUMNS base table" in the current specification with the following new base table and view: TRIGGER_COLUMN_USAGE base table This table would consist of 8 columns (instead of the 7 columns in the existing TRIGGERED_COLUMNS base table). 3 columns to identify the Catalog, Schema, and Name of a Column. 1 column to indicate whether the named column is an explicit UPDATE column (specified in the <trigger column="" list=""> of an UPDATE <trigger event=""> of this trigger), an explicit "Contained" column (iontained in the <triggered action=""> of this trigger), or an "Implicit" column (implicitly referenced because it happens to be a column in the subject table of an UPDATE Trigger specified without an explicit <trigger column="" list="">). This 8-th column Could also be used later to identify other kinds of column usage that may be the basis of a <trigger event="">, e.g. SELECT (if triggers are extended to SELECT actions), or the actual column (or columns) that get updated by an INSTEAD OF trigger. TRIGGER_COLUMN_USAGE view This view would consist of the same 8 columns as in the base table, but would return only columns owned by the CURRENT_USER that are "referenced" in some trigger (either owned by the CURRENT_USER that are "referenced" in some trigger (either owned by the CURRENT_USER th</trigger></trigger></triggered></trigger></trigger>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					UPDATE Trigger columns of a trigger (possibly owned by some other user) defined in this catalog. This information is not derivable from the existing TRIGGERED_COLUMNS view because that view only returns triggers owned by the CURRENT_USER. The TRIGGERED_COLUMNS view (redefined over the new TRIGGER_COLUMN_USAGE base table) and the new TRIGGER_COLUMN_USAGE view could be used separately to answer all of a users legitimate Trigger questions. The TRIGGERED_COLUMNS view would return the UPDATE columns of triggers owned by the CURRENT_USER and the TRIGGER_COLUMN_USAGE view would return all catalog triggers that explicitly or implicitly "reference" a column owned by the CURRENT_USER. The first view would return the names of columns owned by other people that the given user had UPDATE privileges on, but never the names of triggers owned by other people, and the second view would return the names of triggers owned by other people but never the names of columns owned by other people. Both views are valuable to the user and contain information that a user has legitimate reason to know. Solution None provided with comment.	
	NLD-P02-124		2-Minor Technical	P02-No specific location	FND-521 The following Language Opportunity has been noted: Source: DBL:MCI-098/X3H2-96-111 Language Opportunity: The trigger descriptor defined in GR 2 of Subclause 11.39, " <trigger definition>", maintains an explicit collection of all column names referenced by the <triggered action=""> of the <trigger definition="">. This makes the trigger descriptor different in style from a table constraint descriptor (see Subclause 11.6, "", GR2) or a view descriptor (see Subclause 11.22, "<view definition="">", GR1), which only maintain this information implicitly. A table check consraint maintains the entire <search condition=""> of the Check and a view descriptor maintains the entire <query expression=""> that determines the view. It may be desirable to treat constraint, view, and trigger descriptors in a more homogeneous fashion. Alternatively, a trigger descriptor may just maintain the <triggered action=""> as part of the descriptor, rather than the "triggered action column set". If this is done instead, then Syntax Rule 5 and General Rule 1 of Subclause 11.18, "<drop column="" definition="">", would have to be re-written to accommodate <triggered action=""> instead of "triggered action column set". Solution None provided with comment.</triggered></drop></triggered></query></search></view></trigger></triggered></trigger 	
	NLD-P02-125		2-Minor Technical	P02-No specific location	FND-587 The following Language Opportunity has been noted: Source: Hugh Darwen, 27 January, 1997 Language Opportunity:	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
			č		Currently, all <routine invocation="">s that return values are deemed to be able to</routine>	
					return a null. Hence, such results are automatically tagged as "possibly null".	
					Wouldn't it be nice if you could say, when you define a function, "NEVER	
					RETURNS NULL" or words to that effect? Then its invocations would have the	
					nice "not nullable" characteristic.	
					Solution	
					None provided with comment.	
	NLD-P02-126		2-Minor	P02-No specific	FND-593 The following Language Opportunity has been noted:	
			Technical	location	Source: DBL:LGW-063/X3H2-97-077, point 46.	
					Language Opportunity:	
					There are no provisions in SQL3 for packaging ADT families. This type of	
					packaging is needed to support the creation of a package of ADTs and	
					associated subtypes and routines. It would be useful to define access control at	
					the package level rather than the individual ADTs or routines. It would also be	
					useful to be able to isolate the package so that subject routine resolution of	
					routines inside the package can be restricted to only other routines within the	
					package.	
					This packaging could be accomplished with schemas or SQL-server modules,	
					but neither mechanism is complete at this point.	
					Solution	
	NU D D02 127		2.16	DO2 Manual Ca	None provided with comment.	
	NLD-P02-127		2-Minor	P02-No specific	FND-613 The following Language Opportunity has been noted:	
			Technical	iocuion	Source: DBL:LGW-140/X3H2-97-349	
					The concernt of substitute bility is corntal to the ADT extension of SOL:	
					currently partiant information is scattered over a multitude of subclauses. It	
					needs to be summarized in a separate subclause of the Concepts section	
					Solution	
					None provided with comment	
	NI D P02 128		2 Minor	P02-No specific	FND 624 The following Language Opportunity has been noted:	
	NLD-102-120		Z-Ivinior Technical	location	Source: DRI I GW_1/6/X3H2_97_3/9	
			reennear		Language Onnortunity	
					Viewed tables allow the owner of a table to define a subset of its rows and/or	
					columns. The owner may then grant access to the viewed table to other users	
					without giving access to the base table itself. There is no corresponding	
					capability provided with reeference types. To access a column of a row for	
					which a user has a reference, the user is required to have SELECT privilege on	
					the column of the base table. To alter such a column, the user must have	
					UPDATE privilege on the column of the base table.	
					A mechanism analogous to views on base tables is extremely desirable for	
					adequate granualrity of access control.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					Solution	
					None provided with comment.	
	NLD-P02-129		2-Minor	P02-No specific	FND-626 The following Language Opportunity has been noted:	
			Technical	location	Source: DBL:LGW-146/X3H2-97-349	
					Language Opportunity:	
					The <dereference operation=""> is a very nice syntactic shorthand to avoid the</dereference>	
					writing of a join. This operation should be extended to allow the use of existing	
					referential constraints.	
					CREATE TABLE enrollments (
					student_Iname CHAR VARYING (30),	
					course REFERENCES courses (id)	
					grade CHAR VARYING (2),	
					FOREIGN KEY (student_lname, student_fname) REFERENCES	
					students (lname, fname)	
					(student lname student fname) -> address	
					FROM enrollments	
					WHERE grade = 'A+' ;	
					Solution	
					None provided with comment.	
	NLD-P02-130		2-Minor	P02-No specific	FND-627 The following Language Opportunity has been noted:	
			Technical	location	Source: DBL:LGW-146/X3H2-97-349	
					Language Opportunity:	
					A reference type should be able to refer to a cell of a table and not just the entire	
					row.	
					Solution	
					None provided with comment.	
	NLD-P02-131		2-Minor	P02-No specific	FND-629 The following Language Opportunity has been noted:	
			Technical	location	Source: DBL:LGW-080/X3H2-97-???	
					Language Opportunity:	
					The SQL3 specifications for <attribute definition="">, <routine specification="">, and</routine></attribute>	
					<abstract body="" data="" type=""> prohibit the ability to define an explicit mutator</abstract>	
					function on a single attribute of an ADT with the same signature as the implicit	
					This spectrice in <aurioute definition=""> (incredy over-right names both for the</aurioute>	
					attributes of an ADT and for its associated mutator functions. For example, with	
					the comment attribute of the SL Stillmage ADT it is not possible to define both	
					an attribute name and an explicit mutator function on that attribute with the	
					same name e g COMMENT cannot beused for both names	
					It is an SOL3 Language Opportunity to provide new syntax in the SOL3	
					<pre><attribute definition=""> to allow the implicit mutator function to be explicitly</attribute></pre>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
					renamed (e.g. similar to the way the CONSTRUCTOR option allows the implicit constructor function of an ADT to be renamed) so that the more desirable attribute name can then be used to define an explicit mutator function for that attribute. Example Usage: <attribute name=""> <data type=""> [MUTATOR <mutator name="">]. This new syntax might then be used to allow definition of a comment attribute in the SI_StillImage ADT, with its implicit mutator function renamed to be commentOnly, thereby allowing COMMENT to be used as the name of an</mutator></data></attribute>	
					explicit mutator function that modifies both the comment and the updateTime	
					attributes of the ADT.	
					Solution None provided with comment	
	NLD-P02-132		2-Minor Technical	P02-No specific location	None provided with comment. FND-630 The following Language Opportunity has been noted: Source: DBL:LGW-081/X3H2-97-??? Language Opportunity: Would it be possible to allow very limited Type Templates in SQL3 like DECLARE TYPE TEMPLATE Pixel (n SMALLINT) AS BIT (n) where an upper limit on the value of n is implementation-defined, but with the ability to specify the value of n as an integer <value expression=""> whenever Pixel(n) is declared as a parameter in an SQL-invoked routine or as an SQL variable in a compound statement.</value>	
	NI D D02 122		2 Miner	DO2 No specific	None provided with comment.	
	NLD-P02-133		2-Minor Technical	P02-No specific location	FND-6/6 The following Language Opportunity has been noted: Source: DBL:LGW-152/X3H2-97-352 (also DBL:LGW-023/X3H2-97-044, SEQ# 469, FRANCE-F-015*) Language Opportunity: Some types can be named by themselves (distinct types ADTs and named row types) while others only by defining domains on them (collections row types). This unorthogonality should be removed by allowing any type to be associated to a name through type declaration. Solution None provided with comment.	
	NLD-P02-134		2-Minor Technical	P02-No specific location P02-No specific	FND-696 The following Language Opportunity has been noted: Source: DBL:BBN-128/X3H2-98-354 (BBN-029R1, SEQ#149, USA-P02-034) Language Opportunity: The restriction that only rows of persistent base tables can be referenced should be lifted to allow references to nested (un-named) row types. Solution None provided with comment. FND-707 The following Language Opportunity has been noted:	
	100			location		

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
			Technical		Source: Email from Mike Ubell 5 August, 1998	
					Language Opportunity:	
					In X3H2-98-016, the ability to dynamically dispatch a function was eliminated	
					in favor of method based dispatch. This was done to bring SQL more in line	
					with Java and therefore, presumably, make it easier to import non-SQL written	
					shrink wrap applications into the database. Unfortunately many existing	
					applications (and data type packages) are not written in Java today, or even in	
					C++. By removing the multi-argument dispatch data types that support	
					comparison and inheritance must dispatch on one argument within the method.	
					If the method is implemented in a language that does not support inheritance,	
					then new subtypes may not be added to the shrink-wrapped data type.	
					Solution	
	NH D D02 126		2.16	D02 N :C	None provided with comment.	
	NLD-P02-136		2-Minor	P02-No specific	FND-719 The following Language Opportunity has been noted:	
			Technical	iocuiion	Source: WG3: YGJ-021	
					The reference type and the dereference operator have been added to SOL 3. The	
					ability to update a column or delete a row via a reference must be supplied as	
					well	
					Solution	
					None provided with comment	
	NLD-P02-137		2-Minor	P02-No specific	FND-720 The following Language Opportunity has been noted:	
	1,222 1 02 107		Technical	location	Source: WG3:YGJ-021	
					Language Opportunity:	
					SQL3 requires that a table have an associated user-defined type in order to be	
					referenceable. The combination of user-defined type and base table is now very	
					difficult to change in any way. The two would have to be disassociated, each	
					altered separately, and then associated again. Neither the disassociation of user-	
					defined type and base table nor the altering of a user-defined type are supported.	
					Solution	
					None provided with comment.	
	NLD-P02-138		2-Minor	P02-No specific	FND-721 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:YGJ-021	
					Language Opportunity:	
					Constraints are not a part of a user-defined type. This means that the constraints	
					that are intended for each table of a user-defined type must be explicitly copied	
					Solution	
	NI D D02 120		2 16	DO2 No mooif :	INORE provided with comment.	
	NLD-P02-139		2-iviinor	location	FIND-722 The following Language Opportunity has been noted:	1
1		1	recinical		JULICE. W UJ.I'IA-U72IA	1

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					Language Opportunity: The table defining features in Core SQL should be examined to ensure that the features exhaust all ov Core (perhaps by showing that all BNF nonterminals that are available to Core have been assigned to some faeture) and that they are rigorously stated. Solution	
	NLD-P02-140		2-Minor Technical	P02-No specific location	FND-758 The following Language Opportunity has been noted: Source: WG3:BHX-149 Language Opportunity: If might be useful to add to SQL the ability to use explicit character set names taken from the public registry for character set names (an IANA [Internet Assigned Numbers Authority] registry available at ftp://ftp.isi.edu.in-notes/iana/assignments/character- sets). Solution None provided with comment.	
	NLD-P02-141		2-Minor Technical	P02-No specific location	FND-773 The following Language Opportunity has been noted: Source: WG3:BHX-107/H2-2000 Language Opportunity: It is desirable to provide the capability on CREATE TABLE to change options (scope, reference checking, NOT NULL specification, default values, datalink control definitions, and so on) that are associated with components nested inside row types, collection types, and structured types. Solution None provided with comment.	
	NLD-P02-142		2-Minor Technical	P02-No specific location	FND-778 The following Language Opportunity has been noted: Source: WG3:BHX-117/H2-2000 Language Opportunity: WG3:SLD-046 added several new fields to the CLI descriptor area: CURRENT_TRANSFORM_GROUP, SPECIFIC_TYPE_CATALOG, SPECIFIC_TYPE_SCHEMA, and SPECIFIC_TYPE_NAME. The same fields could profitably be added to the SQL descriptor area, too. Solution None provided with comment.	
	NLD-P02-143		2-Minor Technical	P02-No specific location	 FND-780 The following Language Opportunity has been noted: Source: WG3:HEL-047/H2-2000 Language Opportunity: 2. Insurmountable (?) problem for query generators The unfriendliness described in FND-779 causes a certain difficulty to general purpose applications, such as query generators, that appears to be insurmountable. Given two arbitrary character string expressions of character set 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					CS, there is no guaranteed way of having them compared under the default collation of CS without knowing what that collation is. Moreover, the default collation can be looked up in the Information Schema only if the character set CS itself is known. There is no sure way that we are aware of whereby the character set of an arbitrary string expression can be determined by an SQL application.	
					None provided with comment	
	NLD-P02-144		2-Minor Technical	P02-No specific location	FND-787 The following Language Opportunity has been noted: Source: WG3:PER-146/H2-2001-??? (FCD1/2000 WG3-P01-011) Language Opportunity: [Jake Knoppers] saw that with respect to "normative references" point 1p that ISO 8601:2001 version is to be referenced. This is good; [he works] on that standard. [His]comment is that serious consideration should also be given to referencing ISO 19108:2000 "Geographic information Temporal schema". ISO 8601 deals mainly with Gregorian calendar referencing. Increasingly, various areas of business application such as banking/financial services, geomatics, intelligent transportation systems, etc. use other calendar referencing systems, such as the GPS clock, which is use for synchronization among the global position satellites and provides for a "common" single world wide date/time referencing among IT systems of autonomous organizations (one then maps the GPS date/time stamp to one's local time, whatever it is). It is likely that many SQL based implementations will do the same. [He does] not know whether you want to treat this as a "comment" an "informative note/footnote", etc. but [he thinks] that it is important for SQL users. Solution	
	NLD-P02-145		2-Minor Technical	P02-No specific location	None provided with comment.FND-788 The following Language Opportunity has been noted:Source: WG3:PER-146/H2-2001-??? (FCD1/2000 WG3-P01-018)Language Opportunity:Allow implementations to be able to represent year numbers outside of 0001-9999 (0000 is 1 B.C, etc.).The restriction of YEAR to be between 0001 and 9999 is unsupportable. Notealso that ISO/IEC 8601:2001 does not have any such restriction; 0000 andnegative years are allowed (year 0000 is year 1 BC, -0001 is year 2 BC,), asare year indications with more than 4 digits.Further, sub-second precision should be possible to use (i.e. required by thestandard). (Note: The CD Editing Meeting believes that this sentence means thatimplementations should be mandated to supply significant digits, other thanzero, to the right of the decimal point, although there may be hardware that doesnot support "clock ticks" at such a fine granularity.)Solution	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					None provided with comment.	
	NLD-P02-146		2-Minor	P02-No specific	FND-789 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:PER-146/H2-2001-??? (FCD1/2000 WG3-P02-010)	
					Language Opportunity:	
					Allow decimal numbers to be expressed using any one (for each numeral) of the	
					decimal number category (Nd) ranges in the UCS. Conversely, there should also	
					be a way of getting out formatted numbers using a specified range (by script	
					name or similar) of Nd characters.	
					Allow the character MINUS as an allas to HYPHEN-MINUS in arithmetic	
					expressions. Anow LESS THAN OR EQUAL, GREATER-THAN OR EQUAL,	
					GREATER-THAN OR SLANTED EQUAL (Unicode 3.2), and	
					comparison semantics Allow DOT OPERATOR for multiplication	
					Solution	
					None provided with comment	
	NLD-P02-147		2-Minor	P02-No specific	FND-791 The following Language Opportunity has been noted:	
	1122 102 117		Technical	location	Source: WG3:PER-146/H2-2001-??? (FCD1/2000 USA-P02-010)	
					Language Opportunity:	
					There is no discussion of the relationship between determinism and isolation	
					level. Two read transactions starting at the exact same time working on the	
					"same" SQL data can still have different results if they operate on different	
					isolation levels.	
					The May, 2001 CD Editing Meeting in Perth observed that describing such	
					interactions is extremely difficult and all such descriptions known to the Editing	
					Meeting participants rely heavily (perhaps exclusively) on the locking paradigm,	
					which the standard does not require. Because of this, the Editing Meeting	
					unlikely	
					Solution	
					None provided with comment	
	NLD-P02-148		2-Minor	P02-No specific	FND-807 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:PER-171/H2-2001-??? (FCD1/2000 USA-P02-010)	
					Language Opportunity:	
					It may be useful to have a notion of "hereditary property" of BNF nonterminals.	
					A hereditary property P would work like this: If $A ::= B$, then $P(A) = P(B)$,	
					unless there is an explicit syntax rule to the contrary.	
					Examples of hereditary properties would be declared type, scale, precision, most	
					specific type, value.	
					This is already the haphazard approach of the standard, for example, to say in	
					one SK that "the data type of B is D1" and then later assume that the data type of A is DT since A $\mu_{\rm e}$ B	
					Of A is D1 since A ::= B. $C_{abc} d^{a} = d^{a}$	
					Solution	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
					None provided with comment.	
	NLD-P02-149		2-Minor	P02-No specific	FND-827 The following Language Opportunity has been noted:	
			Technical	location	Source: FCD1/2000 WG4-P02-001	
					Language Opportunity:	
					It should be allowed to invoke a method using a <routine invocation=""> with a</routine>	
					signature that is identical to the <method selection=""> specified in Subclause 6.16,</method>	
					" <method invocation="">", and in Subclause 6.17, "<static invocation="" method="">",</static></method>	
					respectively.	
					Solution	
					None provided with comment.	
	NLD-P02-150		2-Minor	P02-No specific	FND-830 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:PER-188/H2-2001-???	
					Language Opportunity:	
					In the mathematical community, multiset union of M1 and M2 is defined as	
					consisting of every element that is an element of either <i>M1</i> or of <i>M2</i> , occurring	
					either as many times as it does in $M1$ or as many times as it does in $M2$,	
					whichever is the greater. (The SQL operator called UNION ALL, and also	
					called MULTISET UNION after acceptance of WG3:PER-098 is referred to as	
					"union plus", denoted thus: U+.) The mathematical definition of multiset union	
					seems just as good a counterpart of the multiset intersection we already have as	
					union plus does, because intersection can be expressed by just changing "either"	
					to "both", "or" to "and", and "greater" to "lesser" in the above informal	
					definition of multiset union.	
					Solution	
					None provided with comment.	
	NLD-P02-151		2-Minor	P02-No specific	FND-831 The following Language Opportunity has been noted:	
			Technical	iocation	Source: The merger of X3H2-95-178/DBL:YOW-048, X3H2-95-	
					201/DBL: YOW-049R, and X3H2-95-1/9R2/DBL: YOW-050R (Was Language	
					Opportunity PSM-061)	
					Exampliance Opportunity:	
					traceable. The list of <i>z</i> routine invocations that they were propagated back	
					through should be made available somewhere, such as in the Diagnostics Area	
					Solution	
					Solution None provided with comment	
	NI D P02 152		2 Minor	P02-No specific	FND 848 The following Language Opportunity has been noted:	
	11LD-102-132		Technical	location	Source: WG3:YYL016 (USA_P02_113)	
			reennear		Language Onnortunity	
					A number of DBMS products support materialized views whose results are	
					stored in the database and subsequently maintained by the system whenever any	
					of the generally underlying base tables of the views changes. Materialized views	
					play an important role in offering significant performance gains for complex	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					queries, especially in Data Warehouse applications.	
					The next edition of the SQL standard should standardize the syntax and	
					semantics of materialized views.	
					Solution	
					None provided with comment.	
	NLD-P02-153		2-Minor	P02-No specific	FND-849 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:YYJ-016 (USA-P02-114)	
					Language Opportunity:	
					In [FoundationCD], it is possible write insert and update statements where the	
					value of one or more fields are not immediately known by the updater. This	
					includes columns populated by subqueries, functions, system values, etc. In	
					some cases, the updater needs to know the values after the insert/update has	
					occurred. In some cases, this can be accomplished by requerying the data after	
					the update. In other cases, the updater cannot easily requery the data. This	
					includes cases such as when a function is used to generate the primary key. For	
					example: Insert into T1 (c1, c2, c3) values (fn_generate_pk('T1'), :var 2,	
					:var 3); It would be useful to have a mechanism that allows an insert or update	
					statement to return the inserted or updated rows in a singleton select or a cursor.	
					Solution	
					None provided with comment.	
	NLD-P02-154		2-Minor	P02-No specific	FND-850 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:YYJ-016 (USA-P02-117)	
					Language Opportunity:	
					SQL should be enhanced to support EJB Query Language.	
					Information about the EJB Query Language can be found the public document	
					available at:	
					http://java.sun.com/aboutJava/communityprocess/first/jsr019/ejb2-finaldraft.pdf	
					particularly in Chapter 10.	
					Solution	
					None provided with comment.	
	NLD-P02-155		2-Minor	P02-No specific	FND-876 The following Language Opportunity has been noted:	
	_		Technical	location	Source: WG3:DRS-128	
					Language Opportunity:	
					SQL/Foundation, as currently written, prohibits the creation and invocation of	
					multiple polymorphic routines whose parameters differ only by character set or	
					by interval class (year-month or day-time). This is clearly unacceptable for	
					many users' needs.	
					This Opportunity has been "narrowed" by acceptance of WG3:FRA-120R1. It	
					was formerly PSM-127.	
					Solution	
					None provided with comment.	

SEQ	Cmnt	See	a	De	Description	Addressed
#		Also	Severity	Reference		Ву
	NLD-P02-156		2-Minor	P02-No specific	FND-913 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:ZSH-155 = H2-2002	
					Language Opportunity:	
					Only SOL Elagor must detect	
					Solution	
					Solution None mention with comment	
	NI D D02 157		2 Minor	PO2 No specific	FND 014 The following Language Opportunity has been noted:	
	NLD-P02-137		2-Minor Technical	location	FND-914 The following Language Opportunity has been noted:	
			Technical	location	Language Opportunity:	
					Suppose you have defined a structured UDT with 50 attributes. In order to grant	
					somebody else the right to retrieve and set the values of each of those attributes	
					you must execute no fewer than 101 GRANT statements! First, you must grant	
					USAGE on the type itself. Then, you must grant EXECUTE on each of the 50	
					observer methods and EXECUTE on each of the 50 mutator methods. The	
					process is particularly cumbersome, because granting EXECUTE on the	
					observer methods requires something like "GRANT EXECUTE ON	
					INSTANCE METHOD attribute_n FOR typename TO username" (which is	
					easy enough), but granting EXECUTE on the mutator methods requires	
					something like "GRANT EXECUTE ON INSTANCE METHOD attribute_n	
					(argument-type-1, argument-type-2,argument-type-n) FOR typename TO	
					username". Of course, you could choose to use the <specific name=""> for the</specific>	
					methods, but those names are likely to be awkward and/or non-intuitive.	
					The process of entering all of those GRANTs is incredibly unfriendly to type	
					definers and grows worse as UDTs get more complex.	
					Contrast this with the process of granting retrieval and modification privileges	
					on a table with 1000 columns: "GRANT SELECT ON tablename TO username"	
					and "GRANT UPDATE ON tablename TO username". That's it.	
					Granting (and revoking!) access privileges to attributes of UD1s should be made	
					more user-iriendiy.	
					Solution	
	NI D D02 159		2 Minor	PO2 No masifia	INORE PROVIDED WITH COMMENT.	
	NLD-P02-138		2-Minor Technical	location	FND-913 The following Language Opportunity has been noted:	
			recinical	1.5 Curron	Janguage Onnortunity:	
					Instead of trying to discover and remember all the possible dependencies	
					between schema objects what we should do is create the dependences at the time	
					of creating the dependent object. This should do is create the dependency at the time	
					for DROP and REVOKE, as well as making them more intelligible and easier to	
					maintain.	
					Solution	
					None provided with comment.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
	NLD-P02-159		2-Minor Technical	P02-No specific location	FND-916 The following Language Opportunity has been noted: Source: WG3:ZSH-155 = H2-2002 Language Opportunity: The character string "associate" occurs 373 times in [FoundFCD], mostly in the phrase "associated with". In many cases the meaning, or effect, of an association between two objects can be found only by finding all the places where it is mentioned. In a number of such cases the phrase could be avoided altogether, in others the significance of the association could be more explicitly explained. We give one or two examples where it does not appear difficult to avoid the phrase. Subclause 03.03.01.01, "Other terms", <sql statement="" variable=""> that was associated with an <sql statement<br="">name> by a <prepare statement=""> Subclause 04.02.01, "Character strings and collating sequences", Each collation known in an SQL-environment is applicable to one or more character sets, and for each character set, one or more collations are applicable to it, one of which is associated with it as its character set collation. The words in bold are unnecessary, and could well be deleted altogether. The word "default" could be added, between "its" and "character set". Subclause 05.04, "Names and identifiers", Syntax Rule 17) 17) An <identifier> that is a <correlation name=""> is associated with a table within a particular scope. The scope of a <correlation name=""> is either a <select statement: single row>, <subquery>, or <query specification=""> (see Subclause 7.6, ""), or is a <trigger definition=""> (see Subclause 11.39, "<trigger definition="">"). Scopes may be nested. In different tables or with the same table. Solution None provided with comment</trigger></trigger></query></subquery></select </correlation></correlation></identifier></prepare></sql></sql>	
	NLD-P02-160		2-Minor Technical	P02-No specific location	FND-917 The following Language Opportunity has been noted: Source: WG3:ZSH-153R1 = H2-2002-153R1 Language Opportunity: The concepts section needs to explain that CAST AS is the mechanism to translate datetime and interval data types to and from host data parameters. Solution None provided with comment.	
	NLD-P02-161		1-Major Technical	P02-No specific location	FND-974 The following Possible Problem has been noted: Source: WG3:SIA-031 = H2-2004-??? Possible Problem: The General Rules applying to <rollback statement=""> are incomplete, and inconsistent with the text of Subclause 4.35.2, "Savepoints". General Rule 2) is, in part:</rollback>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					 1) If a <savepoint clause=""> is not specified, then:</savepoint> a) Every valid locator is marked invalid. b) All open cursors in any SQL-client module associated with the current SQLtransaction are closed. Nothing is said about locators or cursors held from the previous transaction. Nor is anything said about prepared statements. General Rule 3) is, in part: 1) If a <savepoint clause=""> is specified, then:</savepoint> a) Every valid locator that was generated in the current SQL-transaction subsequent to the establishment of S is marked invalid. b) For every open cursor CR in any SQL-client module associated with the current SQL-transaction that was opened subsequent to the establishment of S, the following statement is implicitly executed: CLOSE <i>CR</i> c) The status of any open cursors in any SQL-client module associated with the current SQL-transaction that were opened by the current SQL-transaction before the establishment of S is implementation-defined. NOTE 497 — The current SQL-transaction is not terminated, and there is no other effect on the SQL-data or schemas. This General Rule says nothing corresponding to the content of the final paragraph of Subclause 4.35.2, "Savepoints", nor about restoring some settable elements of the SQL-session context, e.g. current role name. Solution: Specify what happens in terms of the contents of the SQL-session context. Solution 	
	NLD-P02-162		2-Minor Technical	P02-No specific location	 FND-980 The following Language Opportunity has been noted: Source: WG3:STX-020 Language Opportunity: A <set role="" statement=""> always raises an exception if there is no current user identifier. This prevents the use of definer's rights routines (where the definer is a role) to handle the setting of roles. If this is desired functionality, then one of the following alternatives should be chosen and implemented.</set> a) Allow a <set role="" statement=""> if there is a current user identifier and the role is an applicable role for that user identifier or, if there is no current user identifier.</set> b) Allow a <set role="" statement=""> if the role is an applicable role for the current user identifier.</set> c) Allow a <set role="" statement=""> if the role is an applicable role for the most rhecent current user identifier. That is the user identifier with the highest stack</set> 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					position.	
					Solution	
					None provided with comment.	
	NLD-P02-163		1-Major	P02-No specific	NLD-P02-052	
			Technical	location	When the General Rules of another Subclause are invoked the specification of	
					the paramater passing is not always correct.	
					sometimes the arguments are not correctly identified. In the latter case is is	
					sometimes because the invoked Subclause does not itself given clearly	
					identifiable names to its arguments.	
					All the calling and called Subclauses should be checked and corrected.	
					See also FND-948	
					Solution	
					None provided with comment.	
	NLD-P02-164		3-Minor	P02-No specific	WG3-P02-006	
			Technical	location	A (so far unknown) possible problem is identified and resolved. Currently,	
					certain DDL statements do not have a restriction to disallow a <host parameter<="" td=""><td></td></host>	
					specification>, an <sql parameter="" reference="">, a <dynamic parameter<="" td=""><td></td></dynamic></sql>	
					specification>, or an <embedded specification="" variable=""> (and, a <sql td="" variable<=""><td></td></sql></embedded>	
					Solution	
					Solution None provided with comment	
		le l			SQL/CLI	
	NLD-P03-001		1-Major	P03-05.04, CLI	WG3-P03-001	
			Technical	Implicit Cursor	GR 8) b) "The General Rules of Subclause 14.1, " <declare cursor="">", in ISO/IEC</declare>	
					90/3-2 are applied to CK. Doesn't work because it doesn't say what CP is equivalent to in <declare< td=""><td></td></declare<>	
					Cursor> and in any case I don't think there is anything to be equivalent to - CR	
					is not defined in <declare cursor=""> only in <open cursor="">.</open></declare>	
					Solution	
					None provided with comment.	
	NLD-P03-002		1-Major	P03-05.06,	WG3-P03-002	
			Technical	Implicit	In GR 4) p) 2) 1) B) II) "If the <cast specification=""></cast>	
				EXECUTE	CAST (SV AS TDT)	
				OSING ana OPEN USING	does not conform to the General Rules of Subclause 6.12, " <cast< td=""><td></td></cast<>	
				clauses	specification>", in ISO/IEC 9075-2, then an exception condition is raised in	
					accordance with the General Rules of Subclause 6.12, " <cast specification="">", in</cast>	
					ISO/IEC 70/J-2.	
					I don't think you can conform to "General Rules"	
					Also why is this rule not covered by the following subrule (III) which	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					(effectively) invokes the syntax and General Rules of 6.12, " <cast< th=""><th></th></cast<>	
					specification>".	
					Solution	
	NH D D02 002		1.26.1	D02.06.10	None provided with comment.	
	NLD-P03-003		1-Major	P03-06.19, ExacDirect	WG3-P03-003	
			Technical	ExecDirect	I his subclause needs to be examined to see if it needs similar treatment to that	
					statements" A similar comment applies to Subclause 6.20 "Execute"	
					The drafting of this P.P. completes Action Item k) recorded in STX-012	
					"Action Items from Xi'an".	
					Solution	
					None provided with comment.	
	NLD-P03-004		4-Minor	P03-06.33	CLI-053 The following Possible Problem has been noted:	
			Editorial	GetDiagField	Source: WG3:HBA-042 = H2-2003-	
					Possible Problem:	
					Hanging between In General Rule 12) i), hanging between Subrules ii) C) and	
					iii), is the sentence:	
					If the value of TABLE_NAME identifies a declared local temporary table, then	
					the value of CATALOG_NAME is <space>s and the value of</space>	
					SCHEMA_NAME is 'MODULE'.	
					I ne second sentence of General Rule 12(1) (iii) (1) is:	
					CATALOC NAME is spaces and SCHEMA NAME contains 'MODULE'	
					It rather looks as though the former was intended to replace the latter since the	
					style of wording it uses seems to be more prevalent	
					SOL:1999 contains the same error.	
					Solution Replace the second sentence of General Rule 12) i) iii) 1) (as quoted	
					above) with:	
					If the value of TABLE_NAME identifies a declared local temporary table, then	
					the value of CATALOG_NAME is <space>s and the value of</space>	
					SCHEMA_NAME is 'MODULE'.	
					Solution	
					None provided with comment.	
	NLD-P03-005		I-Major	P03-No specific	NLD-P03-002	
			Technical	ioculion	when the General Rules of another Subclause are invoked the specification of	
					Sometimes the name of the argument(s) is (arg) not explicitly given and	
					sometimes the arguments are not correctly identified. In the latter case is is	
					sometimes because the invoked Subclause does not itself given clearly	
					identifiable names to its arguments.	
					All the calling and called Subclauses should be checked and corrected.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					Solution	
					None provided with comment.	
	NLD-P03-006		2-Minor	<i>P03-06.17</i> ,	CLI-055 The following Possible Problem has been noted:	
			Technical	EndTran	Source: WG3:STX-053 Addressing Action Item n) from Xian on SIA031	
					Possible Problem:	
					SIA-031 Para 2.1.1 made Changes to Working Draft SQL/Foundation Subclause	
					16./, " <commit statement="">"</commit>	
					Solution	
	NI D D02 007		2 Minor	P03 04 03	None provided with comment.	
	NLD-P05-007		2-Milloi Technical	Diagnostics	Source: WG3:STX 053 Addressing Action Item n) from Yian on SIA031	
			Teennear	areas in SQL/CLI	Possible Problem.	
				and P03-0A.2,	SIA-031 Para 2.2.3 made Changes to SOL/Foundation Annex B.	
				COBOL library	"Implementation-defined elements"	
				item SQLCLI	The diagnostics areas in CLI are different than those of embedded/module SQL.	
					But there does not appear to be any indication in CLI of whether it covers the	
					relationship of the CLI diagnostics area to the impact of a ROLLBACK to a	
					SAVEPOINT.	
					We may need to modify CLI to make the impact on the CLI diagnostics area	
					implementation-defined but before doing this we might want to check what CLI	
					implementations do today.	
					Solution	
	NI D D02 009		2 Minor	P03 05 04	None provided with comment.	
	NLD-P05-008		2-Minor Technical	Implicit cursor	Source: X3H2 98 077P1/DBL: BBN 222 and Paul Cotton March 1, 1998	
			Teennear		Language Onnortunity	
					General Rule 7)e) "Case" i) "If CR specifies INSENSITIVE" carries out the	
					same functionality as expressed in the General Rules of SQL/Foundation <open< td=""><td></td></open<>	
					cursor>. It is a Language Opportunity to reference the appropriate rules in	
					SQL/Foundation instead of repeating them here.	
					Solution	
					None provided with comment.	
	NLD-P03-009		2-Minor	P03-08.01,	CLI-026 The following Language Opportunity has been noted:	
			Technical	Claims of	Source: X3H2-98-077R1/DBL:BBN-??? and Source: Paul Cotton, March 1,	
				SOL/CLI	1998	
					Language Opportunity: Would it make sense to have a CLI flagger which discovers perpertable	
					extensions? One way to do this would be to set an environment attribute (if there	
					is such a thing) saying that any use of a nonportable argument will return a	
					special status code. CLI should support this requirement only if it is also	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					required for conformance to dynamic SQL.	
					Solution	
					None provided with comment.	
	NLD-P03-010		2-Minor	P03-No specific	CLI-047 The following Language Opportunity has been noted:	
			Technical	location	Source: FCD (1999) ballot comment USA-P03-024	
					Language Opportunity:	
					WG3:SLD-010/X3H2-98-027R3 provides for fetching multiple rows in one CLI	
					routine invocation. It would be appropriate to be able to provide an array of	
					input parameter values to a single statement execution in a similar fashion.	
					Solution	
					None provided with comment.	
	NLD-P03-011		2-Minor	P03-No specific	CLI-048 The following Language Opportunity has been noted:	
			Technical	location	Source: FCD (1999) ballot comment USA-P03-025	
					Language Opportunity: WC2.SLD 010/V2112 08 027D2 grouides for fataking multiple roug into an	
					wG5.SLD-010/ASH2-98-02/R5 provides for felching multiple rows into an	
					the application can bind to an array of record structures, where fields of the	
					record structure are the input or output parameters	
					Solution	
					None provided with comment	
	NLD-P03-012		2-Minor	P03-06.21. Fetch	CLI-049 The following Language Opportunity has been noted:	
	1122 100 012		Technical		Source: WG3:SLD-010/X3H2-99-027R3	
					Language Opportunity:	
					The arrays that receive the results from multi-row fetches must be contiguous.	
					For example, if you are performing	
					SELECT EMPNO, NAME FROM EMP	
					the application cannot create a record structure with fields for EMPNO and	
					NAME, and then create an array of these structures. The reason is that all the EMDNOs will be delivered in a single contiguous error, and all of the NAMEs	
					will be delivered in a separate array. It would be useful to provide for an offset	
					with a record structure or a "stride" (distance between successive elements of an	
					array) This is a method of binding known as row-wise binding	
					Row-wise binding was deliberately not part of the paper that proposed multi-	
					row fetch since it is an orthogonal enhancement and therefore benefits by being	
					considered in a separate proposal. We note in passing that row-wise binding can	
					be accomplished simply and elegantly by introducing a new descriptor field that	
					informs whether the buffers are laid out as 'regular' (or column-wise) binding, or	
					as row-wise binding.	
					Solution	
					None provided with comment.	
	NLD-P03-013		2-Minor	P03-06.34,	CLI-054 The following Language Opportunity has been noted:	
			Technical	GetDiagRec	Source: WG3:STX-001 Action Item n) Mark Ashcroft. Additional to SIA-	

SEQ	Cmnt	See				Addressed	
#	ID	Also	Severity	Reference	Description	Ву	
					025R1 Language Opportunity: SIA-025R1, "A Shorthand for Getting ALL Diagnostics" proposes to add a new diagnostics option to embedded SQL but ignores the question of whether it should also be added to CLI or the SQLJ binding. Solution None provided with comment.		
SQL/PSM							
	_NLD-P04-001		1-Major Technical	P04-05.02, Names and identifiers	PSM-153 The following Possible Problem has been noted: Source: WG3:HBA-042 = H2-2003 Possible Problem: SQL/Foundation, Subclause 5.4, "Names and identifiers", Syntax Rule 5) a) says: a) If LSQ [a <local or="" qualifier="" schema="">] is "MODULE", then TN [a <table name>] shall be contained in an <sql-client definition="" module=""> M and the <module contents=""> of M shall contain a <temporary declaration="" table=""> TT whose has a <qualified identifier=""> equivalent to QI. This does not cater for the case of a <temporary declaration="" table=""> referenced by a contained in a <module routine="">. Solution Perhaps PSM should replace the cited Syntax Rule 5) a) with: a) If LSQ [a <local or="" qualifier="" schema="">] is "MODULE", then TN [a <table name>] shall be contained either in an <sql-client definition="" module="">, without an intervening <sql-schema statement="">, or in a <sql-client module<br="">definition> that contains a <temporary declaration="" table=""> TT whose <table name> has a <qualifier> equivalent to QI</qualifier></table </temporary></sql-client></sql-schema></sql-client></table </local></module></temporary></qualified></temporary></module></sql-client></table </local>	See Comment	
	NLD-P04-002		1-Major Technical	P04-05.02, Names and identifiers	PSM-154 The following Possible Problem has been noted: Source: WG3:HBA-042 = H2-2003 Possible Problem: Whatever is said, in the Syntax or General Rules of SQL/Foundation, Subclause 5.4, "Names and identifiers", about how a identifies a created local temporary table must be replaced by PSM to cater for <module routine="">s. Let TT be a local created temporary table; let R1 be SQL-routine in SQL-server module M1 and let R2 be SQL-routine in SQL-server module M2. Both R1 and R2 contain references to TT. It is understood that, regardless of the pattern of invocation, each of R1 and R2 has its own instance of TT. None of the foregoing is specified in any General Rule, and is described inadequately Concepts. It needs to be properly specified. Solution None provided with comment.</module>		
	NLD-P04-003		1-Major	P04-08.01,	PSM-155 The following Possible Problem has been noted:		

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
			Technical	<routine< th=""><th>Source: WG3:HBA-042 = H2-2003</th><th></th></routine<>	Source: WG3:HBA-042 = H2-2003	
				invocation>	Possible Problem:	
					SQL/Foundation, Subclause 10.4, " <routine invocation="">", General Rule 5) d) i)</routine>	
					is:	
					1) If R is an SQL routine, then remove from RSC the identities of all instances of	
					created local temporary tables,	
					This doesn't look good for an SQL routine RT contained in an SQL-server	
					Solution:	
					PSM must modify the cited subrule in some way	
					Solution	
					None provided with comment	
	NLD-P04-004		1-Major	P04-08.01.	PSM-156 The following Possible Problem has been noted:	
			Technical	<routine< th=""><th>Source: WG3:HBA-042 = $H2-2003$-</th><th></th></routine<>	Source: WG3:HBA-042 = $H2-2003$ -	
				invocation>	Possible Problem:	
					It is not clear whether the rows of a temporary table (whether declared or	
					created) that is local to an SQLserver module survive from one invocation of	
					SQL-routines in that module.	
					Let M be a SQL-server module and TT either a declared temporary table local to	
					it, or a local created temporary table; let R be an SQL-routine in M that accesses	
					TT. R is invoked twice during a transaction from by some invoker INV.	
					It seems to be intended that, provided INV does not, between the invocations of	
					R1, access TT, the second invocation of R will find TT as the first invocation left it.	
					Furthermore, if INV is an SQL-routine in M, and INV accesses TT, then INV	
					and R access the same (instance of) TT.	
					On the other hand, if INV is in an SQL-server module MI, different from M,	
					then, whether TT is a declared temporary table or a local created temporary	
					table, a reference to TT in INV, if valid, identifies a different table from the one	
					accessed by R.	
					I nen again, if K invokes a further SQL-routine KS, that, like livy, is in MI, then	
					The foregoing is not specified in any General Rule, nor is it clearly described in	
					Concents	
					Solution	
					None provided with comment.	
	NLD-P04-005		1-Major	P04-10.03,	PSM-149 The following Possible Problem has been noted:	
	10.000		Technical	<revoke< th=""><th>Source: CD1-2000 comments USA-P04-005</th><th></th></revoke<>	Source: CD1-2000 comments USA-P04-005	
				statement>	Possible Problem:	
					Because PSM expands the possibilities of <sql procedure="" statement="">, the</sql>	
					capabilities for the <triggered action=""> of a trigger are much increased.</triggered>	
					Consequently the rules regarding dependencies of a trigger on a privilege or	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					schema object must be extended in PSM. For example, in 9075-2	
					(SQL/Foundation), Subclause 12.7, " <revoke statement="">", SR 24) subrules g)</revoke>	
					through j) deal with when SELECT privilege is required to define a trigger.	
					None of these rules cover the possibility of a <scalar subquery=""> in a <case< th=""><th></th></case<></scalar>	
					statement>. Likewise the rules for SELECT WITH HIERARCHY OPTION are	
					inadequate.	
					The commenter does not believe that the solution is to run around trying to find	
					need a general mechanism that constructs a dependency graph relating arbitrary	
					schema objects and privileges, so that as features and parts are added, each new	
					feature or part need only specify its contribution to the dependency graph	
					algorithm. For example, dependencies on privileges can be declared in the	
					Access Rules, so that whenever an Access Rule is used, a dependency is	
					automatically created. That way <revoke statement=""> would not need to duplicate</revoke>	
					information that is already implicit in the Access Rules. Similarly, dependencies	
					on schema objects can be generated in the rules of , <column< th=""><th></th></column<>	
					reference>, etc. Then <revoke statement=""> and the drop statements would not</revoke>	
					need to generate dependencies, they could simply assume that they are defined.	
					Solution	
					None provided with comment.	
	NLD-P04-006		1-Major	P04-11.02, <sql< th=""><th>PSM-158 The following Possible Problem has been noted:</th><th></th></sql<>	PSM-158 The following Possible Problem has been noted:	
			Technical	proceaure statement>	Source: DCOR/2004 WG3-P04-001	
				sidiemeni>		
					General Rule 3) calls for the General Rules of, " <nandler declaration="">" to be</nandler>	
					and there does not seem to be any clear context available in either <sql< th=""><th></th></sql<>	
					procedure statements or in $<$ handler declarations to make an implicit choice	
					Solution	
					None provided with comment	
	NLD-P04-007		1-Major	P04-No specific	PSM-152 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:HBA-040	
					Language Opportunity:	
					The scope of an extended name that contains LOCAL is not adequately	
					specified in the following cases:	
					a) Where the extended name is contained in an <sql control="" statement=""></sql>	
					immediately contained in an <externally-invoked procedure="">.</externally-invoked>	
					b) Where the extended name is contained in a <schema routine="">.</schema>	
					c) Where the extended name is contained in a <module routine=""></module>	
					The determination of what object, if any, is identified by an extended name	
					should not depend on the statement that contains it being contained in a	
					particular <5QL-server module definition>, still less a particular <sql-client< th=""><th></th></sql-client<>	
					module definition>.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					For consistency, if a local scope is to be permitted in these cases, it should	
					follow the precedent of <cursor name="">, provided <statement name=""> also follows</statement></cursor>	
					it.	
					The preferred solution, however, is to make all extended names global, by	
					deleting <scope option="">.</scope>	
					Solution	
					None provided with comment.	
	NLD-P04-008		2-Minor	P04-No specific	PSM-088 The following Language Opportunity has been noted:	
			Technical	location	Source: Steve Cannan, during the course of discussing DBL:MCI-060	
					Language Opportunity:	
					Need some syntax to do an ALTER VIEW or similar to "rebind" subject	
					routines, * column references, etc. for all objects that contain statically-bound	
					references of any sort.	
					Solution	
					None provided with comment.	
	NLD-P04-009		2-Minor	P04-No specific	PSM-095 The following Language Opportunity has been noted:	
			Technical	location	Source: Ed Dee, while discussing DBL:MCI-132 ballot comments	
					Language Opportunity:	
					FOR statements terminate (with a closed cursor exception) if the statement list	
					of the <for statement=""> list contains a COMMIT or ROLLBACK. Further, no</for>	
					statement contained in the <for statement=""> can set any transaction attributes.</for>	
					It is desirable that an application programmer be able to initiate or terminate	
					transactions within a <10r statement>.	
					Solution None provided with comment.	
	NLD-P04-010		2-Minor	P04-08.01,	PSM-106 The following Language Opportunity has been noted:	
			Technical	<routine< td=""><td>Source: DBL:MCI-161, point 2.5, item 5</td><td></td></routine<>	Source: DBL:MCI-161, point 2.5, item 5	
				invocation>	Language Opportunity:	
					In Subclause 8.1, " <routine invocation="">", the prohibitions on SQL-transaction</routine>	
					staements and SQL-connection statements in SQL-invoked routines might be	
					lifted, if a way can be found to make sure that SQLinvoked routines end SQL-	
					sessions and SQL-transactions that they start, don't end SQL-transactions and	
					SQL-sessions that they didn't start, and don't switch SQL-connections without	
					restoring the SQL-connection with which they started.	
					Solution	
					None provided with comment.	
	NLD-P04-011		2-Minor	P04-08.01,	PSM-107 The following Language Opportunity has been noted:	
			Technical	<routine< td=""><td>Source: Discussion of DBL:MCI-161, point 2.5, item 5</td><td></td></routine<>	Source: Discussion of DBL:MCI-161, point 2.5, item 5	
				invocation>	Language Opportunity:	
					In Subclause 8.1, " <routine invocation="">", the prohibitions on SQL-transaction</routine>	
					statements and SQL-connection statements in SQL-invoked routines might be	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
					lifted by changing "SQL-connection statement" to "SQL-connection statement and the implementation does not support the execution of that SQL-statement in an invoked SQL-routine that is a procedure" in each of the two rules that make this prohibition, and making an appropriate entry in Annex B, "Implementation- defined elements", saying something like "It is implementation-defined whether or not an SQL-implementation supports the execution of SQL-transaction statements and/or SQL-connection statements in an invoked SQL-routine; if it does so, then the effects are implementation-defined." Solution None provided with comment.	
	NLD-P04-012		2-Minor Technical	P04-No specific location	PSM-124 The following Language Opportunity has been noted: Source: DBL:MCI-040/X3H2-96-169:UK-017 Language Opportunity: No way of obtaining the associated sqlstate of a condition name. We think the <condition name=""> feature is a nice idea, but we suspect it will generate a requirement, akin to the observation in the preceding comment, for a built-in function to return the associated sqlstate value of a given condition name. Furthermore, it might even be required to hold condition names in variables or arguments, in which case they have to become character strings. We would be happy to hold this feature over for SQL3, in the interests of simplification and early progression of PSM2 and to give time for the requirements to be fully thought through and appropriately addressed in the language. Solution None provided with comment.</condition>	
	NLD-P04-013		2-Minor Technical	P04-No specific location	PSM-140 The following Language Opportunity has been noted: Source: DBL:LGW-081/X3H2-97-??? Language Opportunity: Is it possible in SQL3 to relax the specification of string data types such as <character stringtype=""> and <bit string="" type=""> so that the declared length of these types (with appropriate usage restrictions) can be specified at execution time rather than at compile time? Can I declare avariable in an outer block of a compound statement and then use that variable as the <length> of a bit string variable declaration in an inner block? Solution None provided with comment.</length></bit></character>	
	NLD-P04-014		2-Minor Technical	P04-No specific location	 PSM-152 The following Language Opportunity has been noted: Source: WG3:HBA-040 Language Opportunity: The scope of an extended name that contains LOCAL is not adequately specified in the following cases: a) Where the extended name is contained in an <sql control="" statement=""> immediately contained in an <externally-invoked procedure="">.</externally-invoked></sql> 	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					 b) Where the extended name is contained in a <schema routine="">.</schema> c) Where the extended name is contained in a <module routine=""></module> The determination of what object, if any, is identified by an extended name should not depend on the statement that contains it being contained in a particular <sql-server definition="" module="">, still less a particular <sql-client definition="" module="">.</sql-client></sql-server> For consistency, if a local scope is to be permitted in these cases, it should follow the precedent of <cursor name="">, provided <statement name=""> also follows it.</statement></cursor> The preferred solution, however, is to make all extended names global, by deleting <scope option="">.</scope> 	
	NI D D04 015		2 Minor	DOA 14 Dumannia	None provided with comment.	
	NLD-P04-015		2-Minor Technical	P04-14, Dynamic SQL	PSM-157 The following Language Opportunity has been noted: Source: WG3:HBA-048 = H2-2003 Language Opportunity: SQL/Foundation, Subclause 19.10, " <input clause="" using=""/> ", Syntax Rule 1) is: 1) The <general specification="" value=""> immediately contained in <using argument> shall be either a <host parameter="" specification=""> or an <embedded variable specification>. and SQL/Foundation, Subclause 19.11, "<output clause="" using="">", Syntax Rule 1) is: 1) The <target specification=""> immediately contained in <into argument=""> shall be either a <host parameter="" specification=""> or an <embedded variable<br="">specification>. Without these being modified by SQL/PSM, it is thus not currently possible for an SQL parameter to be either a <using argument=""> or an <into argument="">. Solution None provided with comment.</into></using></embedded></host></into></target></output></embedded </host></using </general>	
					SOL/MED	
	NLD-P09-001		2-Minor Technical	P09-24.10, ROUTINE_MAP PING_OPTIONS view	MED-067 The following Possible Problem has been noted: Source: DCOR/2004, WG3-P09-003 Possible Problem: The View ROUTINE_MAPPING_OPTIONS has no privilege check and no restriction to the actual catalog.	
					Solution	
	NLD-P09-002		2-Minor Technical	P09-24.11, ROUTINE_MAP PINGS view	None provided with comment. MED-068 The following Possible Problem has been noted: Source: DCOR/2004, WG3-P09-004 Possible Problem: The View ROUTINE_MAPPINGS has no privilege check and no restriction to	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					the actual catalog.	
					Solution	
					None provided with comment.	
	NLD-P09-003		2-Minor	P09-24.12,	MED-069 The following Possible Problem has been noted:	
			Technical	USER_MAPPIN	Source: DCOR/2004, WG3-P09-005	
				G_OPTIONS	Possible Problem:	
				view	The View USER_MAPPINGS_OPTIONS has no privilege check and no	
					restriction to the actual catalog.	
					Solution	
					None provided with comment.	
	NLD-P09-004		2-Minor	P09-24.13,	MED-070 The following Possible Problem has been noted:	
			Technical	USER_MAPPIN	Source: DCOR/2004, WG3-P09-006	
				GS view	Possible Problem:	
					The View USER_MAPPINGS has no privilege check and no restriction to the	
					actual catalog.	
					Solution	
					None provided with comment.	
	NLD-P09-005		2-Minor	P09-25.02,	MED-071 The following Possible Problem has been noted:	
			Technical	DATA_TYPE_DE	Source: DCOR/2004, WG3-P09-008	
				SCRIPTOR base	Possible Problem:	
				table	The Constraint	
					DATA_TYPE_DESCRIPTOR_DATA_TYPE_CHECK_COMBINATIONS of	
					the table DATA_TYPE_DESCRIPTOR is out of synch with its definition in Part	
					11 (Schemata).	
					Solution	
					None provided with comment.	
	NLD-P09-006		2-Minor	<i>P09–25.10</i> ,	MED-072 The following Possible Problem has been noted:	
			Technical	ROUTINE_MAP	Source: DCOR/2004, WG3-P09-009	
				PINGS base lable	Possible Problem:	
					The constraint ROUTINE_MAPPINGS_PRIMARY_KEY requires that the	
					value of the column ROUTINE_MAPPING_NAME is unique across all catalogs	
					in a given DEFINITION_SCHEMA. This seems not be reasonable.	
					Solution	
				D 00.06.0 2	None provided with comment.	
	NLD-P09-007		3-Major	P09-06.02, <cast< td=""><td>MED-065 The following Possible Problem has been noted:</td><td></td></cast<>	MED-065 The following Possible Problem has been noted:	
			Editorial	specification>	Source: FCD1/2002, USA-P09-041	
					Possible Problem:	
					The table in SK 2) is an inappropriate way to add new data types to the casting table in ISO/IEC 0075.2	
					table III 150/1EC 90/3-2.	
					A unretent approach would be preferable to avoid problems caused by adding data types in multiple incremental parts (e.g. DATALINK in SOL/MED and	
					data types in multiple incremental parts (e.g., DATALINK III SQL/MED and	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					XML in SQL/XML). A better approach would be to use a new SR 2) that says something like "Add a new rightmost column to the table following SR 6) in ISO/IEC 9075-2", followed by a table that looks something like this: <data type=""> SD of <data type=""> of TD <value expression=""> DL EN N AN N</value></data></data>	
					RW M Then another new SR would be specified, something like this: "Add a new row at the end of the table following SR 6) in ISO/IEC 9075-2" <data type=""> SD of <data type=""> of TD <value expression=""> EN AN VC FC D T TS YM DT BO UDT CL BL RT CT RW DL DL N N N N N N N N N N N N N N Y This approach has the advantage of correctly inserting a column and a row, rather than replacing the entire table. However, it leaves the disadvantage that insertion of a column and a row by SQL/MED and another by SQL/XML causes two cells of the table to be unspecifiedthe cell concerning casting of the data type added by SQL/MED to and from the data type added by SQL/XML. That disadvantage might be resolved by adding (e.g., in Foundation) a statement that such "unspecified cells" are implicitly filled with "N", so that no such casting is supported.</value></data></data>	
					Solution	
	NLD-P09-008		3-Major Editorial	P09-No specific location	MED-064 The following Possible Problem has been noted: Source: FCD1/2002, DEU-P09-980 Possible Problem: A look at Clause 4, "Concepts", and associated Subclauses seems to suggest that many columns defined in Clause 25, "Definition Schema", that are presently optional (meaning that a value of null is permitted) should be mandatory. Thus, a careful examination of all column definitions is required, and some of them may require NOT NULL constraints to be added. Solution None provided with comment.	
	NLD-P09-009		1-Major	P09-No specific	When the General Rules of another Subclause are invoked the specification of	
			Technical	location	the paramater passing is not always correct.	
					sometimes the arguments are not correctly identified. In the latter case is is	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
			, i i i i i i i i i i i i i i i i i i i		sometimes because the invoked Subclause does not itself given clearly	
					identifiable names to its arguments.	
					All the calling and called Subclauses should be checked and corrected.	
					Solution	
					None provided with comment.	
	NLD-P09-010		2-Minor	P09-No specific	MED-028 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:RTM-017R3/X3H2-99-255R2, Comment WG3-P09-005	
					Language Opportunity:	
					Acceptance of WG3:YGJ-082 made it prohibited to link a single external file	
					more than once. This has been identified as an undesirable restriction in at least	
					some situations.	
					Solution	
					None provided with comment.	
	NLD-P09-011		2-Minor	P09-No specific	MED-033 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:RTM-017R3/X3H2-99-255R2, Comment WG3-P09-011	
					Language Opportunity:	
					It is desirable to provide the capability to deal with character sets and collations	
					for character string columns of foreign tables.	
					Solution	
					None provided with comment.	
	NLD-P09-012		2-Minor	P09-No specific	MED-045 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:BHX-148/H2-2000	
					Language Opportunity:	
					WG3:BHX-148R1 proposed the use of only U1F-16 to communicate character	
					strings between the SQL server and the foreign-data wrapper. This limitation	
					implementation defined character sets	
					Solution	
					Solution None provided with comment	
	NI D 200 013		2 Minor	POQ No specific	MED 046 The following Language Opportunity has been noted:	
	NLD-F09-015		Z-Millor Technical	location	Source: WG3:BHX_108P1/H2_2000_ and ECD1 2000_GBR_P09_0/1	
			reennear		Language Onnortunity	
					Generic options — some requirements are obvious and should be standardized	
					— for example the name by which the FT is known at the FS may be different	
					from that in the SQL Environment. If the server is SQL-aware, then the foreign	
					table could be defined by a <query specification="">. There is a need for discussion</query>	
					of the costs/benefits/opportunities/mechanisms for further standardization.	
					Solution	
					None provided with comment.	
	NLD-P09-014		2-Minor	P09-No specific	MED-047 The following Language Opportunity has been noted:	
			Technical	location	Source: FCD1 2000, GBR-P09-043, FCD1 2000, GBR-P09-001, and FCD1	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					Language Opportunity: There is a need to acknowledge current implementations: Make FOREIGN DATA WRAPPER optional and add options USE INTERFACE <name> and USE PROTOCOL <name> Use of Standard Interfaces.</name></name>	
					Where standard interfaces already exist for accessing foreign data, it should be possible to reference the interfaces without requiring Wrappers. Example:	
					Let A and B be RDBMS Vendors; Let X and Y be video specialists. If AX is an implementation of Video using SQL MED and a wrapper WX designed by X and BY is an implementation of Video using SQL MED and a wrapper WY designed by Y then SQL MED does not guarantee that the WY wrapper will work with A or that WX will work with B or that a user of AX can easily port their application to BY.	
					Suppose both X and Y support a standard interface VAPI, then it would be possible to write wrappers that map to VAPI. This might achieve some ability to change video suppliers, but only if the wrapper writers use the VAPI interface with portability in mind. Actual interchangeability is most likely if the wrappers are written by the vendors A and B and supported by them. But in this case the	
					SQL-MED interface becomes an internal one of no interest to users. Use of Protocols. Where foreign data is remote and protocols exist for accessing the remote	
					servers, it should be possible to reference the protocols without requiring wrappers. Example:	
					Let A and B be RDBMS vendors; Let AP and BP be protocols used for accessing remote servers by A and B.	
					Most vendors have a proprietary protocol and many have also implemented their competitors' protocols too. Hence there is already a well defined means of accessing remote data.	
					If these protocols are implemented through wrappers then interchangeability of components could be achieved at three levels: — SOL-MED	
					- A protocol API - The protocol itself	
					Of these, the SQL-MED interface is the most complex, the latest to appear and the most incomplete. It seems to add no value.	
					We think it would be more appropriate to let the foreign server supporting the foreign tables be directly associated with the Protocol	
					Solution	
	NLD-P09-015		2-Minor	P09-No specific	MED-055 The following Language Opportunity has been noted:	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
#	ID	Also	Severity Technical	Reference location	DescriptionSource: FCD1 2000, AUS-P09-007Language Opportunity:There are a number of places in the 'Sequence of actions during foreign serverrequest executions' where the same routine may be called multiple times toreturn information about options etc. In addition there are some places whereMultiple routines are called each returning one item of information at a timefrom about the particular object.Each of these calls requires a 'context' switch in most operating systems whichmay in some circumstances end up incurring a substantial operating systemoverhead in terms of CPU etc.Thus it would be preferable if there were additional alternative methods bywhich this information could be passed between the SQL Server and the foreignwrapper routines.One mechanism may be to use a structure for various components that may bepassed directly to the wrapper routine. Alternatively more than one item ofinformation may be returned by a single callThus for example, in addition to the following routines— GetServerName— GetServerVersion~ dispervice for construction~ dispervice for construction	Ву
					a single routine GetServerInfo may return all the information. Or in the case where multiple calls would be made to a single routine (for example GetWrapperOption) to return multiple options either an array or a formatted XML document may be used so that a single call may return multiple options. We would like to see some discussion on the possibility of adding optimal multi-return-value procedures to reduce the possible overhead of excessive multiple procedural calls.	
					None provided with comment.	
	NLD-P09-016		2-Minor Technical	P09-No specific location	MED-056 The following Language Opportunity has been noted: Source: FCD1 2000, AUS-P09-007 Language Opportunity: The SQLSTATE corresponding to <i>FDW-specific condition — unable to create</i> <i>reply</i> is not sufficiently precise or informative. More specific diagnostic information is required.	
					Solution	
	NLD-P09-017		2-Minor Technical	P09-No specific location	MED-061 The following Language Opportunity has been noted: Source: WG3:YYJ-016 (USA-P09-018) Language Opportunity: MED's facility for communicating between the "local" SQL-server and the	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					foreign-data wrapper can be significantly enhanced by providing the ability to pass pre-parsed SQL statements or fragments of them from the SQL-server to the foreign-data wrapper. The most obvious choice for representing this information is in an XML format of some sort, preferably a parse tree or analog. Solution	
	NLD-P09-018		2-Minor Technical	P09-No specific location	None provided with comment. MED-066 The following Language Opportunity has been noted: Source: WG3:DRS-119 Language Opportunity: SQL/MED currently only provides read-only access for foreign tables. However, there are applications which require the ability to update data stored in those tables, this includes the ability to create new data and to delete existing data (UID - update, insert, delete). If and when this LO is addressed, changes applied to different foreign tables (possibly residing on different foreign servers) need to be handled according to ACID principles (atomicity, consistency, isolation, durability). Additionally, Subclauses are needed along the lines of those in Foundation, headed "Effect of inserting/ replacing/ deleting", plus extensions to existing DML Subclauses in Foundation that will cause these new Subclauses to be invoked when appropriate, to handle UID operations correctly. Furthermore, the underlying foreign-data wrapper interface needs to be enhanced to enable UID. It might also be desirable to be able to specify constraints as well as triggers on foreign tables.	
					None provided with comment.	
					SOL/OLB	
	NLD-P10-001		2-Minor Technical	P10-09.09 EntryInfo overview	OLB-002 The following Language Opportunity has been noted: Source: First FCD ballot, comment CAN-P10-017 Language Opportunity: The exact set of class of statements that Table 3, "Association of roles with SQLJ <executable clause="">s" refers to could be more explicitly defined. Solution None provided with comment.</executable>	
	NLD-P10-002		2-Minor Technical	P10-No specific location	OLB-003 The following Language Opportunity has been noted: Source: First FCD ballot, comments CAN-P10-023, CAN-P10-024, and CAN-P10-025 Language Opportunity: There may be many opportunities to replace D&Rs in SQL/OLB with an informative Note that says something like "Conformance to SQL/OLB requires	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					support only for the keywords", when referencing statements or other syntax	
					defined in Foundation or other parts.	
					Solution	
					None provided with comment.	
	NLD-P10-003		2-Minor	P10-No specific	OLB-004 The following Language Opportunity has been noted:	
			Technical	location	Source: First FCD ballot, comment CAN-P10-026, reinstated by WG3:ZSH-	
					04/=H2-2003-028	
					Canguage Opportunity:	
					SQL/OLD could belief in the supporting the optional LOCAL keyword in SQL/OLD could belief it in supporting the optional LOCAL keyword in	
					Solution	
					None provided with comment	
	NLD-P10-004		2-Minor	P10-09.07.03.	OI B-009 The following Language Opportunity has been noted:	
			Technical	Profile	Source: First FCD ballot, comment DEU-P10-014	
				customizer	Language Opportunity:	
				interface	Something needs to be said about how the operations "acceptsConnention" in	
					this subclause and in subclause 5.6, 'Customization interface', relate to each	
					other.	
					Solution	
					None provided with comment.	
	NLD-P10-005		2-Minor	P10-09.09,	OLB-010 The following Language Opportunity has been noted:	
			Technical	EntryInfo overview	Source: First FCD ballot, comment CAN-P10-018 and WG3:PER-098R1/H2-2001-059	
					Language Opportunity:	
					Table 4, "SQLJ type properties", must be extended to support the new SQL-99	
					data types (e.g., ARRAY, MULTISET, and ROW).	
					Support for ARRAY has been provided by WG3:DRS-080/H2-2002-458. It is	
					not anticipated that support for either MULTISE1 or ROW will be provided by	
					SQL/OLB until such time as JDBC provides such support.	
					Solution None provided with comment	
	NI D_P10.006		2-Minor	P10-09.09	OI B-011 The following Language Opportunity has been noted:	
	NLD-1 10-000		Technical	EntryInfo	Source: First FCD ballot comment CAN-P10-019	
			reennear	overview	Language Opportunity:	
					Table 4, "SQLJ type properties", must be extended to support the SQL-92 data	
					types not mentioned (e.g. DECIMAL, BIT, BIT VARYING, and INTERVAL).	
					Support for DECIMAL is provided via the java.sql.Types values NUMERIC	
					and DECIMAL. Further, per SQL/Foundaiton, Annex E, "Incompatibilities with	
					ISO/IEC 9075-2:2003", ISO/IEC 9075-2:1999 defined data types BIT and BIT	
					VARYING, but those types have been deleted from this edition of ISO/IEC	
					9075. It is not anticipated that support for INTERVAL will be provided by	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					SQL/OLB until such time as JDBC provides such support.	
					Solution	
-					None provided with comment.	
	NLD-P10-007		2-Minor	P10-No specific	OLB-014 The following Language Opportunity has been noted:	
			Technical	location	Source: First FCD ballot, comment DEU-P10-015	
					Deformed: P10 SOL /OL P. 06.01 Grammar notation (which no longer existed)	
					This subclause contains conventional material that has traditionally been	
					provided in other parts of 9075 as a subclause of Clause 3. In order to avoid a	
					major rewrite, such a Conventions Subclause should be added to this part; it	
					should merely outline where and how the information one would have expected	
					at that clause is actually provided in this part of 9075.	
					Solution	
					None provided with comment.	
	NLD-P10-008		2-Minor	P10-No specific	OLB-015 The following Language Opportunity has been noted:	
			Technical	location	Source: First FCD ballot, comment GBR-P10-019	
					Deference: P10, SOL/OLP, 00.05 (no title given)	
					"Binary portability" more properly "portability of intermediate object code	
					representation", is an objective of the originators of the Java language. Clause 5	
					does not sufficiently distinguish between statements of intent, tutorial matter and	
					concrete specification. The clause should be merged into the general Concepts	
					clause, and should be further revised to clarify the distinction between things	
					that are part of the SQLJava binding and features of those things that are part of	
					Java.	
					Solution	
	NLD-P10-009		2-Minor	P10-No specific	OI B-017 The following Language Opportunity has been noted:	
			Technical	location	Source: First FCD ballot, comment USA-P10-025	
					Language Opportunity:	
					This document contains "Definitions and Rules" clauses that sometimes appear	
					analogous to SQL "Syntax Rules" and sometimes like "General Rules".	
					However, unlike "Syntax Rules" and "General Rules" there is no general	
					specification of the effect of violating a "Definition and Rules" nor of when the	
					Definition and Kules are validated/performed. The validation time of and the	
					Solution	
					None provided with comment	
	NLD-P10-010		2-Minor	P10-No specific	OLB-018 The following Language Opportunity has been noted:	
			Technical	location	Source: First FCD ballot, comment DEU-P10-020	
					Language Opportunity:	

SEQ	Cmnt	See		-		Addressed
#	ID	Also	Severity	Reference	Description	Ву
					To improve readability, more cross-references are needed. E.g., when the	
					interfaces are specified that are implemented by some class definition (see for	
					instance 10.2.1) it would be helpful to have reference to the subclause defining	
					that interface. The author of this comment is aware that there is abundant	
					precedence for such cross-references in the document (see "See also" sections).	
					Solution	
					None provided with comment.	
	NLD-P10-011		2-Minor	P10-No specific	OLB-025 The following Language Opportunity has been noted:	
			Technical	location	Source: Email from Fred Zemke, 2001-11-05, from unknown source	
					Language Opportunity:	
					There are many paragraphs that say "An SQLException will be thrown" without	
					saying what that condition is! Is the implementation free to raise any exception	
					that it feels like, possibly even one chosen randomly?	
					If not, then the standard must say what condition is thrown! These places are	
					usually accompanied by an editor's note, which should be removed whenever the	
					problem at that location is solved.	
					Solution	
					None provided with comment.	
	NLD-P10-012 2-Minor	2-Minor	P10-04.09,	OLB-028 The following Language Opportunity has been noted:		
			Technical	Default connection context	Source: WG3:ZSH-153R1 = H2-2002-153R1	
					Language Opportunity:	
					A problem arises because both SQL/OLB and SQL/JRT have mechanisms for	
					referencing their default SQL-environment. In SQL/OLB, the JNDI registered	
					"jdbc/defaultDataSource" name will, if present, identify the default data source	
					for SQL operations to be performed against. In SQL/JRT, the JDBC URL	
					"jdbc:default:connection" identifies a JDBC connection to the current	
					SQLimplementation, SQL-session, and SQL-transaction. This raises the	
					question: When, if ever, are the following logically equivalent?	
					1) Connection con = DriverManager.getConnection("jdbc:default:connection");	
					2) Connection con =	
					sqlj.runtime.ref.DefaultContext.getDefaultContext().getConnection();	
					Connection();	
					4) Context ctx = new InitialContext(); DataSource ds = (DataSource)	
					ctx.lookup("jdbc/defaultDataSource"); Connection con = ds.getConnection();	
					That is, when is the java.sql.Connection con, appearing in the above code	
					sequences providing a JDBC connection to the same SQL-implementation? We	
					believe it is desirable, if not required, for an application to be able to run either	
					inside a database as a "stored procedure" or outside as a regular application	
					without having to be recoded, so we ask that above be issue be clarified.	
					Solution	
					None provided with comment.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
	NLD-P10-013		2-Minor	P10-No specific	OLB-029 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:ZSH-153R1 = H2-2002-153R1	
					Language Opportunity:	
					SQL/OLB should make it possible for an SQL/OLB application to use the JDBC	
					3.0 support of what JDBC 3.0 refers to as 'Auto Generated Keys' without having	
					to use JDBC to do so. This capability is often used to access what many DBMSs	
					refer to as a Row ID of a just inserted or updated row. And while	
					SQL/Foundation doesn't standardize a Row ID, the facility would have utility by	
					anowing access to what SQL/Foundation refers to as identity columns of	
					Solution	
					Solution	
	NI D D10 014		2 Minor	P10 No specific	OI P 020 The following Language Opportunity has been noted:	
	NLD-F10-014		Z-Millor Technical	location	Source: Email from Mark Ashworth 2004-07-22 SIA Action Item (see minutes)	
			reennear		for SIA-025)	
					Language Opportunity:	
					WG3:SIA-025R1. "A Shorthand for Getting ALL Diagnostics" proposes to add	
					a new diagnostics option to embedded SQL but ignores the question of whether	
					it should also be added to CLI or the SQLJ (SQL/OLB) binding.	
					Solution	
					None provided with comment.	
					SQL/Schemata	
	NLD-P11-001		1-Major	P11-05.09,	SCHEM-029 The following Possible Problem has been noted:	
			Technical	APPLICABLE_R	Source: WG3: SIA-026r3	
				OLES view	Possible Problem:	
					The function and definition of the Information Schema view	
					APPLICABLE_ROLES are given in Subclause 5.9, "APPLICABLE_ROLES	
					view":	
					Function	
					Identifies the applicable roles for the current user.	
					ROLE NAME.	
					IS GRANTABLE) AS	
					((SELECT GRANTEE, ROLE_NAME, IS_GRANTABLE	
					FROM DEFINITION SCHEMA.ROLE AUTHORIZATION DESCRIPTORS	
					(CURRENT USER 'PUBLIC')	
					OR	
					GRANTEE IN	
					(SELECT ROLE NAME	
					FROM ENABLED ROLES)))	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					(SELECT RAD.GRANTEE, RAD.ROLE NAME, RAD.IS GRANTABLE FROM DEFINITION SCHEMA.ROLE AUTHORIZATION DESCRIPTORS RAD JOIN APPLICABLE ROLES R ON RAD.GRANTEE = R.ROLE NAME)); The text shown underlined is redundant. It was proposed by ICN-039 as a replacement for the CURRENT_ROLE that had previously been the second element of the first IN list, having been proposed -erroneously, we believe - by PER-193. Before PER-193, that IN list was merely "(CURRENT_USER, "PUBLIC')", which was consistent with the stated Function of Subclause 5.9, "APPLICABLE_ROLES view", "Identifies the applicable roles for the current user". Of course the current role, if there is one, is a role that is applicable for the current user, if there is one. In Part 2 SQL/Foundation, Subclause 18.3, " <set role statement>", GR4) makes sure of that (and in fact applies, possibly erroneously, an even stronger condition). It appears, then, that the text shown underlined should be deleted. However, we hesitate to propose that because we are uncertain as to the real purpose of the APPLICABLE_ROLES view, considering that there isn't always a current user. What roles, if any, are deemed to be applicable, "for" what, when the top cell of the authorization stack of the current SQL-session contains a role name and no user identifier? Is that role name included in the answer? Solution None provided with comment.</set 	,
	NLD-P11-002		2-Minor Technical	P11-06.21, DATA_TYPE_DE SCRIPTOR base table	SCHEM-033 The following Possible Problem has been noted: Source: WG3:STX-050 Comment WG3-P11-023 Possible Problem: There is no foreign key check for the columns SCOPE_CATALOG, SCOPE_SCHEMA, and SCOPE_NAME to the tables table. It is not clear to the Author, if this reference has to be in the same CATALOG or not Solution None provided with comment.	
	NLD-P11-003		2-Minor Technical	P11-06.41, SCHEMATA base table	SCHEM-031 The following Possible Problem has been noted: Source: WG3:STX-050 Comment WG3-P11-020 Possible Problem: A foreign key between the table SCHEMATA and the table CHARACTER_SETS is missing. It is not clear to the author, if this Character Set has to reside in the same catalog. If this is the case, the following constraint would resolve the problem: CONSTRAINT SCHEMATA_FOREIGN_KEY_CHARACTER_SETS FOREIGN KEY (DEFAULT_CHARACTER_SET_CATALOG,	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					DEFAULT_CHARACTER_SET_SCHEMA,	
					DEFAULT_CHARACTER_SET_NAME)	
					REFERENCES CHARACTER_SETS	
					Solution	
					None provided with comment.	
	NLD-P11-004		2-Minor	<i>P11–06.62</i> ,	SCHEM-032 The following Possible Problem has been noted:	
			Technical	USER_DEFINE	Source: WG3:STX-050 Comment WG3-P11-021	
				D_TYPES base	Possible Problem:	
				table	The value list and the select list of the last query of the constraint	
					USER_DEFINED_TYPES_CHECK_SOURCE_TYPE do not match. They have	
					different number of elements. It reads:	
					(USER_DEFINED_TYPE_CATALOG, USER_DEFINED_TYPE_SCHEMA,	
					USER_DEFINED_TYPE_NAME, SOURCE_DTD_IDENTIFIER) IN	
					(SELECT OBJECT_CATALOG, OBJECT_SCHEMA, OBJECT_NAME,	
					OBJECT_TYPE, DTD_IDENTIFIER	
					FROM DATA_TYPE_DESCRIPTOR	
					Solution	
					None provided with comment.	
	NLD-P11-005		2-Minor	<i>P11-06.11</i> ,	SCHEM-002 The following Language Opportunity has been noted:	
			Technical	CHARACTER_S	Source: DCOR 2000, SWE-STC-030	
				ETS base table	Language Opportunity:	
					This base table contains a bare minimum of information. It could be enhanced to	
					indicate relationships among character sets, for example whether the character	
					set is standard, implementation-defined, or userdefined, and what characater set	
					a user-defined charater set is based on.	
					Solution	
					None provided with comment.	
	NLD-P11-006		2-Minor	<i>P11-06.44</i> ,	SCHEM-008 The following Language Opportunity has been noted:	
			Technical	SQL_IMPLEME	Source: WG3:PER-118/H2-2001-???	
				NTATION_INFO	Language Opportunity:	
				buse tuble	Subclause 6.44, "SQL_IMPLEMENTATION_INFO base table", is defined to	
					contain information about SQL-implementation information items (identified by	
					name and number) but these items are not defined in the other parts of the	
					standard.	
					Solution	
					None provided with comment.	
	NLD-P11-007		2-Minor	<i>P11-06.46</i> ,	SCHEM-009 The following Language Opportunity has been noted:	
			Technical	SQL_SIZING	Source: WG3:PER-118/H2-2001-???	
				base table	Language Opportunity:	
					Subclause 6.46, "SQL_SIZING base table", is defined to contain information	
					about SQL sizing items (identified by name and number) but these items are not	
SEQ	Cmnt	See				Addressed
-----	--------------	---------	-----------------------	--	---	-----------
#	ID	Also	Severity	Reference	Description	By
					defined in the other parts of the standard. (Subclause 6.44,	
					"SQL_IMPLEMENTATION_INFO base table", has the same problem.).	
					Solution	
					None provided with comment.	
	NLD-P11-008		2-Minor	<i>P11-06.21</i> ,	SCHEM-013 The following Language Opportunity has been noted:	
			Technical	DATA_TYPE_DE	Source: WG3:PER-100R2/H2-2001-062R2	
				SCRIPTOR base	Language Opportunity:	
				table	Paper WG3:PER-100r2 noted the following Language Opportunity:	
					The user may wish to recover the original type declaration, rather than the	
					equivalent type declaration that is used by the SQL-server. This concern could	
					be met by adding columns such as ORIGINAL_DATA_TYPE and	
					ORIGINAL_PRECISION to the DATA_TYPE_DESCRIPTORS base table, as	
					well as all views that draw from it. These new columns should be part of a new	
					conformance feature, to make them optional, since not every implementation	
					will be able to display the original type declaration.	
					Solution	
					None provided with comment.	
	NLD-P11-009		2-Minor	<i>P11-06.21</i> ,	SCHEM-014 The following Language Opportunity has been noted:	
			Technical	DATA_TYPE_DE	Source: WG3:PER-100R2/H2-2001-062R2	
				scripion base	Language Opportunity:	
				luble	Paper WG3:PER-100r2 noted the following Language Opportunity:	
					Users might be interested to know the largest and smallest exponents	
					accomodated by the approximate numeric types.	
					Solution	
					None provided with comment.	
	NLD-P11-010	2-Minor	P11-06.21,	SCHEM-015 The following Language Opportunity has been noted:		
			Technical	SCRIPTOR hase	Source: WG3:PER-100R2/H2-2001-062R2	
				table	Language Opportunity:	
					Paper wG3:PER-100r2 noted the following Language Opportunity:	
					A capability would be a table that simply listed all the data type equivalences of the numeric data types	
					ne numeric data types.	
					Solution	
	NI D D11 011		2 Minor	P11.05.11	None provided with comment.	
	NLD-P11-011		2-ivinor Technical	ATTRIBUTES	SUPEN-016 The following Language Opportunity has been noted:	
			recinical	view	Source, $DCON/2002$ USA-51C-046 and $WOS.ZSH-155K1 = H2-2002-155K1$	
					Language Opportunity. The function statement says that this view lists the attributes of structured types	
					that the user has access to "Access" is ambiguous. As the view is currently	
					defined it appears to mean "has USAGE or UNDER privilege on" This could	
					be solved by changing "that are accessible" to "that the user has USAGE or	
					UNDER privilege for". However this comment will not suggest that solution.	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					Instead, this comment will point out that there are ways to define access to an attribute other than USAGE or UNDER privilege on the attribute's type. First, there are other ways to access the type than through USAGE privilege. The type might be the parameter type of an SQL-invoked routine that the user can execute, it might be the return type of a regular function or method that the user can execute, it might be the type of a column that the user can SELECT, or the type of a selectable nested site such as the field of a row, the element type of a collection, or the attribute type of a different structured type. All of these constitute "access" to a structured type. An analogy can be drawn between user-defined types and domains. Note that the DOMAINS view shows not just the domains that the user can access. After defining accessible types, you have the question of what makes an attribute accessible. Is it EXECTUE privilege on the observer? Or perhaps EXECUTE on either the observer or the mutator? Or some other criterion?	
					None provided with comment.	
	NLD-P11-012		2-Minor Technical	P11-05.73, USER_DEFINE D_TYPES view P11-05 20	SCHEM-019 The following Language Opportunity has been noted: Source: DCOR/2002 USA-STC-059 and WG3:ZSH-153R1 = H2-2002-153R1 Language Opportunity: The word "accessible" in the function is ambiguous. What is meant is those user-defined types for which the user has USAGE or UNDER privilege. However, it is questioned in a separate comment on the ATTRIBUTES view whether "accessible" should be limited to types with USAGE or UNDER privilege. Note that DIRECT_SUPERTYPES view will reveal type T's existence if T is the direct supertype of a type T2 for which the user has USAGE or UNDER privilege, even if the user does not have USAGE or UNDER privilege on T itself. This seems inconsistent. Also, COLUMNS view will display the type T if there is a column whose type is T. It is suspected, but not verified, that the type will also be visible in other views of the Information Schema, wherever the type of a site is displayed (for example, ATTRIBUTES view, FIELDS view, ROUTINE view, PARAMETERS view). Note that DOMAINS view shows a domain if either the user has USAGE privilege on the domain or the user has SELECT privilege on a column whose type is the domain; this provides a precedent that "accessible" is not limited to "has a privilege on". Solution None provided with comment. SCHEM-021 The following Language Opportunity has been noted:	
	NLD-P11-013		2-Minor Technical	P11-05.20, "COLUMN_UD T_USAGE view	SCHEM-021 The following Language Opportunity has been noted: Source: WG3:ZSH-153R1 = H2-2002-153R1 Language Opportunity: DCOR comment USA-STC-049 pointed out that the join condition joining	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					DEFINITION_SCHEMA.COLUMNS with DEFINITION_SCHEMA.SCHEMATA assumes that the former table has columns named USER_DEFINED_TYPE_CATALOG and USER_DEFINED_TYPE_SCHEMA, which it does not. That comment goes on to suggest that perhaps the intent was to join in DATA_TYPE_DESCRIPTORS, which does. However, if the suggestion in USA-STC-049 were followed, this would not really be sufficient to find all columns that are dependent on a user- defined type. What about columns that are row types with a field that is a user- defined type? Or collection types with an element type that is a user- defined type? Or collection types with an element type that is a user- defined type? See the notion of usage-dependent added to Foundation by WG3:YYJ- 083r1. Note that in that paper, it is argued that the notion of usage-dependency does not need to recurse through attributes of a structured type. While this argument is sufficient for the purpose of enforcing RESTRICT or CASCADE semantics, and justifiable for Access Rule checking, does it make sense for this view? For example, if type T1 has an attribute of type T2, and column C1 is of type T1, does C1 depend on T2 in the meaning of this view? If the user is using the view to find all columns to drop before dropping type T2, then the user wants to see C1 in this view. The alternative is that the user must do his own recursion (find all UDTs that depend on T2, then find all columns that depend on any of them.) Solution	
	NLD-P11-014		2-Minor Technical	P11-06.44, SQL_IMPLEME NTATION_INFO base table	None provided with comment. SCHEM-022 The following Language Opportunity has been noted: Source: WG3:ZSH-153R1 = H2-2002-153R1 Language Opportunity: There is no list of values for IMPLEMENTATION_INFO_ID and IMPLEMENTATION_INFO_NAME. It seems that many of these were intended to be the codes used in CLI by GetInfo (see for example CLI GetInfo GR 10) subrules b), c), p) and q).) The writer of this comment does not know if there are codes that are necessary for CLI or other parts of SQL. But see CLI subclause 7.1 SQL_IMPLEMENTATION_INFO base table. Solution None provided with comment.	
	NLD-P11-015		2-Minor Technical	P11-No specific location	SCHEM-023 The following Language Opportunity has been noted: Source: WG3:ZSH-153R1 = H2-2002-153R1 Language Opportunity: Implementations should not be required to expose columns about optional features that they don't support. For example, in Subclause 5.22, "COLUMNS view", the IS_SELF_REFERENCING column is meaningful only if Feature S051, "Create tables of type", is implemented. If conformance to that feature is not claimed,	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					then references to the column should be prohibited.	
					Solution	
					None provided with comment.	
	NLD-P11-016		2-Minor	P11-05,	SCHEM-026 The following Language Opportunity has been noted:	
			Technical	Information	Source: DBL:LGW-152/X3H2-97-352 (also DBL:LGW-023/X3H2-97-044,	
				Schema	SEQ# 406, USA-102*)	
					Language Opportunity: The DOUTINES view and base table have columns that contain the timestamn	
					of when the routine was CREATED and LAST ALTERED. These are	
					analogous to the file creation and modification timestamps typically provided by	
					a file system These timestamps are useful for comparing the creation and	
					modification timestamps of the database objects with the timestamps in an	
					external source code control and configuration management utility. Since SQL3	
					supports extensive programmatic capabilities this configuration management	
					support is extremely useful. However it does not go far enough. Created and	
					Last_altered timestamps would also be useful in the following base tables and	
					their associated views:	
					- ABSTRACT_DATA_TYPES	
					- DUMAINS	
					- IADLES VIEWS	
					- COLUMNS	
					- ASSERTIONS	
					- CHARACTER SETS	
					- COLLATIONS	
					- TRANSLATIONS	
					— TRIGGERS	
					— SUB_TABLES	
					Solution	
					None provided with comment.	
	NLD-P11-017		2-Minor	P11-05,	SCHEM-027 The following Language Opportunity has been noted:	
			Technical	Schema	Source: DBL:LGW-152/X3H2-97-352 (also DBL:LGW-023/X3H2-97-	
				Senema	044, SEQ#409, USA-105)	
					Language Opportunity:	
					Many "information discovery" products depend upon full text searches of	
					document databases to feed the indexing mechanisms used in their search	
					engines. It is very difficult to extend this technique to "structured"	
					relational databases especially if they have high numeric content unless	
					there is some textual description of the semantics associated with data	
					values and schema objects. Sometimes "information discovery" agents	
					will search the INFORMATION_SCHEMA Catalog Schema Table and	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
					Column names looking for relevant key word "stems" to feed to the	
					search engine. It would be very helpful to users of such agents if there	
					were a "standard" way to read and write textual descriptions of what each	
					schema object represents. Certainly Information Resource Dictionary	
					Systems (IRDS) could help in this task or users could define a special	
					schema for this purpose but at present there is no dependable standard	
					mechanism to make such information available to outside agents. One	
					easy-to-implement yet very helpful facility would be to associate a	
					"COMMENT" or "DESCRIPTION" column with each relevant table in	
					the INFORMATION_SCHEMA together with a "SET SCHEMA	
					COMMENT statement" (or other appropriate syntax) that would allow	
					the owner of a schema object to set and/or modify the COMMENT	
					column associated with it. The normal Information Schema view	
					definition would then determine which users are able to read the	
					COMMENT column so information discovery agents would be able to	
					"discover" whatever comments exist for PUBLIC schema objects and	
					report back to their creators any interesting database content.	
					In addition to information discovery agents comment or description	
					information is crucial to support the reusability of ADTs. An SQL	
					programmer must know what an ADT is supposed to do in order to	
					correctly utilize or subtype it. This information can only be provided by	
					the ADT creator in a text format and is much more likely to be useful if	
					stored in the INFORMATION_SCHEMA than if stored in paper	
					documentation at the bottom of a stack on someone else's desk. This	
					could be accomplished by adding syntax to the ADT definition to	
					support a large amount of text.	
					The SQL objects for which comment/description information would be	
					useful include: DOMAINS, TABLES, VIEWS, COLUMNS,	
					ASSERTIONS, CHARACTER_SETS, COLLATIONS,	
					TRANSLATIONS, TRIGGERS, SUB_TABLES, as well as distinct	
					types, abstract data types, and SQL-invoked routines.	
					Solution	
					None provided with comment.	
	NLD-P11-018		1-Major	P11-No specific	SCHEM-028 The following Language Opportunity has been noted:	
			Technical	location	Source: WG3:HBA-034R2 = H2-2003-343R4	
					Language Opportunity:	
					We have at least the following kinds of SQL-schema objects that might be	
					involved in a dependency relationship	
					— cneck constraints	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
			~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~		- assertions	
					— generated columns	
					— SOL-invoked routines	
					- triggers	
					— views	
					character sets	
					— collations	
					— transliterations	
					— domains	
					— non-generated columns	
					— base tables	
					— sequences	
					— user-defined types	
					We do not have Information Schema views to report all possible dependencies	
					between these kinds of SQL-schema objects, as seen in this table:	
					(See LO)	
					In the preceding table,	
					• A blank cell means that the dependency cannot occur.	
					• 'y' means that an Information Schema view exists to report such dependencies.	
					• 'N' means that such dependencies can occur, but there is no reporting	
					mechanism currently.	
					Most of the possible dependencies are explained as follows:	
					• A <value expression=""> can be a <scalar subquery="">, which can be a grouped</scalar></value>	
					query, which can depend on a check constraint, assertion, or unique constraint in	
					order to deduce a functional dependency.	
					Thus, anything that permits a <value expression=""> can be dependent on a check</value>	
					constraint, assertion or unique constraint (but only if Feature 1301, "Functional	
					dependencies is supported).	
					• CAST to a character string type with a user-defined character set implies a	
					might be dependent on a character set. Thus anything permitting a <value expression=""></value>	
					A collection can be used in comparison productor, and thus most kinds of SQL	
					- A contation can be used in comparison predicates, and thus most kinds of SQL-	
					Schema objects might depend on a containa. • $\Lambda < value expression > can contain a CONVERT expression, which depends on$	
					a transliteration so most kinds of SOL scheme objects might depend on a	
					transliteration. Conversely, a transliteration uses an SOL invoked routing, so a	
					transliteration can be dependent on anything that a routine can be dependent on	
					Speaking in orders of magnitude, if we have n kinds of SOL scheme objects	
					and we add one more, then we have $(n + 1)^2 - n^2 = 2n + 1$ new kinds of	
					dependency to think about. Thus the cost of adding one kind of SQL-schema	

SEQ	Cmnt	See	a	D		Addressed
#	ID	Also	Severity	Reference	Description	Ву
					object is potentially $2n + 1$ new kinds of dependency. A side from the fact that so many kinds of dependency are currently	
					unsupported [Fred Zemke thinks] that the technique of creating one Information	
					Schema view for each kind of dependency has become unmanageable for our	
					users, and unmaintainable for ourselves. Therefore, [Fred thinks] it is time to	
					come up with a different model for dependency tracking and reporting. [Fred	
					thinks] the correct approach is to define a base table to track immediate	
					dependencies between all kinds of SQL-schema objects, and a recursive	
					Information Schema view that shows all deducible dependencies.	
					Solution	
				<b>D11</b> 00	None provided with comment.	
	NLD-P11-019		2-Minor	P11-00, Definition	SCHEM-030 The following Language Opportunity has been noted:	
			Technical	Schema	Source: wG3:SIA-018 = ANSI NC11S H2-2003-429 / email from Joern Bartels	
					SIA-018 adds the Subclause 10.11 "Determination of view and view component	
					privileges" to Part 2.	
					This subclause introduces the new view privilege dependency descriptor.	
					There is no corresponding base table in Clause 6, "Definition Schema" of Part	
					11 defined. As this descriptor is created in the Subclause 12.1, " <grant< td=""><td></td></grant<>	
					statement>" of Part 2 and used in the Subclause 12.7, " <revoke statement="">" of</revoke>	
					Part 2, it needs to be stored somewhere.	
					Solution	
					None provided with comment.	
					SQL/JRT	
	NLD-P13-001		2-Minor	P13-No specific	JRT-001 The following Language Opportunity has been noted:	
			Technical	location	<b>Source:</b> WG3:YYJ-041 = H2-2001-405	
					Language Opportunity:	
					java io Serializable	
					or any Java equivalent". This would also permit, for example, implementing	
					Externalizable, which can often be done with better performance and space	
					usage than Serializable.	
					Solution	
					None provided with comment.	

## UK Ballot Comments on ISO/IEC 9075-2 CD — 2005-03-09

SEQ	Cmnt	See				Addressed By
#	ID	Also	Severity	Reference	Description	
				{ <b>CI</b>	D} SQL/Foundation	
	GBR-P02-001		2-Minor	P02-02-01, JTC1	Some of the standards identified in this subclause have been superseded by later	
			Technical	standards	editions. For example:	
					ISO/IEC 1539-1:2004, also ISO/IEC1539-2:2000 may be relevant	
					ISO/IEC 1989:2002	
					ISO 8001:2004 ISO/IEC 8640 now has Amd1:1007 and Amd2:1008	
					ISO 9899 1999 now has a Cor ² 2004	
					ISO/IEC 10646:2003 – now a single-part standard	
					ISO/IEC 14651:2001 now has Amd1:2003	
					These later versions should be reviewed and, if found appropriate, replace the	
					versions referenced in this subclause. Other referenced standards should also be	
					reviewed for continued suitability.	
					Solution	
					None provided with comment.	
	GBR-P02-002		2-Minor	<i>P02-03-01-03</i> ,	Since several editions of The Unicode Standard are included in subclause 2.2, a	
			Technical	from Unicode	specific edition containing the definitions should be identified here.	
				from oncoue	Alternatively, it may now be more appropriate to reference ISO/IEC 10646.	
					Solution	
	GDD D02 002		0.15	D02.04.02.05	None provided with comment.	
	GBR-P02-003		2-Minor	P02-04-02-05, Character	Since ISO/IEC 10646 is now a single-part standard, the references to its Part 1	
			Technical	encoding forms	and that part's annexes in the sixth and seventh bullets of the third paragraph	
				00	fifth bullet could probably also be replaced by a reference to ISO/IFC 10646	
					Solution	
					None provided with comment.	

SEQ	Cmnt	See				Addressed By
#	ID	Also	Severity	Reference	Description	
	GBR-P02-004	GBR-P02-015 GBR-P10-005	4-Minor Editorial	P02-04-23, Embedded syntax	The in-line list of expansions of <embedded host="" program="" sql=""> that occupies most of the first three lines of the first paragraph is both unnecessary and difficult to read. Removing the list now would not cause any problem to the reader and would also be a minor reduction to the maintenance problem should the list of standard embeddings ever change.</embedded>	
					Solution	
	GBR-P02-005		2-Minor Technical	P02-05-04, "Names and identifiers"	See comment.         SR28) attempts to deal with descriptor names and "extended" names:         Figure 3       In a <descriptor name="">, <extended name="" statement="">, or         <extended cursor="" name="">, if a <scope option="">is not         specified, then a <scope option="">of LOCAL is implicit. If         a <scope option="">is contained in an <sql-schema< td="">         statement&gt; then it shall not contain LOCAL.         The last sentence does not explicitly deal with the case of <scope option=""> being         missing in a <descriptor name="">, <extended name="" statement="">, or <extended< td="">         cursor name&gt; contained in an <sql-schema statement="">. As the default is         LOCAL, it appears that such an <sql-schema statement=""> is invalid—in other         words, if one wants to write one of these names inside an <sql-schema< td="">         statement&gt;, one can do so only if the key word GLOBAL is given for <scope< td="">         option&gt;. We doubt if that is the intention and we doubt if any existing products         claiming conformance enforce such a rule.         The wording being commented on was introduced by our own change proposal,         STX-036, "Follow-up to SAI-033, Action Item STX-012 m)". We regret that         we have no solution to offer at this time, but we would hope to be able to         collaborate on such a solution during the CD ballot resolution process.</scope<></sql-schema<></sql-schema></sql-schema></extended<></extended></descriptor></scope></sql-schema<></scope></scope></scope></extended></extended></descriptor>	
					Solution None provided with comment.	

SEQ	Cmnt	See				Addressed By
#	ID	Also	Severity	Reference	Description	
	GBR-P02-006	GBR-P02-007	2-Minor Technical	P02-06-01, <data type=""></data>	Syntax Rule 23 specifies that a DECIMAL numeric item may be maintained by the implementation with a precision greater than that specified in the declaration of the item. There may be good practical reasons for this behavior in some SQL implementations. However, this implementation-determined value is used when the corresponding descriptor is created, and the value specified by the user is lost. This may not be of any significance within the context of a single SQL implementation, but could have adverse effects in environments where schemata and SQL data are exported from one SQL implementation and imported into another. It is possible that after a round trip the precision of a data item may be increased and may be much larger than the originator ever expected. There should be some facility whereby both the user-specified and implementation-determined precision are available to an export tool (or to anything else) from the Information Schema.	
	GBR-P02-007	GBR-P02-006	3-Major Editorial	P02-06-01, <data type=""></data>	Subrule a) in each of General Rules 11 and 12 defines specifies in a fairly long- winded way that there is a normalised data type corresponding to each of the numeric data types. The mapping appears to be intended to be a static one, fixed for each SQL implementation. It would be better if the subrules were moved to become Syntax Rules, or sub-rules associated with the Syntax Rules that describe the meanings of, and constraints on <precision> and <scale>. Solution None provided with comment.</scale></precision>	

SEQ	Cmnt	See				Addressed By
#	ID	Also	Severity	Reference	Description	
	GBR-P02-008		1-Major Technical	P02-07.06, " <table reference&gt;"</table 	SQL supports no immediate counterpart of the relational RENAME operator, which, given a relation $rl$ returns a relation that is identical to $rl$ except that one or more attributes differ in name with respect to their counterparts in $rl$ . RENAME is a very useful way of setting up appropriate operands for invocation of operators like JOIN and UNION, whose SQL counterparts are (near enough) NATURAL JOIN and UNION CORRESPONDING.	
					That SQL has no adequate counterpart to RENAME becomes apparent when one considers the problem of selecting all the columns from a 300-column table except that one of those 300 is to be renamed. One has to write out the names of the other 299, which is not only tiresome and error-prone, but also renders oneself vulnerable to the later addition of further columns to the table in question.	
					A tentatively suggested solution is to support syntax in  as in this example:	
					EMP RENAME ( NAME AS EMPNAME, SALARY AS PAY ) AS E	
					The example shows a simple table name (EMP), but this could be replaced by a <query expression=""> as in <derived table="">s. The provision of the correlation name E is optional, here, of course, but we include it in the example to draw attention to the fact that any new syntax to follow a table name needs to be carefully thought out, especially if it uses the key word AS.</derived></query>	
					If this comment is closed as a language opportunity but arouses sufficient sympathy from other national bodies, it could be addressed by a change proposal submitted during FCD ballot resolution.	
					Solution	
					None provided with comment.	

SEQ	Cmnt	See				Addressed By
#	ID	Also	Severity	Reference	Description	
	GBR-P02-009		1-Major Technical	P02-07.12, " <query specification&gt;"</query 	The syntax allows for a <select list=""> to be an <asterisk> but does not allow it to include an <asterisk> along with other things. Instead one has to write a <qualified asterisk=""> when one wants all the input columns plus one or more calculated columns, for example.</qualified></asterisk></asterisk></select>	
					In other words, the suggestion is that one should be able to write, for example:	
					SELECT *, SALARY + BONUS AS PAY FROM EMP E, DEPT D WHERE E.D# = D.D#	
					instead of having to write	
					SELECT E*, D.*, SALARY + BONUS AS PAY FROM EMP E, DEPT D WHERE E.D# = D.D#	
					If this comment is closed as a language opportunity but arouses sufficient sympathy from other national bodies, it could be addressed by a change proposal submitted during FCD ballot resolution.	
					Solution	
					None provided with comment.	
	GBR-P02-010		1-Major Technical	P02-07.12, " <query specification&gt;"</query 	SQL's nearest counterpart of the relational projection operator is SELECT DISTINCT <i>list of column names</i> FROM <i>t</i> . Unfortunately this supports only projection by explicitly naming the attributes to be included. In practice, it is sometimes more convenient to list the attributes that are not required, especially when the operation is being used for some intermediate result.	
					The existing construct * might provide a convenient place to slot in projection by exclusion, as in SELECT * EXCEPT ( <i>column-name-list</i> ) FROM <i>t</i> or SELECT <i>t</i> .* EXCEPT ( <i>column-name-list</i> ) FROM <i>t</i>	
					If this comment is closed as a language opportunity but arouses sufficient sympathy from other national bodies, it could be addressed by a change proposal submitted during FCD ballot resolution.	
					Solution	
					None provided with comment.	

SEQ #	Cmnt ID	See Also	Severity	Reference	Description	Addressed By
	GBR-P02-011		1-Major Technical	P02-14, "Data manipulation"	At the WG3 meeting in St Croix, 2004, the UK's national body opening comments (see TXL-001, Agenda item 7.18) included this:	
					One requirement in particular that has been briefly discussed in the UK is an extension of the existing "multiple assignment" capability to encompass database updating operations. The idea is to be able to submit a sequence of two or more INSERT, UPDATE, DELETE, or MERGE operations as a single, atomic statement. The motivation for the idea is to overcome updating difficulties in the presence of "cyclic" constraints without any need to invoke the suspect deferred constraint checking mechanism (though we would not seek to remove that suspect mechanism from the standard, having failed in that endeavour several times already). We would like to gauge the interest of other national bodies in this idea. In particular, if there <i>is</i> any such interest, we would welcome any suggestions as to how triggers and cascaded referential actions might be handled—these are already a major problem with the deferred constraint checking mechanism and they would seem to pose problems for the multiple assignment approach too.	
					As we heard of some interest in, and no opposition to, this idea, we repeat it here. Some have questioned whether the idea really solves the problem, wrongly perceiving the proposed syntax as being equivalent to an atomic compound statement. To clarify, we point out that multiple assignment requires all the source expressions to be effectively evaluated before any of the constituent single assignments is executed. Thus, in the following example:	
					INSERT INTO T1 VALUES 1, 2, 3, INSERT INTO T2 VALUES foo ( 1 ) ;	
					if the body of foo references T1, the value yielded by that reference is the value prior to the insertion of VALUES 1, 2, 3. Now, it might be that if the two INSERTs are executed as statements in their own right in the order given, the first one fails with a foreign key violation. It might also be that if they are executed in reverse order, the insert into T2 fails with a foreign key violation. Finally, it might also be—and this is the point—that executing them in the order shown but with checking of the foreign key constraints deferred, the body of foo runs into some exception because it has been coded on the assumption that the database is consistent with the declared constraints.	
					The "multiple assignment" approach ensures that no constituent statement fails on account of an integrity check and also that no evaluation of a source expression that references the database is against an inconsistent state of the database.	
					If this comment is closed as a language opportunity but arouses sufficient	

SEQ	Cmnt	See				Addressed By
#	ID	Also	Severity	Reference	Description	
	GBR-P02-012		2-Minor	P02-14-01,	A recent GBR paper that was accepted for WD 9075-2 concerned the	
			Technical	<declare< td=""><td>specification of what were called 'dynamic results sets', which resulted in a new</td><td></td></declare<>	specification of what were called 'dynamic results sets', which resulted in a new	
				cursor>	definition for result set: "A sequence of rows brought into existence by opening	
					a cursor and ranged over by that cursor". The main reason for this addition was	
					that the object created by opening a cursor was not a table, as it could have	
					additional properties, such as ordering. However the consequences of this	
					realisation have not been fully applied.	
					This clause describes how to produce a 'result table', which should be a 'result	
					set'. However, declaring a cursor does not of itself produce a result set (or a	
					table), so some parts of this clause should be moved to clause 14.28, "Effect of	
					opening a cursor".	
					Solution	
					None provided with comment.	
	GBR-P02-013	GBR-P02-011	2-Minor	<i>P02-14-03</i> ,	The Function is given as "Position a cursor on a specified row of a table and	
			Technical	<fetch< td=""><td>retrieve values from that row." This reference to 'table' should be to 'result set'.</td><td></td></fetch<>	retrieve values from that row." This reference to 'table' should be to 'result set'.	
				sidiemeni>	The Syntax Rules and General Rules of this statement should also be expressed	
					in terms of result set rather than table.	
					Solution	
-					None provided with comment.	
	GBR-P02-014	GBR-P02-011	2-Minor	P02-14-28, Effect	The rules of this clause should be expressed in terms of result set rather than	
			Technical	of opening a	table.	
				curbor	Solution	
					None provided with comment.	
	GBR-P02-015	GBR-P02-004	1-Major	<i>P02-20-01</i> ,	The restriction to support of a defined set of standard programming languages	
		GBR-P03-001	Technical	<embedded sql<="" th=""><th>should be relaxed, and where there are sets of rules that are specific to particular</th><th></th></embedded>	should be relaxed, and where there are sets of rules that are specific to particular	
				nost program>	languages, there should be an additional member of the set that for other	
					languages these things are implementation-defined.	
					Solution	
					None provided with comment.	

# USA Comments on SC32 N 1199: ISO/IEC CD 9075-2

SEQ	Cmnt	See	•	5.4		Addressed
#	ID	AISO	Severity	Reference	Description	Ву
	USA-P02-001		3-Major Editorial	P02-00.00, Foreword	The 6 th paragraph seems to be out of date. <b>Solution</b> Modify the 6 th paragraph as shown here: This sixth fifth edition cancels and replaces the fifth fourth edition (ISO/IEC 9075:2003 1999). This problem might (probably does) extend to other parts as well.	See comment
					single list in Framework.	
	USA-P02-010		4-Minor Editorial	P02-04.18.03, Known functional dependencies in a base table	In the 5 th paragraph, replace "parameteric" with "parametric" <b>Solution</b> Provided with comment.	See comment
	USA-P02-020		1-Major Technical	P02-05.04, Names and identifiers	SQL restricts a <descriptor name=""> to <simple specification="" value=""> (<literal>, <host names="" parameter="">, <sql parameter="" references="">, and <embedded variable<br="">name&gt;). We believe an <identifier> should also be a valid choice for a <descriptor name="">, just as an <identifier> is a valid choice for a <statement name&gt; and a <cursor name="">. Solution None provided with comment.</cursor></statement </identifier></descriptor></identifier></embedded></sql></host></literal></simple></descriptor>	
	USA-P02-030		1-Major Technical	P02-07.13, <query expression&gt;</query 	Sometimes it is necessary/sufficient for an application to only retrieve the first $n$ rows of a query (where $n > 0$ ). Therefore an option on <query expression=""> is necessary to indicate that only a fixed number of rows should be returned. We suggest syntax such as "FETCH FIRST n ROWS ONLY", "LIMIT TO n ROWS", or "TOP n" to accomplish this. The order of the rows that are the result of a <query expression=""> is implementation-dependent and therefore two executions of the same <query expression="">s including this new functionality would be non-deterministic. Thus, another option is needed to order the result before the first $n$ rows are retrieved. This option would be the <order by="" clause=""> as it is currently defined for a <cursor specification="">. Of course, specifying either one or both of the new options makes the <query expression=""> read-only (<i>i.e.</i>, not updatable). Many times only the very first row is important. For this case, additional syntactic sugar would be appropriate.</query></cursor></order></query></query></query>	

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					This could be accomplished by modifying the syntax of <query expression=""> as follows: <query expression=""> ::= [ <with clause=""> ] <query body="" expression=""> [ [ <order by="" clause=""> ] <row clause="" counting=""> ] <row clause="" counting=""> ::= Syntax to be determined Solution None provided with comment.</row></row></order></query></with></query></query>	
	USA-P02-035		1-Major	P02-08.06,	The <similar predicate=""> was proposed in 1988 in a paper known as CPH-37a.</similar>	
	0511102 055		Technical	<similar predicate&gt;</similar 	Over the next few years, the predicate was refined into what we see today, but the basis for its design was always the regular expression syntax specified in Unix and in the Posix standard that was then popular.	
					Unfortunately, in later years, the Posix standard has been viewed as not very successful, and its regular expression design has not stood the test of time. Most modern standards and other specifications (notably the W3C's XML Schema and, by extension, XQuery) have chosen to use a different regular expression "standard": that of the Perl language (version 5 seems most popular). The differences between Posix regular expressions and Perl regular expressions are not huge, but they are important differences. It is undoubtedly true that SQL implementations have faithfully implemented the <similar predicate=""> as specified, using Posix-style regular expressions. However, there is visible market demand for Perl-style regular expressions.</similar>	
					that uses Perl-style regular expressions (or a corresponding subset thereof), or an additional predicate (or even more than one predicate) should be defined to provide Perl-style regular expressions, or the syntax of the <similar predicate=""> should be enhanced to allow an application program to specify which of the two styles of regular expressions it wishes to use. Solution None provided with comment.</similar>	
	USA-P02-040		1-Major Technical	P02-11.12, <alter column<br="">definition&gt;</alter>	While it is possible to modify the properties of identity columns, similar functionality is missing for generated columns. Of course, consideration would have to be given to possible side effects, such as what would happen to values of generated columns stored in existing rows in tables. <b>Solution</b> None provided with comment.	

#### INCITS H2-2005-084R1 = WG3:TXL-040 — USA COMMENTS ON CD 9075-1, -2, -3, -4, -9, -10, -11, AND -13

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Bv
	USA-P02-050		1-Major Technical	P02-11.12, <alter column<br="">definition&gt;</alter>	It should be possible to modify the data type of a column under certain conditions. One criterion should be that the existing and new data types are compatible. Another one should be that the size, precision and scale, where applicable, of the new data type are at least as large as the ones of the existing data type. For example, it should be possible to increase the size of a VARCHAR column (on the other hand, decreasing the size should not be allowed). Other possible restrictions need to be considered as well. <b>Solution</b> None provided with comment.	
	USA-P02-060		1-Major Technical	<i>P02-11.39,</i> <trigger <i>definition&gt;</i></trigger 	Not all views are updatable. However every view could be made updatable if an <i>INSTEAD OF</i> trigger were defined for it. Similar to the existing BEFORE and AFTER triggers on base tables, an INSTEAD OF trigger could be defined for views for INSERT, UPDATE, and DELETE operations. However, for INSTEAD OF triggers it is sufficient to be on a per row basis and no more than one such trigger for each action of insert, update, and delete. Solution None provided with comment.	
	USA-P02-070		4-Minor Editorial	P02-13.04, Calls to an <externally- invoked procedure&gt;</externally- 	Modify General Rule 3) e) as shown here: e) If <i>DT</i> identifies INT, DEC, or REAL and the caller language of <i>EP</i> is M, then a reference to <i>PN</i> that assigns some value <i>SV</i> to <i>PN</i> implicitly assigns the value CAST ( <i>SV</i> AS CHARACTER VARYING( <i>ML</i> )) to <i>PI</i> , where <i>ML</i> is the implementation-defined maximum length of variable- length of character strings. <b>Solution</b> Provided with comment.	See comment
	USA-P02-080		1-Major Technical	P02-14.09, <merge statement&gt;</merge 	Currently, the syntax of MERGE statement allows either updating the existing rows in the target table or inserting new rows into the target table, by branching to either <merge clause="" matched="" when=""> or <merge clause="" matched="" not="" when=""> based on the truth value of a <search condition="">. We believe it would be useful to allow multiple insert/update operations based on the truth value of additional <search condition=""> in both <merge clause="" matched="" when=""> and <merge not<br="" when="">matched clause&gt;. Solution None provided with comment.</merge></merge></search></search></merge></merge>	
	USA-P02-090		3-Major Editorial	P02-14.11, <update statement: searched&gt;</update 	Modify Conformance Rule 1) as shown here: 1) Without Feature F781, "Self-referencing operations", conforming SQL language shall not contain an <update positioned="" searched="" statement:=""> in which a leaf generally underlying table of <i>T</i> is an underlying table of any <query expression=""> generally contained in the <search condition="">.</search></query></update>	See comment

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	Ву
					Solution	
					Provided with comment.	
	USA-P02-100		3-Major	P02-19.06,	General Rule 6) contains the phrase:	See comment
			Editorial	<prepare statement&gt;</prepare 	the implementation-defined maximum <del>value of <length> for the</length></del> CHARACTER VARYING data type	
					The standard usually says:	
					implementation-defined maximum length of variable-length character strings	
					Solution	
					Provided with comment.	
	USA-P02-110		3-Major Editorial	P02-19.06, <prepare statement&gt;</prepare 	There seems to be a major problem with GR 6) a) xv) that starts off "If <i>DP</i> is the <cast operand="">" and GR 6) a) xvi) that starts off "The General Rules of Subclause 14.22", perhaps caused by an editorial mishap somewhere along the way. Our research indicates that GR 6) a) xvi) was actually part of GR 6) a) xv) until 5WD-02-Foundation-2002-12R1.pdf, but got split off into a separate rule (with the addition of some possibly spurious text) by the time 5WD-02- Foundation-2003-12R1.pdf came along. Could the Editor please research this</cast>	
					bug further and restore the correct rule?	
					Solution	
					None provided with comment.	
	USA-P02-150		3-Major Editorial	P02-22.01, <get diagnostics statement&gt;</get 	Modify the lead-in of GR 6) a) ii) 1) as shown here (see also SIA-025r1): 1) Let <b>S</b> be	See comment
					Solution	
	$USA_P02_120$		4-Minor	P02-22 01 < get	Modify Syntax Rule 2) as shown here:	See comment
	057-102-120		Editorial	diagnostics statement>	<ul> <li>2) The declared type of <all info="" target=""> shall be a character string type.</all></li> </ul>	See comment
					Solution	
			4 14:00 - 00	P02 22 01 (act	Provided with comment.	Can annual t
	USA-P02-130		4-Minor Editorial	P02-22.01, < get diagnostics statement>	General Rule 6) contains the phrase:         implementation-defined maximum value for the <length> contained in         a <data type="">         The standard usually says:         implementation-defined maximum length of variable-length         character strings</data></length>	See comment
					Solution	
			4.3.6	D02 22 01	Provided with comment.	
	USA-P02-140		4-Minor	P02-22.01, <get< td=""><td>Delete one of the two "where"s in GR 6) a) i).</td><td>See comment</td></get<>	Delete one of the two "where"s in GR 6) a) i).	See comment
			Editorial	augnosites	Solution	

#### INCITS H2-2005-084R1 = WG3:TXL-040 — USA COMMENTS ON CD 9075-1, -2, -3, -4, -9, -10, -11, AND -13

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
				statement>	Provided with comment.	
	USA-P01-999		1-Major	P02-No specific	All Possible Problems and Editor's Notes must be satisfactorily resolved and all	
			Technical	location	problems discovered during the course of the ballot resolution process must be	
					satisfactorily resolved.	
					Solution	
					None provided with comment.	
	USA-P02-180		1-Major	P02-No specific	SQL should provide a way to insert/update/delete rows from a table and to	
			Technical	location	retrieve the rows that were inserted/updated/deleted as a single operation. See	
					also Language Opportunity FND-849, which asks for a similar functionality.	
					Solution	
					None provided with comment.	
	USA-P02-190		1-Major	P02-No specific	Applications have a need to temporarily disable (table) constraints (across	
			Technical	location	transaction boundaries).	
					For example, a user may want to disable constraints because he is going to be	
					reloading data from multiple sites and he does not want to go through the work	
					of making sure that the data across N tables is loaded in the appropriate order.	
					Though the standard has the ability to defer constraint checking until the end of a	
					transaction this would not be a complete solution as most users would not want	
					to perform the entire set of loads all under one gigantic transaction and in the	
					case of loads from multiple data sources, it is likely impractical if not	
					impossible.	
					So the standard should provide syntax to let the user specify whether a given	
					constraint is enforced or not enforced.	
					Solution	
					None provided with comment.	
	USA-P02-200		1-Major	P02-No specific	Many applications have a need to deal with binary data. Not always are	
			Technical	location	BINARY LARGE OBJECTs (BLOBs) appropriate for these applications since	
					BLOBs have certain restrictions, such as they are not allowed in <general set<="" td=""><td></td></general>	
					function>, <group by="" clause="">, and <order by="" clause="">, etc. For such applications,</order></group>	
					it would be beneficial if the standard supports a "regular" binary string type that	
					does not have the restrictions associated with BLOB type. Just as the standard	
					has the kinds of character string types: CHAK, VAKCHAK, and CLOB, we	
					VAPRIMARY, and RLOR (all measure in lengths of octats)	
					Note that this new data type needs to avoid the problems that led to the removal	
					of the BIT and BIT VARYING data types	
					Solution	
					None provided with comment.	
	USA-P02-210		1-Major	P02-No specific	Exact numeric types are used when exact results are desired for arithmetic	
	USA-P02-200 USA-P02-210		1-Major Technical	P02-No specific location P02-No specific	Though the standard has the ability to defer constraint checking until the end of a transaction this would not be a complete solution as most users would not want to perform the entire set of loads all under one gigantic transaction and in the case of loads from multiple data sources, it is likely impractical if not impossible. So the standard should provide syntax to let the user specify whether a given constraint is enforced or not enforced. <b>Solution</b> None provided with comment. Many applications have a need to deal with binary data. Not always are BINARY LARGE OBJECTs (BLOBs) appropriate for these applications since BLOBs have certain restrictions, such as they are not allowed in <general function="" set="">, <group by="" clause="">, and <order by="" clause="">, etc. For such applications, it would be beneficial if the standard supports a "regular" binary string type that does not have the restrictions associated with BLOB type. Just as the standard has three kinds of character string types: CHAR, VARCHAR, and CLOB, we believe the standard should support three kinds of binary types: BINARY, VARBINARY, and BLOB (all measure in lengths of octets). Note that this new data type needs to avoid the problems that led to the removal of the BIT and BIT VARYING data types. <b>Solution</b> None provided with comment. Exact numeric types are used when exact results are desired for arithmetic</order></group></general>	

#### USA COMMENTS ON CD 9075-1, -2, -3, -4, -9, -10, -11, AND -13 — INCITS H2-2005-084R1 = WG3:TXL-040

SEQ	Cmnt	See				Addressed
#	ID	Also	Severity	Reference	Description	By
			Technical	location	operations, but they suffer from the fact that the range of values supported by these types is much smaller than the range supported by approximate numeric types. The upcoming revision of ANSI/IEEE standard 754, "IEEE Standard for Binary Floating-Point Arithmetic", includes a new kind of exact numeric type called "Decimal Floating Point" that offers a much bigger range of values for a given precision while providing exact results for arithmetic operations. Many popular programming languages are in the process of adding this new type to their type systems. We believe SQL should also add this new type to the list of predefined types. Solution None provided with comment.	
	USA-P02-220		1-Major Technical	P02-No specific location	When a user is granted one or more roles, he may expect to be allowed to access not only those objects on which he has been granted privileges, but also the privileges granted to every role with which he is associated. Unfortunately, this is not currently the case in SQL. When a user starts an SQL-session, he is not allowed to access any of the objects on which a role he has been granted has privileges (unless, of course, he has the privilege directly). If he wants to access objects on which a role he has been granted has privileges, he has to first perform a SET ROLE statement. It would be useful to have the ability to specify that a ROLE is enabled by default. This would prevent a user from having to execute a SET ROLE statement to get the an initial ROLE. Solution None provided with comment.	
	USA-P02-225		1-Major Technical	P02-No specific location	The current definition of ROLEs supports the ability to use a SET ROLE statement for only one role at a time. There are situations where it would be useful to allow a user to enable multiple roles concurrently. <b>Solution</b> None provided with comment.	
	USA-P02-235		1-Major Technical	P02-No specific location	Currently it is possible to grant a user both a "Teller" role and an "Auditor" role even though the two roles would be considered mutually exclusive in most organizations. Similarly, it is possible to grant a user an "Auditor" role even if that user already has the "Teller" role. We believe that SQL should provide a mechanism to prevent a user from acquiring or setting such mutually exclusive roles, similar to what the ANSI standard on Role Based Access Control, ANSI INCITS 359:2004 refers to as "static separation of duty relations". <b>Solution</b> None provided with comment.	
	USA-P02-240		Technical	location	users to manipulate entire database objects. However, SQL does not provide any	

### INCITS H2-2005-084R1 = WG3:TXL-040 — USA COMMENTS ON CD 9075-1, -2, -3, -4, -9, -10, -11, AND -13

SEQ #	Cmnt ID	See Also	Severity	Reference	Description	Addressed Bv
					sort of mandatory access control—that is, there is no way in the current security model to restrict access to subsets of data in database objects such as tables, based on labels such as "Top Secret", "Board of Directors only", etc. only to those users that possess appropriate authorizations. Such an extended security model is a requirement in many organizations. We believe the security model of SQL should be extended to offer such a capability. <b>Solution</b> None provided with comment.	_,
	USA-P02-160		3-Major Editorial	P02-No specific location	<ul> <li>Subclause 9.8, "Determination of identical values", General Rule 2) reads (with emphasis added by the comment author):</li> <li>2) Case: <ul> <li>a) If V1 and V2 are both null, then V1 is identical to V2.</li> <li>b) If V1 is null and V2 is not null, or if V1 is not null and V2 is null, then V1 is not identical to V2.</li> <li>c)</li> </ul> </li> <li>This GR and other rules in Foundation and other parts use the terms "is/are null", "is not null", etc. loosely. What is preferable is to use the terms "is the null value", etc. instead.</li> <li>One needs to carefully examine Foundation and possibly all other parts to replace all these offending phrases.</li> </ul> Solution	

<<< End >>>